

A new lower bound on the independence number of a graph

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Abstract

For a given connected graph G on n vertices and m edges, we prove that its independence number $\alpha(G)$ is at least $((2m+n+2) - ((2m+n+2)^2 - 16n^2)^{1/2})/8$.

Keywords : independence number, min algorithm, connected graph.

Introduction

Let $G=(V,E)$ be a connected graph G on $n=|V|$ vertices and $m=|E|$ edges.

For a subgraph H of G and for a vertex $i \in V(H)$, let $d_H(i)$ be the degree of i in H and let $N_H(i)$ be its neighbourhood in H . Let $\delta(H)$ and $\Delta(H)$ be the minimum degree and the maximum degree of H , respectively. A subset X of V is called independent if its vertices are mutually non-adjacent. The independence number $\alpha(G)$ is the largest cardinality among all independent sets of G .

The problem of finding an independent set of maximum cardinality is known to be NP –complete[1]. Some approximation algorithms were designed to tackle this problem, among them, the well known MIN algorithm [4], which can be implemented in time linear in n and m :

$G_1 := G, j := 1$

While $V(G_j) \neq \emptyset$ do

 Begin

 Choose $i_j \in V(G_j)$ with $d_{G_j}(i_j) = \delta(G_j)$, delete $\{i_j\} \cup N_{G_j}(i_j)$ to obtain G_{j+1} and set

$j := j+1$;

 End ;

$k := j-1$

Let k_{MIN} be the smallest k the algorithm MIN provides for a given connected graph G .

Harant [3] proved that $\alpha(G) \geq k_{\text{MIN}} \geq ((2m+n+1) - ((2m+n+1)^2 - 4n^2)^{1/2})/2$.

The purpose of the present note is to improve this lower bound.

Lemma : For a given connected graph G on n vertices and m edges,

$$\alpha(G) \geq ((2m+n+2) - ((2m+n+2)^2 - 16n^2)^{1/2})/8.$$

The proof starts with the inequality (1) proved by Harant [3] :

$$k_{\text{MIN}} \geq n^2 / (2m+n - \sum_{i \in I_n} (d_G(i) - \delta(G_i))) \tag{1}$$

and uses a variation of the one given by Halldorson [2].

For $j=1, \dots, k_{\text{MIN}}$, let $d_{G_j}(i_j)$ be the degree in the remaining graph of the j -th vertice choosed at the j -th iteration of the algorithm MIN. The number of vertices deleted in the j -th iteration is thus $1 + d_{G_j}(i_j)$ and the sum of the degrees of the $1 + d_{G_j}(i_j)$ vertices deleted is at least $(1 + d_{G_j}(i_j))d_{G_j}(i_j)$. Thus the number of edges removed in the j -th iteration is at least $(1 + d_{G_j}(i_j))d_{G_j}(i_j)/2$.

Let X be an independent set of G of maximum cardinality α , and let k_j be the number of vertices among the $1 + d_{G_j}(i_j)$ vertices deleted in the j -th iteration that are also contained in X .

$$\sum_{j=1}^{k_{\text{MIN}}} k_j = \alpha$$

Since X is edgless, and G is connected then the number of edges removed in the j -th iteration ($j=1, \dots, k_{\text{MIN}} - 1$) is at least :

$$\binom{1 + d_{G_j}(i_j)}{2} + \binom{k_j}{2} + 1$$

(for $j=1, \dots, k_{\text{MIN}} - 1$, there is at least one edge between $N_{G_j}(i_j)$ and G_{j+1} , because G is supposed connected).

In the k_{MIN} -th iteration, at least

$$\binom{1 + d_{G_j}(i_{k_{\text{MIN}}})}{2} + \binom{k_{k_{\text{MIN}}}}{2}$$

edges are removed.

Hence we obtain the following inequality :

$$m \geq \sum_{j=1}^{k_{\text{MIN}}-1} \left(\binom{1 + d_{G_j}(i_j)}{2} + \binom{k_j}{2} + 1 \right) + \binom{1 + d_{G_j}(i_{k_{\text{MIN}}})}{2} + \binom{k_{k_{\text{MIN}}}}{2}$$

then :

$$2m \geq 2k_{\text{MIN}} - 2 + \sum_{j=1}^{k_{\text{MIN}}} \left((1 + d_{G_j}(i_j)) d_{G_j}(i_j) \right) + \sum_{j=1}^{k_{\text{MIN}}} k_j + \sum_{j=1}^{k_{\text{MIN}}} (k_j)^2$$

consequently :

$$2m \geq 4k_{\text{MIN}} - 2 + \sum_{j=1}^{k_{\text{MIN}}} \left((1 + d_{G_j}(i_j)) d_{G_j}(i_j) \right) \tag{2}$$

On the other hand :

Since $\forall (j, j') \in \{1, \dots, k_{\text{MIN}}\}, j \neq j' \Rightarrow (\{i_j\} \cup N_{G_j}(i_j)) \cap (\{i_{j'}\} \cup N_{G_{j'}}(i_{j'})) = \emptyset$ and

$$I_n = \bigcup_{j=1}^{k_{\text{MIN}}} (\{i_j\} \cup N_{G_j}(i_j)) = \{1, \dots, n\}$$

then

$$\sum_{i \in I_n} \delta(G_j) = \sum_{j=1}^{k_{\text{MIN}}} \sum_{i \in \{i_j\} \cup N_{G_j}(i_j)} \delta(G_j)$$

and

$$\sum_{i \in I_n} \delta(G_i) = \sum_{j=1}^{k_{\text{MIN}}} (1 + d_{G_j}(i_j)) \delta(G_j) \leq \sum_{j=1}^{k_{\text{MIN}}} ((1 + d_{G_j}(i_j)) d_{G_j}(i_j))$$

thus

$$\sum_{i \in I_n} (d_G(i) - \delta(G_i)) = 2m - \sum_{i \in I_n} \delta(G_i) \geq 2m - \sum_{j=1}^{k_{\text{MIN}}} ((1 + d_{G_j}(i_j)) d_{G_j}(i_j))$$

by using inequality (2) we have :

$$\sum_{i \in I_n} (d_G(i) - \delta(G_i)) \geq 4k_{\text{MIN}} - 2$$

then inequality (1) implies :

$$k_{\text{MIN}} \geq n^2 / (2m + n + 2 - 4k_{\text{MIN}})$$

and consequently :

$$k_{\text{MIN}} \geq ((2m + n + 2) - ((2m + n + 2)^2 - 16n^2)^{1/2}) / 8.$$

Conclusion

This note presented an improved lower bound on the independence number of a connected graph, and as a future work, it will be proved that this bound is optimale for a particular class of graphs.

References :

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