

Lifted Inequalities for 0–1 Mixed-Integer Bilinear Covering Sets*

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March 1, 2011

Abstract

In this paper, we study 0–1 mixed-integer bilinear covering sets. We derive several families of facet-defining inequalities via sequence-independent lifting techniques. We then show that these sets have polyhedral structures that are similar to those of certain fixed-charge single-node flow sets. As a result, we obtain new facet-defining inequalities for these sets that generalize well-known lifted flow cover inequalities from the integer programming literature.

1 Introduction and motivation

Nonlinear branch-and-bound is a method to solve mixed-integer nonlinear programming (MINLP) problems to global optimality; see [9, 17]. This method has been implemented in commercial solvers such as BARON [27] and LINDO Global [18]. It requires that convex relaxations of the problem be recursively solved over smaller and smaller subsets of the feasible region obtained by branching on variables. Most existing commercial software use a method proposed by McCormick [21] to obtain these convex relaxations for factorable problems. McCormick’s relaxation is an instantiation of a more general technique that relaxes (nonconvex) constraints of the form $g(x) \geq r$ into (convex) constraints of the form $\bar{g}(x) \geq r$ where $\bar{g}(x)$ is a concave overestimator of $g(x)$. This technique does not use the right-hand-side of the inequality in the process. As a result, the relaxation obtained is typically not the strongest possible.

Some of the functional forms that appear most frequently in the formulation of nonlinear programs are probably multilinear inequalities and equalities. In particular, bilinear inequalities of the covering type

$$\sum_{j \in N} a_j x_j y_j \geq d, \quad (1)$$

where $a_j > 0$, $x_j \in S \subseteq \mathbb{R}_+$, and $y_j \in S' \subseteq \mathbb{R}_+$ appear in the formulation of various practical problems (including trimloss applications; see Harjunkoski et al. [16] for an example), and are among the simplest nonconvex inequalities that can be studied. Therefore, sets of the form (1) provide an important test bed for the derivation of new, stronger convexification methods that use right-hand-side information. When variables do not have upper bounds, we have derived in [30] closed-form expressions for the convex hull of feasible solutions of (1) over various subsets of the nonnegative orthant. For problems where variables are continuous and have finite upper bounds, we also derived in [29] convex relaxations of (1) that are stronger than McCormick’s.

In this paper, we study further the convex hull of feasible solutions to (1) when variables are bounded. In particular, we consider 0–1 mixed-integer bilinear covering sets of the form

$$B = \left\{ (x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{j=1}^n a_j x_j y_j \geq d \right\},$$

*This work was supported by NSF CMMI grants 0856605 and 0900065.

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34 where $n \in \mathbb{Z}_{++}$, $a_j > 0 \forall j \in N := \{1, \dots, n\}$, and $d > 0$. Results similar to those derived in this paper can
35 also be obtained for sets defined through constraints of the form $\sum_{j=1}^k (a_j x_j y_j + b_j x_j) + \sum_{j=k+1}^n a_j y_j \geq d$.
36 This generalization allows us to extend the applicability of our study to problems where the bounds on y are
37 not 0 and 1 and, in addition, to problems where some of the x variables are fixed. Our proofs extend easily
38 to such a setup because the two sets share strong relationships that are described in Proposition 5.1 and the
39 discussion following it.

40 In order to guarantee that B is not empty, we impose

41 **Assumption 1.** $\sum_{j=1}^n a_j \geq d$.

42 On the theoretical side, we are interested in studying relaxation techniques for B that will take both
43 the right-hand-side d and upper bounds on the variables into account. On the one hand, it follows from
44 the separability of $\sum_{j \in N} a_j x_j y_j$ over j that $\text{conv}\{(x, y, z) \in \{0, 1\}^n \times [0, 1]^n \times \mathbb{R} \mid z \leq \sum_{j \in N} a_j x_j y_j\}$ is
45 described by the McCormick constraints that overestimate each bilinear term separately [1]. Therefore, the
46 tightest relaxation of the type $\bar{g}(x) \geq d$, where $\bar{g}(x)$ is a concave overestimator of $\sum_{j \in N} a_j x_j y_j$ restricted
47 to $\{0, 1\}^n \times [0, 1]^n$ over $[0, 1]^{2n}$, is the relaxation described above that uses McCormick constraints. On the
48 other hand, if upper bounds on the variables are absent, the convex hull of the bilinear covering set can be
49 obtained explicitly [30]. Yet, as we will see in Proposition 1.1, it is difficult to optimize linear functions over B
50 and therefore the study of PB will help us understand better the difficulties that arise from the simultaneous
51 presence of a right-hand-side and upper-bounds on the variables.

52 On the practical side, we are interested in deriving convex relaxations of B since they directly yield convex
53 relaxations for problems with constraints of the form $\sum_{j=1}^n f_j(z) x_j \geq d$, where $z \in \mathbb{R}^p$, by replacing $f_j(z)$ with
54 $a_j y_j + b_j$ where $y_j \in [0, 1]$. We are also interested in studying B because of its relations to some important
55 mixed-integer linear sets. In particular, since the set B is a relaxation of the fixed-charge single-node flow
56 set without inflows

$$57 \quad F = \left\{ (x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{j=1}^n a_j y_j \geq d, x_j \geq y_j \forall j \in N \right\},$$

58 see Lemma 4.1 for a proof, valid inequalities for B will also be valid for F . Further, we will show in Section 4
59 that facets of either F or B can be easily identified if facet-defining inequalities for the other set are known.
60 As a result, the inequalities we derive for B also reveal new families of facet-defining inequalities for the
61 convex hull of F .

62 We next argue that it is typically difficult to find globally optimal solutions to problems containing B as
63 a constraint by showing that it is NP-hard to optimize a linear function over B . To this end, consider the
64 following optimization problem (Q) that seeks to minimize a linear objective function over the bilinear set
65 B :

$$66 \quad (Q) \quad \min \left\{ \sum_{j=1}^n b_j x_j + \sum_{j=1}^n c_j y_j \mid (x, y) \in B \right\}$$

67 where $b \in \mathbb{R}^n$ and $c \in \mathbb{R}^n$.

68 **Proposition 1.1.** *Problem (Q) is NP-hard.*

69 *Proof.* The proof is by reduction from the 0–1 knapsack problem, which is proven to be NP-hard in [10].
70 Consider the following 0–1 knapsack instance:

$$71 \quad (K) \quad z^K = \min \left\{ \sum_{j=1}^n b_j x_j \mid \sum_{j=1}^n a_j x_j \geq d, x_j \in \{0, 1\} \forall j \in N \right\}.$$

72 We define a corresponding instance of (Q) by setting $c_j = -1$ for all $j \in N$, i.e.

$$73 \quad (P) \quad z^P = \min \left\{ \sum_{j=1}^n b_j x_j - \sum_{j=1}^n y_j \mid \sum_{j=1}^n a_j x_j y_j \geq d, x_j \in \{0, 1\}, y_j \in [0, 1] \forall j \in N \right\}.$$

74 The reduction from (K) to (P) is clearly polynomial. Observe further that if x^* is a feasible solution to (K) ,
75 then $(x^*, \mathbf{1})$ is feasible to (P) , therefore showing that $z^P \leq z^K - n$. Similarly, if (x^*, y^*) is an optimal solution
76 to (P) , then x^* is feasible to (K) as $\sum_{j=1}^n a_j x_j^* \geq \sum_{j=1}^n a_j x_j^* y_j^* \geq d$. Therefore $z^K \leq z^P + \mathbf{1}^\top y^* \leq z^P + n$.
77 We conclude that $z^P = z^K - n$ and that x^* is an optimal solution to (K) if and only if $(x^*, \mathbf{1})$ is an optimal
78 solution to (P) . □

80 In this paper, we are interested in studying the convex hull of B , $\text{conv}(B)$, that we denote by PB . Since
81 B is a finite union of polytopes, PB is polyhedral.

82 **Proposition 1.2.** *PB is a polytope.*

83 It follows that, when studying PB , it is sufficient to consider linear inequalities. Proposition 1.1 suggests
84 that finding a complete closed-form expression for the convex hull of B is difficult. As a result, we will
85 focus our efforts on constructing families of strong cutting planes for optimization problems containing the
86 constraints of B by studying the convex hull of B . To construct these inequalities, we will use lifting. Lifting is
87 a well-known integer programming technique that generates strong inequalities for a given set by transforming
88 an inequality valid for a restricted subset of the feasible region into a globally valid constraint. Early work on
89 lifting in integer programming can be found in Wolsey [32, 33]. A generalization to nonlinear programming
90 is given in Richard and Tawarmalani [24]. In particular, lifting is said to be sequence-independent if the
91 order in which the restrictions are removed does not change the derived inequality. Subadditivity of a certain
92 perturbation function, called the lifting function, is a sufficient condition for lifting to be sequence-independent
93 when the restrictions involve fixing the variables at their bounds; see Proposition 3.2 and [24]. In this paper,
94 we derive new tools to verify that functions are subadditive that we exploit to derive large families of facet-
95 defining inequalities for PB . These results illustrate that lifting can successfully use bounds on variables in
96 the generation of cuts for MINLPs. Further, the results have implications for fixed-charge flow models, a
97 family of problems both theoretically and practically important in mixed-integer linear programming.

98 The paper is structured as follows. In Section 2, we derive basic polyhedral results about PB . We provide
99 necessary and sufficient conditions for trivial inequalities to be facet-defining. Then, we derive a linear
100 description of PB for the special case where $n = 2$. This result is used to identify the seed inequalities that
101 will be used in lifting procedures. In Section 3, we show that for a general class of multi-dimensional functions,
102 it suffices to check the subadditivity condition at certain points to establish the subadditivity of the function
103 everywhere. Then, using this result, we derive, in closed-form, three families of facet-defining inequalities for
104 PB using sequence-independent lifting techniques. One requires the use of a subadditive approximation of the
105 lifting function. In Section 4, we prove that there are some tight connections between the facets of PB and
106 those of PF . In particular, we show that the lifted inequalities developed for PB generalize certain families
107 of flow cover cuts and yield new facet-defining inequalities for the fixed-charge single-node flow set without
108 inflows F . We summarize the contributions of our work and conclude with directions of future research in
109 Section 5.

110 2 Basic polyhedral results

111 In this section, we derive basic results about the polyhedral structure of PB . First, we provide necessary and
112 sufficient conditions for PB to be full-dimensional.

113 **Proposition 2.1.** *PB is a full-dimensional polytope if and only if $\sum_{j=1}^n a_j - a_i \geq d$ for all $i \in N$.*

114 *Proof.* First, we show that if $\sum_{j=1}^n a_j - a_i \geq d$ for all $i \in N$, then PB is full-dimensional. For all $i \in N$,
115 construct $p^i = (\mathbf{1} - e_i, \mathbf{1})$ and $q^i = (\mathbf{1}, \mathbf{1} - e_i)$. Also define $r = (\mathbf{1}, \mathbf{1})$. The points p^i , q^i , and r belong to
116 B . These points are affinely independent because $r - p^i$ and $r - q^i$ for all $i \in N$ are linearly independent.
117 Since we have described $2n + 1$ affinely independent points in PB , we have shown that PB is full-dimensional.
118 Next, we prove that if PB is a full-dimensional polyhedron, then $\sum_{j=1}^n a_j - a_i \geq d$ for all $i \in N$. Assume by
119 contradiction that $\sum_{j=1}^n a_j - a_i < d$ for some $i \in N$. Since $\sum_{j=1}^n a_j \geq d$ from Assumption 1, B is nonempty
120 and so $x_i = 1$ in every feasible solution of B , showing that PB is not full-dimensional. This is the desired
121 contradiction. □

122 In the remainder of this paper, we will assume that PB is full-dimensional.

123 **Assumption 2.** $\sum_{j=1}^n a_j - a_i \geq d$ for all $i \in N$.

124 Observe that Assumption 2 strictly dominates Assumption 1 and implies that $n \geq 2$. We next identify
125 some basic characteristics of the facet-defining inequalities of PB .

126 **Proposition 2.2.** *Let*

$$127 \quad \sum_{j=1}^n \alpha_j x_j + \sum_{j=1}^n \beta_j y_j \geq \delta \quad (2)$$

128 *be a facet-defining inequality for PB that is not a scalar multiple of $x_i \leq 1$ for $i \in N$ or $y_i \leq 1$ for $i \in N$.
129 Then, (i) $\alpha_i \geq 0$, $\forall i \in N$, (ii) $\beta_i \geq 0$, $\forall i \in N$, and (iii) $\delta \geq 0$.*

130 *Proof.* Select $i \in N$. Since (2) is a facet-defining inequality for PB that is not a scalar multiple of $x_i \leq 1$,
131 there exists $(x^*, y^*) \in B$ with $x_i^* < 1$ such that

$$132 \quad \sum_{j=1}^n \alpha_j x_j^* + \sum_{j=1}^n \beta_j y_j^* = \delta. \quad (3)$$

133 Consider now $(\bar{x}, \bar{y}) = (x^*, y^*) + (1 - x_i^*)(e_i, 0)$. This point belongs to B and therefore satisfies (2), *i.e.*,

$$134 \quad \sum_{j=1}^n \alpha_j \bar{x}_j + \sum_{j=1}^n \beta_j \bar{y}_j \geq \delta. \quad (4)$$

135 Subtracting (3) from (4), we obtain that $\alpha_i \geq 0$. The proof that $\beta_i \geq 0$ for all $i \in N$ is similar. The fact that
136 $\delta \geq 0$ then follows from (3) after noting that all terms in the left-hand-side are nonnegative. \square

137 The following proposition further studies facet-defining inequalities whose right-hand-sides are zero.

138 **Proposition 2.3.** *Let*

$$139 \quad \sum_{j=1}^n \alpha_j x_j + \sum_{j=1}^n \beta_j y_j \geq 0 \quad (5)$$

140 *be a facet-defining inequality for PB . Then, (5) is a scalar multiple of $x_j \geq 0$ for $j \in N$ or of $y_j \geq 0$ for
141 $j \in N$.*

142 *Proof.* Assume for a contradiction that (5) is not a scalar multiple of $x_j \geq 0$ for $j \in N$ or of $y_j \geq 0$ for $j \in N$.
143 Then, for each $i \in N$, there exists $(x^i, y^i) \in B$ such that $x_i^i > 0$ and for which

$$144 \quad \sum_{j=1}^n \alpha_j x_j^i + \sum_{j=1}^n \beta_j y_j^i = 0. \quad (6)$$

145 Since we know from Proposition 2.2 that $\alpha_j \geq 0$ and $\beta_j \geq 0$ for all $j \in N$, we obtain from (6) that

$$146 \quad 0 = \sum_{j=1}^n \alpha_j x_j^i + \sum_{j=1}^n \beta_j y_j^i \geq \alpha_i x_i^i \geq 0. \quad (7)$$

147 We conclude that, for each $i \in N$, $\alpha_i = 0$ since $x_i^i > 0$. Similarly, we can establish that $\beta_i = 0 \forall i \in N$. This
148 is a contradiction to the fact that (5) is facet-defining for PB . \square

149 We now focus on these inequalities that play a special role in Propositions 2.2 and 2.3 and characterize
150 when they are facet-defining for PB . We refer to these inequalities as *bound inequalities*.

151 **Proposition 2.4.** *The upper bound inequalities $x_i \leq 1$, $y_i \leq 1$ are facet-defining for PB for all $i \in N$.
152 Further, for $i \in N$, the lower bound inequalities $x_i \geq 0$, $y_i \geq 0$ are facet-defining for PB if and only if
153 $\sum_{j=1}^n a_j - a_i - a_{l(i)} \geq d$ where $l(i) \in \operatorname{argmax}\{a_j \mid j \in N \setminus \{i\}\}$.*

154 *Proof.* The validity of all these inequalities is trivial since they belong to the description of B . Assume for a
155 contradiction that $x_i \leq 1$ is not facet-defining for PB . Then, it follows from Proposition 2.2 that $(e_i, 0)$ is a
156 recession direction of PB , a contradiction to the fact that PB is a polytope; see Proposition 1.2. The proof
157 that $y_i \leq 1$ is facet-defining for PB is similar.

158 Now, we show that $x_i \geq 0$ is facet-defining if $\sum_{j=1}^n a_j - a_i - a_{l(i)} \geq d$ by describing $2n$ affinely independent
159 points in B that satisfy $x_i = 0$. For $k \in N \setminus \{i\}$, we construct the $2(n-1)$ points, $\bar{p}^k = (\mathbf{1} - e_i - e_k, \mathbf{1} - e_i - e_k)$
160 and $\bar{q}^k = (\mathbf{1} - e_i - e_k, \mathbf{1} - e_i)$. We also define $r^1 = (\mathbf{1} - e_i, \mathbf{1} - e_i)$ and $r^2 = (\mathbf{1} - e_i, \mathbf{1})$. Clearly, the points
161 r^1, r^2 and \bar{p}^k, \bar{q}^k for $k \in N \setminus \{i\}$ satisfy $x_i \geq 0$ at equality and are feasible for B since $\sum_{j=1}^n a_j - a_i \geq$
162 $\sum_{j=1}^n a_j - a_i - a_k \geq \sum_{j=1}^n a_j - a_i - a_{l(i)} \geq d$. These points are affinely independent since $\bar{p}^k - r^1, \bar{q}^k - r^1,$
163 and $r^2 - r^1$ can be easily verified to be linearly independent. To prove the reverse direction, assume now
164 that $x_i \geq 0$ is facet-defining for PB . We claim that $\sum_{j=1}^n a_j - a_i - a_{l(i)} \geq d$. Assume for a contradiction
165 that $\sum_{j=1}^n a_j - a_i - a_{l(i)} < d$. This condition implies that every feasible solution (x, y) of PB with $x_i = 0$
166 also must satisfy $x_{l(i)} = 1$. As a result, the dimension of the face defined by $x_i = 0$ is less or equal to
167 $2n - 2$, which is a contradiction. Similarly, it can be proven that $y_i \geq 0$ is facet-defining for PB if and only
168 if $\sum_{j=1}^n a_j - a_i - a_{l(i)} \geq d$. \square

169 Observe that the above proofs are also valid when $y_i \in \{0, 1\}$ instead of $y_i \in [0, 1]$ for some subset $J \subseteq N$.
170 We next study another simple facet-defining inequality for PB .

171 **Proposition 2.5.** *The inequality $\sum_{j=1}^n a_j y_j \geq d$ is facet-defining for PB .*

172 *Proof.* Validity is easily verified since $\sum_{j=1}^n a_j y_j \geq \sum_{j=1}^n a_j x_j y_j \geq d$. To prove that $\sum_{j=1}^n a_j y_j \geq d$ is facet-
173 defining, we present $2n$ points (x^i, y^i) in B that satisfy $\sum_{j=1}^n a_j y_j^i \geq d$ at equality and such that the system
174 $\alpha x^i + \beta y^i = \delta$ for $i = 1, \dots, 2n$ only has solutions (α, β, δ) that are scalar multiples of $(\mathbf{0}, a, d)$. Consider the
175 $2n$ points $p^k = (\mathbf{1}, \Delta_k(\mathbf{1} - e_k))$ and $q^k = (\mathbf{1} - e_k, \Delta_k(\mathbf{1} - e_k))$ where $\Delta_k = \frac{d}{\sum_{j=1}^n a_j - a_k}$ for $k \in N$. Note that
176 because of Assumption 2, $0 < \Delta_k \leq 1$ for all $k \in N$. Clearly, p^k and q^k belong to B and satisfy $\sum_{j=1}^n a_j y_j \geq d$
177 at equality. These $2n$ points yield the system:

$$178 \quad \sum_{j=1}^n \alpha_j + \Delta_k \left(\sum_{j=1}^n \beta_j - \beta_k \right) = \delta \quad \forall k \in N, \quad (8)$$

$$179 \quad \sum_{j=1}^n \alpha_j - \alpha_k + \Delta_k \left(\sum_{j=1}^n \beta_j - \beta_k \right) = \delta \quad \forall k \in N. \quad (9)$$

180 By subtracting (8) from (9), we obtain that $\alpha_k = 0$ for $k \in N$. From (8) and the definition of Δ_k , we then
181 conclude that, for all $k, l \in N$,

$$182 \quad \sum_{j=1}^n \beta_j - \beta_k = \frac{\delta}{d} \left(\sum_{j=1}^n a_j - a_k \right) \quad \text{and} \quad \sum_{j=1}^n \beta_j - \beta_l = \frac{\delta}{d} \left(\sum_{j=1}^n a_j - a_l \right).$$

183 Subtracting these expressions yields $\beta_k - \frac{\delta}{d} a_k = \beta_l - \frac{\delta}{d} a_l$. After defining $\beta_k - \frac{\delta}{d} a_k = \theta$ for $k \in N$ and using
184 these relations in (8), we obtain that $\theta = 0$, which implies $\beta_k = \frac{\delta}{d} a_k$ for $k \in N$. Therefore, we conclude that
185 all solutions (α, β, δ) to the system (8) and (9) are scalar multiples of $(\mathbf{0}, a, d)$. \square

186 In the remainder of this paper, we will often use the term *facet* to refer to a facet-defining inequality. We
187 will also refer to inequalities $x_i \leq 1$, $y_i \leq 1$, and $\sum_{j=1}^n a_j y_j \geq d$ as *trivial facets* of PB . To illustrate the
188 richness of the polyhedral structure of PB , we present an example next. The linear inequalities describing
189 the convex hull of this set were obtained using PORTA; see Christof and Löbel [6].

190 **Example 2.6.** *Consider the 0–1 mixed-integer bilinear covering set*

$$191 \quad B = \left\{ (x, y) \in \{0, 1\}^4 \times [0, 1]^4 \mid 19x_1y_1 + 17x_2y_2 + 15x_3y_3 + 10x_4y_4 \geq 20 \right\}.$$

192 The linear description of PB has 58 inequalities that are presented in the Appendix. They include:

$$\begin{array}{rcll}
193 & 50x_1 & + 90x_3 + 45x_4 + 76y_1 + 153y_2 & \geq 135 & (10) \\
194 & 70x_1 + 90x_2 & + 27x_4 + 38y_1 & + 135y_3 & \geq 117 & (11) \\
195 & 19x_1 + 17x_2 & & + 15y_3 + 10y_4 & \geq 20 & (12) \\
196 & & 17x_2 + 15x_3 & + 19y_1 & + 10y_4 & \geq 20 & (13) \\
197 & & & 19y_1 + 17y_2 + 15y_3 + 10y_4 & \geq 20 & (14) \\
198 & 14x_1 & + 10x_3 + 5x_4 & + 17y_2 & \geq 15 & (15) \\
199 & & 12x_2 + 10x_3 + 5x_4 + 19y_1 & & \geq 15 & (16) \\
200 & & & 10x_3 + 5x_4 + 19y_1 + 17y_2 & \geq 15 & (17) \\
201 & x_1 + x_2 + x_3 & & + 10y_4 & \geq 2 & (18) \\
202 & x_1 + x_2 + x_3 + x_4 & & & \geq 2 & (19) \\
203 & x_1 & & & \geq 0 & (20) \\
204 & & & y_1 & \geq 0 & (21) \\
205 & x_1 & & & \leq 1 & (22) \\
206 & & & y_1 & \leq 1 & (23)
\end{array}$$

207 Among the inequalities in Example 2.6, we recognize the upper bound inequalities (22) and (23) that
208 are shown to be facet-defining for PB in Proposition 2.4. In this example, the lower bound inequalities
209 (20) and (21) are also facet-defining, as can be established from Proposition 2.4. Further, (14) is the trivial
210 facet-defining inequality studied in Proposition 2.5. Our goal is now to discover families of valid inequalities
211 for PB that would explain (10) – (13) and (15) – (19).

212 To derive these nontrivial facet-defining inequalities, we first study the convex hull of B when $n = 2$ with
213 the goal of identifying seed inequalities for subsequent lifting procedures. We show that the linear description
214 of PB has at most three nontrivial inequalities. In this study, Assumption 2 requires that $a_1 \geq d$ and $a_2 \geq d$.

215 **Proposition 2.7.** *Let*

$$216 \quad B^2 = \left\{ (x, y) \in \{0, 1\}^2 \times [0, 1]^2 \mid a_1x_1y_1 + a_2x_2y_2 \geq d \right\},$$

217 where $a_1 \geq d$, $a_2 \geq d$ and $d > 0$. Then,

$$218 \quad \text{conv}(B^2) = X := \left\{ (x, y) \in [0, 1]^2 \times [0, 1]^2 \mid \begin{array}{l} x_1 + x_2 \geq 1 \\ dx_1 + a_2y_2 \geq d \\ a_1y_1 + dx_2 \geq d \\ a_1y_1 + a_2y_2 \geq d \end{array} \right\}.$$

219 *Proof.* We prove the result using disjunctive programming techniques; see [5]. We define

$$\begin{array}{l}
220 \quad X_{10} := B^2 \cap \{x_1 = 1, x_2 = 0\} = \{(1, y_1, 0, y_2) \mid \frac{d}{a_1} \leq y_1 \leq 1, 0 \leq y_2 \leq 1\}, \\
\quad X_{01} := B^2 \cap \{x_1 = 0, x_2 = 1\} = \{(0, y_1, 1, y_2) \mid 0 \leq y_1 \leq 1, \frac{d}{a_2} \leq y_2 \leq 1\}, \\
\quad X_{11} := B^2 \cap \{x_1 = 1, x_2 = 1\} = \{(1, y_1, 1, y_2) \mid a_1y_1 + a_2y_2 \geq d, 0 \leq y_1 \leq 1, 0 \leq y_2 \leq 1\}.
\end{array}$$

221 It is easily verified that $\text{conv}(B^2) = \text{conv}(X_{10} \cup X_{01} \cup X_{11}) = \text{conv}(X_2 \cup X_{11})$ where $X_2 := \text{conv}(X_{10} \cup X_{01})$.
222 We first use disjunctive programming techniques to obtain a linear description of X_2 and then compute
223 $\text{conv}(B^2)$ as $\text{conv}(X_2 \cup X_{11})$. Using Theorem 2.1 in Balas [5], we write

$$224 \quad X_2 = \text{proj}_{(x,y)} \left\{ (x_1, y_1, x_2, y_2, \bar{z}_1, \bar{z}_2, \hat{z}_1, \hat{z}_2, \lambda) \mid \begin{array}{l} (x_1, y_1, x_2, y_2) = (\lambda, \bar{z}_1 + \hat{z}_1, 1 - \lambda, \bar{z}_2 + \hat{z}_2), \\ \frac{d}{a_1} \lambda \leq \bar{z}_1 \leq \lambda, 0 \leq \bar{z}_2 \leq \lambda, \\ 0 \leq \hat{z}_1 \leq 1 - \lambda, \frac{d}{a_2} (1 - \lambda) \leq \hat{z}_2 \leq 1 - \lambda, \\ 0 \leq \lambda \leq 1 \end{array} \right\}.$$

225 We then use Fourier-Motzkin elimination [34] to compute the projection. We first eliminate the variables λ ,
226 \hat{z}_1 and \hat{z}_2 using the equations $\lambda = x_1$, $\hat{z}_1 = y_1 - \bar{z}_1$, and $\hat{z}_2 = y_2 - \bar{z}_2$. We obtain

$$227 \quad x_1 + x_2 = 1, 0 \leq x_1 \leq 1,$$

228 and

$$\begin{aligned}
& \frac{d}{a_1}x_1 \leq \bar{z}_1 \leq x_1, \\
& x_1 + y_1 - 1 \leq \bar{z}_1 \leq y_1, \\
& 0 \leq \bar{z}_2 \leq 1 - x_2, \\
& y_2 - x_2 \leq \bar{z}_2 \leq y_2 - \frac{d}{a_2}x_2,
\end{aligned}$$

230 from which we project variables \bar{z}_1 and \bar{z}_2 to obtain

$$231 \quad X_2 = \text{conv}(X_{10} \cup X_{01}) = \left\{ (x_1, y_1, x_2, y_2) \left| \begin{array}{l} x_1 + x_2 = 1, x_1 \geq 0, x_2 \geq 0, \\ \frac{d}{a_1}x_1 \leq y_1 \leq 1, \frac{d}{a_2}x_2 \leq y_2 \leq 1 \end{array} \right. \right\}$$

232 since $x_1 \leq 1$ and $x_2 \leq 1$ are implied by $x_1 + x_2 = 1$, $x_1 \geq 0$ and $x_2 \geq 0$. Now, compute $\text{conv}(X_2 \cup X_{11})$ as

$$233 \quad \text{proj}_{(x,y)} \left\{ \begin{array}{l} (x_1, y_1, x_2, y_2, \bar{u}_1, \\ \bar{u}_2, \bar{v}_1, \bar{v}_2, \hat{v}_1, \hat{v}_2, \lambda) \end{array} \left| \begin{array}{l} (x_1, y_1, x_2, y_2) = (\bar{u}_1 + (1 - \lambda), \bar{v}_1 + \hat{v}_1, \bar{u}_2 + (1 - \lambda), \bar{v}_2 + \hat{v}_2), \\ \bar{u}_1 + \bar{u}_2 = \lambda, \bar{u}_1 \geq 0, \bar{u}_2 \geq 0, \\ \frac{d}{a_1}\bar{u}_1 \leq \bar{v}_1 \leq \lambda, \frac{d}{a_2}\bar{u}_2 \leq \bar{v}_2 \leq \lambda, \\ a_1\hat{v}_1 + a_2\hat{v}_2 \geq d(1 - \lambda), \\ 0 \leq \hat{v}_1 \leq 1 - \lambda, 0 \leq \hat{v}_2 \leq 1 - \lambda, \\ 0 \leq \lambda \leq 1 \end{array} \right. \right\}.$$

234 We again obtain the projection using Fourier-Motzkin elimination. Using the equations $x_1 = \bar{u}_1 + 1 - \lambda$,
235 $x_2 = \bar{u}_2 + 1 - \lambda$, and $\bar{u}_1 + \bar{u}_2 = \lambda$, we obtain that $\lambda = 2 - (x_1 + x_2)$, $\bar{u}_1 = 1 - x_2$, and $\bar{u}_2 = 1 - x_1$. Using these
236 relations together with $\bar{v}_1 = y_1 - \hat{v}_1$ and $\bar{v}_2 = y_2 - \hat{v}_2$ to eliminate the corresponding variables, we obtain

$$237 \quad x_1 \leq 1, x_2 \leq 1, 1 \leq x_1 + x_2 \preccurlyeq 2,$$

238 and

$$\begin{aligned}
& y_1 + x_1 + x_2 - 2 \leq \hat{v}_1 \leq y_1 - \frac{d}{a_1}(1 - x_2), \\
& -\frac{a_2}{a_1}\hat{v}_2 + \frac{d}{a_1}(x_1 + x_2 - 1) \leq \hat{v}_1, \\
& 0 \leq \hat{v}_1 \leq x_1 + x_2 - 1, \\
& y_2 + x_1 + x_2 - 2 \leq \hat{v}_2 \leq y_2 - \frac{d}{a_2}(1 - x_1), \\
& 0 \leq \hat{v}_2 \leq x_1 + x_2 - 1,
\end{aligned}$$

240 where inequality \preccurlyeq is clearly redundant. Projecting \hat{v}_1 , we obtain

$$241 \quad x_1 \leq 1, x_2 \leq 1, 1 \leq x_1 + x_2, y_1 \leq 1, \frac{d}{a_1}(1 - x_2) \leq y_1, \\
a_1x_1 + (a_1 - d)x_2 \preccurlyeq 2a_1 - d$$

242 and

$$\begin{aligned}
& \frac{d}{a_2}x_1 - \frac{a_1}{a_2}y_1 \leq \hat{v}_2, \\
& \frac{(d - a_1)(x_1 + x_2 - 1)}{a_2} \preccurlyeq \hat{v}_2, \\
& y_2 + x_1 + x_2 - 2 \leq \hat{v}_2 \leq y_2 - \frac{d}{a_2}(1 - x_1), \\
& 0 \leq \hat{v}_2 \leq x_1 + x_2 - 1,
\end{aligned}$$

244 where obviously redundant inequalities have been omitted. Again, inequalities \preccurlyeq are redundant since $x_1 \leq 1$,
245 $x_2 \leq 1$, $x_1 + x_2 \geq 1$, $a_1 \geq d$ and $a_2 \geq d > 0$. Projecting \hat{v}_2 , we obtain the system

$$246 \quad x_1 \leq 1, x_2 \leq 1, 1 \leq x_1 + x_2, y_1 \leq 1, \frac{d}{a_1}(1 - x_2) \leq y_1,$$

247 and

$$\begin{aligned}
& d \leq a_1y_1 + a_2y_2, \\
& a_2 \preccurlyeq (a_2 - d)x_1 + a_1y_1 + a_2x_2, \quad (R) \\
& (a_2 - d)x_1 + a_2x_2 \preccurlyeq 2a_2 - d, \\
& y_2 \leq 1, \\
& \frac{d}{a_2}(1 - x_1) \leq y_2, \\
& 1 \preccurlyeq x_1 + x_2,
\end{aligned}$$

248

249 where inequalities \preceq are either repeated or redundant. In particular, (R) is redundant since it can be obtained
 250 as a conic combination with weights $(a_2 - d)$ and 1 of valid inequalities $x_1 + x_2 \geq 1$ and $a_1 y_1 + dx_2 \geq d$.
 251 Therefore, $\text{conv}(X_2 \cup X_{11})$ is defined by bounds and the four inequalities given in the description of X ,
 252 concluding the proof. \square

253 Next, we give generalizations of the nontrivial facets of $\text{conv}(B^2)$ that we prove are facet-defining for
 254 more general instances of $\text{conv}(B)$. In particular, we give a generalization of inequalities $dx_1 + a_2 y_2 \geq d$
 255 and $a_1 y_1 + dx_2 \geq d$ in Proposition 2.9 and of inequality $x_1 + x_2 \geq 1$ in Proposition 2.11. We will use these
 256 generalizations as seed inequalities for lifting procedures in Section 3.

257 **Lemma 2.8.** *Inequality*

$$258 \quad \sum_{j \in N} \min\{dx_j, a_j x_j, a_j y_j\} \geq d \quad (24)$$

259 *is valid for PB.*

260 *Proof.* We first show that $\sum_{j \in N} \min\{dx_j, a_j x_j, a_j y_j\} \geq d$ is valid for B . Consider $(x, y) \in B$. If there exists
 261 $j \in N$ such that $dx_j < a_j x_j y_j$ then $x_j = 1$ and, consequently, the inequality is satisfied. Otherwise, the
 262 inequality reduces to the defining inequality of B . Since $(x_j, y_j) \in [0, 1]^2$ implies that $x_j y_j \leq \min\{x_j, y_j\}$ and
 263 $a_j \geq 0$, it follows that $\min\{dx_j, a_j x_j, a_j y_j\} \leq \min\{dx_j, a_j x_j, a_j y_j\}$ and, therefore, (24) is valid for PB . \square

264 The set of solutions in $[0, 1]^{2n}$ that satisfy (24) is a subset of the convex relaxations of B discussed in
 265 Section 1. In particular, when each bilinear term is outer-approximated using McCormick envelopes, we
 266 obtain the inequality $\sum_{j \in N} a_j \min\{x_j, y_j\} \geq d$, which is clearly implied by (24). Further, using orthogonal
 267 disjunctions, see [30], it can be shown that

$$268 \quad O := \text{conv}\left\{(x, y) \in \mathbb{R}_+^{2n} \mid \sum_{j \in N} a_j x_j y_j \geq d\right\} = \left\{(x, y) \in \mathbb{R}_+^{2n} \mid \sum_{j \in N} \sqrt{a_j x_j y_j} \geq \sqrt{d}\right\}.$$

269 This convex relaxation is obtained without making use of the bounds or the integrality of the variables
 270 x . It follows from the inequality relating elementary means (see Theorem 5 in [15]) that $\sqrt{da_j x_j y_j} \geq$
 271 $\min\{dx_j, a_j y_j\}$. Therefore, the feasible solutions to (24) are contained in O . However, when $(x, y) \in C \subsetneq \mathbb{R}_+^{2n}$,
 272 a procedure described in [29] permits strengthening relaxation O by restricting attention to C . When one
 273 exploits the fact that $(x, y) \in C = \{0, 1\}^n \times [0, 1]^n$, this construction yields (24).

274 **Proposition 2.9.** *Let $L \subseteq N$ be such that $\sum_{j \in N \setminus L} a_j > d$. Define $\bar{a} = \sum_{j \in N \setminus L} a_j - \max_{i \in N \setminus L} a_i$ and*
 275 *assume that $S = \{(x, \bar{y}) \in \{0, 1\}^{|L|} \times [0, 1] \mid \sum_{i \in L} \min\{a_i, d\} x_i + \bar{a} \bar{y} = d\} \neq \emptyset$. Then,*

$$276 \quad \sum_{j \in L} \min\{a_j, d\} x_j + \sum_{j \in N \setminus L} a_j y_j \geq d \quad (25)$$

277 *is facet-defining for PB. In particular, (25) is facet-defining for PB if (i) $L \cap L^\triangleright \neq \emptyset$, or (ii) $L = \emptyset$, or (iii)*
 278 *$\bar{a} \geq \max_{i \in L} \min\{a_i, d\}$, or as a special case (iv) $\bar{a} \geq d$ where $L^\triangleright := \{j \in N \mid a_j > d\}$.*

279 *Proof.* Validity of (25) for PB follows from Lemma 2.8. We now prove that (25) is facet-defining for PB by
 280 providing $2n$ affinely independent points (x^i, y^i) in B that satisfy (25) at equality. Assume without loss of
 281 generality that $L = \{1, \dots, l\}$. Define $n' = |N \setminus L|$ and denote the points as $(x_L, x_{N \setminus L}, y_L, y_{N \setminus L})$.

282 Let $(x', \bar{y}') \in S$ and define $a' = \sum_{j \in N \setminus L} a_j$. Let $p^0 = (0, \mathbf{1}, 0, \frac{d}{a'} \mathbf{1})$ and $p^j = p^0 + \epsilon(0, 0, 0, \frac{1}{a_j} e_j - \frac{1}{a_{j+1}} e_{j+1})$
 283 for $j = 1, \dots, n' - 1$. For $i \in L$, define $q^i = (e_i, \mathbf{1}, e_i, \frac{d - \min\{a_i, d\}}{a' - \min\{a_i, d\}} \mathbf{1})$, $r^i = p^0 + (0, 0, e_i, 0)$ if $a_i \leq d$, and
 284 $r^i = (e_i, \mathbf{1}, \frac{d}{a_i} e_i, 0)$ if $a_i > d$. For $j \in \{1, \dots, n'\}$, $s^j = (x'_L, \mathbf{1} - e_j, \mathbf{1}, \bar{y}' \frac{\bar{a}}{\sum_{i \in N \setminus (L \cup \{j\})} a_i} (\mathbf{1} - e_j))$. It can be
 285 easily verified that p^0, q^i, s^j and r^i belong to B and that p^j belongs to B when ϵ is sufficiently small.

286 We now show that the above points are affinely independent. Clearly, for $j \in \{0, \dots, n' - 2\}$, p^0, \dots, p^j
 287 satisfy $\sum_{i=1}^{j+1} a_i (\frac{d}{a'} - y_{l+i}) = 0$, whereas p^{j+1} does not. Therefore, p^j are affinely independent. Further, for
 288 $i \in L$, q^i are affinely independent of p^j since the latter satisfy $(x_i, y_i) = (0, 0)$. For $i \in L$, r^i are independent
 289 of p^j and q^i since the latter satisfy $y_i = x_i$. Finally, s^i are affinely independent of p^j, q^j, r^j since the latter
 290 satisfy $x_i = 1$. \square

291 The family of inequalities described in Proposition 2.9 is typically exponential in size. In the case of
 292 Example 2.6, it contains multiple inequalities including (12-14). More generally, it can be verified that
 293 inequalities (10)-(19) in the Appendix are of the form (25).

294 In the remainder of the paper, we use the following notation extensively. For $N_0, N_1 \subseteq N$ such that
 295 $N_0 \cap N_1 = \emptyset$ and $\tilde{N}_0, \tilde{N}_1 \subseteq N$ such that $\tilde{N}_0 \cap \tilde{N}_1 = \emptyset$, we let

$$296 \quad B(N_0, N_1, \tilde{N}_0, \tilde{N}_1) := \left\{ (x, y) \in B \mid \begin{array}{ll} x_j = 0 & \text{for } j \in N_0, & x_j = 1 & \text{for } j \in N_1, \\ y_j = 0 & \text{for } j \in \tilde{N}_0, & y_j = 1 & \text{for } j \in \tilde{N}_1 \end{array} \right\}.$$

297 We also define $PB(N_0, N_1, \tilde{N}_0, \tilde{N}_1) := \text{conv}(B(N_0, N_1, \tilde{N}_0, \tilde{N}_1))$. In particular, $B(\emptyset, \emptyset, \emptyset, N)$ is equivalent to
 298 the classical 0–1 knapsack set

$$299 \quad \left\{ x \in \{0, 1\}^n \mid \sum_{j=1}^n a_j x_j \geq d \right\},$$

300 whose polyhedral structure was first studied by Balas [4], Hammer et al. [14] and Wolsey [31]. The following
 301 proposition describes, among other things, relations between the bilinear set B and the 0–1 knapsack set
 302 $B(\emptyset, \emptyset, \emptyset, N)$.

303 **Proposition 2.10.** *Let*

$$304 \quad \sum_{j \in N} \alpha_j x_j + \sum_{j \in I} \beta_j y_j \geq \delta \tag{26}$$

305 *be an inequality for $PB(\emptyset, \emptyset, \emptyset, N \setminus I)$ that is not a scalar multiple of a bound inequality. Then, (26) is*
 306 *facet-defining for $PB(\emptyset, \emptyset, \emptyset, N \setminus I)$ if and only if (26) is facet-defining for PB .*

307 *Proof.* We first prove that if (26) is facet-defining for $PB(\emptyset, \emptyset, \emptyset, N \setminus I)$, then (26) is facet-defining for PB .
 308 To show that (26) is valid for B , we assume for a contradiction that there exists a point $(x', y') \in B$ with
 309 $\sum_{j \in N} \alpha_j x'_j + \sum_{j \in I} \beta_j y'_j < \delta$. Since $(x', y') \in B$, we have that $\sum_{j \in N} a_j x'_j y'_j \geq d$. Next, we define (\bar{x}, \bar{y})
 310 as $\bar{x} = x'$, $\bar{y}_j = y'_j$ for $j \in I$, and $\bar{y}_j = 1$ for $j \in N \setminus I$. Observe that $(\bar{x}, \bar{y}) \in B(\emptyset, \emptyset, \emptyset, N \setminus I)$ as
 311 $\sum_{j \in I} a_j \bar{x}_j \bar{y}_j + \sum_{j \in N \setminus I} a_j \bar{x}_j \geq \sum_{j \in N} a_j x'_j y'_j \geq d$. Since (26) is valid for $B(\emptyset, \emptyset, \emptyset, N \setminus I)$, (\bar{x}, \bar{y}) satisfies
 312 $\sum_{j \in N} \alpha_j \bar{x}_j + \sum_{j \in I} \beta_j \bar{y}_j = \sum_{j \in N} \alpha_j \bar{x}_j + \sum_{j \in I} \beta_j \bar{y}_j \geq \delta$. This is the desired contradiction.

313 Next, we show that (26) is facet-defining for PB . Since (26) is facet-defining for $PB(\emptyset, \emptyset, \emptyset, N \setminus I)$ and
 314 $\delta \neq 0$ as (26) is not a bound, there exist $n + |I|$ linearly independent points in $B(\emptyset, \emptyset, \emptyset, N \setminus I)$, call them
 315 (x^k, y^k) , that satisfy (26) at equality. Clearly, these points belong to B and satisfy (26) at equality. Now, for
 316 each $j \in N \setminus I$, we construct one new point in $B \setminus B(\emptyset, \emptyset, \emptyset, N \setminus I)$ that satisfies (26) at equality. Choose j
 317 arbitrarily in $N \setminus I$. Since (26) is not a scalar multiple of $x_j \leq 1$, there exists $k_j \in \{1, \dots, n + |I|\}$ such that
 318 $x_j^{k_j} = 0$. Now define $(\bar{x}^{k_j}, \bar{y}^{k_j})$ such that $\bar{x}_i^{k_j} = x_i^{k_j} \forall i \in N$, $\bar{y}_i^{k_j} = y_i^{k_j} \forall i \in N \setminus \{j\}$ and $\bar{y}_j^{k_j} = 0$. Clearly, the
 319 point $(\bar{x}^{k_j}, \bar{y}^{k_j})$ belongs to B and satisfies (26) at equality. Further, it is easily seen that the points (x^k, y^k)
 320 and $(\bar{x}^{k_j}, \bar{y}^{k_j})$ for $j \in N \setminus I$ are linearly independent and therefore show that (26) is facet-defining for PB .

321 To prove the reverse implication, we assume that (26) is a facet-defining inequality for PB that is not
 322 a scalar multiple of a bound. Validity is trivial since for $B(\emptyset, \emptyset, \emptyset, N \setminus I) \subseteq B$. Now, we show that (26) is
 323 facet-defining for $PB(\emptyset, \emptyset, \emptyset, N \setminus I)$. Since $\delta \neq 0$ as (26) is not a bound, the set of $2n$ affinely independent
 324 points (x^k, y^k) in B for $k = 1, \dots, 2n$ that satisfy (26) at equality are also linearly independent. Therefore,

$$325 \quad \begin{vmatrix} x_1^1 & \dots & x_n^1 & y_1^1 & \dots & y_n^1 \\ x_1^2 & \dots & x_n^2 & y_1^2 & \dots & y_n^2 \\ \dots & \dots & \dots & \dots & \dots & \dots \\ x_1^{2n} & \dots & x_n^{2n} & y_1^{2n} & \dots & y_n^{2n} \end{vmatrix} \neq 0.$$

326 Therefore, there must exist $n + |I|$ rows $i_1, \dots, i_{n+|I|}$ where $I = \{j_1, \dots, j_{|I|}\}$ such that

$$327 \quad \begin{vmatrix} x_1^{i_1} & \dots & x_n^{i_1} & y_{j_1}^{i_1} & \dots & y_{j_{|I|}}^{i_1} \\ x_1^{i_2} & \dots & x_n^{i_2} & y_{j_1}^{i_2} & \dots & y_{j_{|I|}}^{i_2} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ x_1^{i_{n+|I|}} & \dots & x_n^{i_{n+|I|}} & y_{j_1}^{i_{n+|I|}} & \dots & y_{j_{|I|}}^{i_{n+|I|}} \end{vmatrix} \neq 0.$$

328 Hence, we see that the $n + |I|$ points $(x_1^{i_k}, \dots, x_n^{i_k}; y_{j_1}^{i_k}, \dots, y_{j_{|I|}}^{i_k})$ for $k = 1, \dots, n + |I|$ are linearly independent.
 329 Now, define the points $(\tilde{x}^{i_k}, \tilde{y}^{i_k})$ for $k = 1, \dots, n + |I|$ such that $\tilde{x}^{i_k} = x^{i_k}$, $\tilde{y}_j^{i_k} = y_j^{i_k}$ for $j \in I$, and $\tilde{y}_j^{i_k} = 1$
 330 for $j \in N \setminus I$. The points $(\tilde{x}^{i_k}, \tilde{y}^{i_k})$ are feasible to $B(\emptyset, \emptyset, \emptyset, N \setminus I)$ and satisfy (26) at equality. Therefore, we
 331 conclude that (26) is facet-defining for $PB(\emptyset, \emptyset, \emptyset, N \setminus I)$. \square

332 Observe that Proposition 2.10 implies that all nontrivial facets of the 0–1 knapsack polytope can be found
 333 in B and that it is sufficient to study the facets of B to obtain the facets of the 0–1 knapsack polytope.
 334 Next, we use Proposition 2.10 to generalize the inequality $x_1 + x_2 \geq 1$ of Proposition 2.7 into an inequality
 335 that will be used as a seed for lifting procedures in Section 3.4.

336 **Proposition 2.11.** *Assume that $\sum_{j \in N} a_j - a_k - a_m < d$ for all $k, m \in N$ with $k \neq m$. The clique inequality*

$$337 \quad \sum_{j \in N} x_j \geq |N| - 1 \quad (27)$$

338 *is facet-defining for PB .*

339 *Proof.* Because of Proposition 2.10, it is sufficient to prove that (27) is facet-defining for $PB(\emptyset, \emptyset, \emptyset, N)$. To
 340 prove validity, assume for a contradiction that there exists $x' \in B(\emptyset, \emptyset, \emptyset, N)$ such that $\sum_{j \in N} a_j x'_j \geq d$ and
 341 $\sum_{j \in N} x'_j \leq |N| - 2$. Since $\sum_{j \in N} a_j x'_j \leq |N| - 2$, there exist $k, m \in N$ with $k \neq m$ such that $x'_k = 0$ and $x'_m = 0$.
 342 Therefore, $\sum_{j \in N} a_j - a_k - a_m \geq \sum_{j \in N} a_j x'_j \geq d$. This contradicts the assumption that $\sum_{j \in N} a_j - a_k - a_m < d$
 343 for all $k, m \in N$ with $k \neq m$. We next show that (27) is facet-defining for $PB(\emptyset, \emptyset, \emptyset, N)$. It can be easily
 344 verified using Assumption 2 that the points $p^k = (\mathbf{1} - e_k, \mathbf{1})$ for $k \in N$ belong to $B(\emptyset, \emptyset, \emptyset, N)$. Since
 345 these points are linearly independent and satisfy (27) at equality, we conclude that (27) is facet-defining for
 346 $PB(\emptyset, \emptyset, \emptyset, N)$. \square

347 **3 Lifted inequalities**

348 In this section, we derive three families of strong valid inequalities for PB via lifting. The first two families
 349 are obtained using sequence-independent lifting from (25) and are facet-defining for PB . In this case, lifting
 350 is simple since the lifting function is subadditive. The third inequality is obtained by lifting (27). Although
 351 the lifting function associated with this seed inequality is not subadditive, we obtain lifted inequalities using
 352 approximate lifting. We then identify conditions under which these inequalities are facet-defining for PB .

353 **3.1 Sequence-independent lifting for bilinear covering sets**

354 Sequence-independent lifting is a well-known technique to construct strong valid inequalities for mixed-integer
 355 linear programs; see Wolsey [33] and Gu et al. [13]. We next give a brief description of how this technique
 356 can be used to derive strong valid inequalities for PB . A more general treatment of lifting in nonlinear
 357 programming is given in Richard and Tawarmalani [24].

358 Given $\emptyset \neq S \subsetneq N$, we consider $B(S, \emptyset, S, \emptyset)$, which is the restriction of B obtained when all variables
 359 (x_j, y_j) for $j \in S$ are fixed to $(0, 0)$. Let $S = \{s, \dots, n\}$ for some $s \geq 2$ and define $S_i = \{i + 1, \dots, n\}$ for
 360 $i \in S$. Assume that the inequality

$$361 \quad \sum_{j=1}^{s-1} \alpha_j x_j + \sum_{j=1}^{s-1} \beta_j y_j \geq \delta \quad (28)$$

362 is facet-defining for $PB(S, \emptyset, S, \emptyset)$. In sequential lifting, we reintroduce the variables (x_j, y_j) for $j \in S$ one at
 363 the time in (28). Assuming that variables (x_j, y_j) have already been lifted in the order $j = s, \dots, i - 1$, we
 364 next review how to lift variables (x_i, y_i) in the inequality

$$365 \quad \sum_{j=1}^{i-1} \alpha_j x_j + \sum_{j=1}^{i-1} \beta_j y_j \geq \delta, \quad (29)$$

366 which is assumed to be facet-defining for $PB(S_{i-1}, \emptyset, S_{i-1}, \emptyset)$. To perform this lifting, we first compute the
 367 *lifting function*

$$\begin{aligned}
 368 \quad P^i(w) = \quad & \max \quad \delta - \left\{ \sum_{j=1}^{i-1} \alpha_j x_j + \sum_{j=1}^{i-1} \beta_j y_j \right\} \\
 369 \quad & \text{s.t.} \quad \sum_{j=1}^{i-1} a_j x_j y_j \geq d - w \\
 370 \quad & x_j \in \{0, 1\}, y_j \in [0, 1] \quad j = 1, \dots, i-1.
 \end{aligned}$$

371 Once the lifting function $P^i(w)$ is computed, the lifting coefficients (α_i, β_i) are obtained from $P^i(w)$ as
 372 follows.

373 **Proposition 3.1** (Richard and Tawarmalani [24]). *Let (29) be a valid inequality for $B(S_{i-1}, \emptyset, S_{i-1}, \emptyset)$.
 374 Assume that there exist $(\alpha_i, \beta_i) \in \mathbb{R}^2$ such that*

$$375 \quad \alpha_i x_i + \beta_i y_i \geq P^i(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{0, 0\}. \quad (30)$$

376 *Then, the inequality*

$$377 \quad \sum_{j=1}^i \alpha_j x_j + \sum_{j=1}^i \beta_j y_j \geq \delta \quad (31)$$

378 *is valid for $B(S_i, \emptyset, S_i, \emptyset)$.*

379 The result of Proposition 3.1 can be applied recursively to construct a valid inequality for PB from
 380 (28). Note that, at each step, the lifting function $P^i(w)$ must be recomputed to account for the changes
 381 in the lifted inequality. Further, if $B(S, \emptyset, S, \emptyset)$ is full-dimensional, the seed inequality (28) is facet-defining
 382 for $B(S, \emptyset, S, \emptyset)$, and for each $i \in S$, the lifting coefficients (α_i, β_i) of the variables (x_i, y_i) are chosen so
 383 that (30) is satisfied at equality by two points (x_i^1, y_i^1) and (x_i^2, y_i^2) such that $(0, 0)$, (x_i^1, y_i^1) and (x_i^2, y_i^2)
 384 are affinely independent (a feature we refer to as *maximal lifting*), then the final lifted inequality will be
 385 facet-defining for PB . In this scheme, (re)computing the lifting functions $P^i(w)$ for each $i \in S$ is often the
 386 most computationally demanding task. However, this computational work is unnecessary when the lifting
 387 function $P^s(w)$ is subadditive. This observation, first made by Wolsey [33], leads to the following result.

388 **Proposition 3.2** (Richard and Tawarmalani [24]). *Assume that (28) is valid for $B(S, \emptyset, S, \emptyset)$. Assume also
 389 that (i) $P^s(w)$ is subadditive over \mathbb{R}_+ , i.e., $P^s(w_1) + P^s(w_2) \geq P^s(w_1 + w_2) \forall w_1, w_2 \in \mathbb{R}_+$, and (ii) there
 390 exist $(\alpha_i, \beta_i) \in \mathbb{R}^2$ for all $i \in S$ such that*

$$391 \quad \alpha_i x_i + \beta_i y_i \geq P^s(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{0, 0\}. \quad (32)$$

392 *Then, the inequality*

$$393 \quad \sum_{j=1}^n \alpha_j x_j + \sum_{j=1}^n \beta_j y_j \geq \delta \quad (33)$$

394 *is valid for PB . Further, if (i) Inequality (28) is facet-defining for $B(S, \emptyset, S, \emptyset)$, (ii) $B(S, \emptyset, S, \emptyset)$ is full-
 395 dimensional and (iii) coefficients (α_i, β_i) are chosen in a way that two linearly independent points satisfy
 396 (32) at equality, then (33) is facet-defining for PB .*

397 The fundamental difference between Proposition 3.1 and Proposition 3.2 lies in equations (30) and (32).
 398 In the latter, the lifting coefficients of all variables (x_i, y_i) are obtained from the same lifting function $P^s(w)$
 399 while in the former, they are obtained from $P^i(w)$ for $i \in S$. Although this difference might seem minor, it
 400 has important practical implications. In particular, the subadditivity of lifting functions typically permits
 401 the derivation of closed-form expressions for lifting coefficients that would otherwise be difficult to obtain.
 402 Observe also that in Proposition 3.2, the subadditivity of $P^s(w)$ is required only over \mathbb{R}_+ since all coefficients
 403 a_i in PB are assumed to be nonnegative.

404 Proposition 3.1 describes how to perform lifting when then variables (x_j, y_j) for $j \in S$ are fixed at $(0, 0)$.
 405 When variables (x_j, y_j) are fixed at $(1, 1)$, similar results can be obtained. In this case, condition (30) must
 406 be changed to

$$407 \quad \alpha_i(x_i - 1) + \beta_i(y_i - 1) \geq P^i(a_i x_i y_i - a_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{1, 1\}. \quad (34)$$

408 Similarly, Proposition 3.2 can be adapted to allow sequence-independent lifting for variables (x_j, y_j) fixed at
 409 $(1, 1)$ by replacing $P^i(w)$ with $P^s(w)$ in (34) and by requiring that the lifting function $P^s(w)$ is subadditive
 410 over \mathbb{R}_- . Subadditive lifting can also be used to generate facets of PB if $B(\emptyset, S, \emptyset, S)$ is full-dimensional, the
 411 seed inequality (28) is facet-defining for $B(\emptyset, S, \emptyset, S)$, and for each $i \in S$, the lifting coefficients (α_i, β_i) of the
 412 variables (x_i, y_i) are chosen so that (34) is satisfied at equality by two points (x_i^1, y_i^1) and (x_i^2, y_i^2) such that
 413 $(1, 1)$, (x_i^1, y_i^1) and (x_i^2, y_i^2) are affinely independent.

414 We next show in the following proposition that all interesting lifted inequalities that can be obtained by
 415 fixing variables (x_i, y_i) at $(0, 1)$ or $(1, 0)$ can also be obtained by fixing variables (x_i, y_i) at $(0, 0)$.

416 **Proposition 3.3.** *Assume that (29) defines a nonempty face of $PB(S_{i-1}, \emptyset, S_{i-1}, \emptyset) = PB(S_{i-1}, \emptyset, \emptyset, S_{i-1}) =$
 417 $PB(\emptyset, S_{i-1}, S_{i-1}, \emptyset)$. Then any inequality obtained from maximally lifting (29) in $PB(S_{i-1}, \emptyset, \emptyset, S_{i-1})$ or
 418 $PB(\emptyset, S_{i-1}, S_{i-1}, \emptyset)$ could have been obtained by maximally lifting (29) in $PB(S_{i-1}, \emptyset, S_{i-1}, \emptyset)$.*

419 *Proof.* First, we consider the case when (x_i, y_i) is fixed at $(1, 0)$. In this situation, valid lifting coefficients
 420 must satisfy

$$421 \quad \alpha_i(x_i - 1) + \beta_i y_i \geq P^i(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1]. \quad (35)$$

422 We next show that maximal lifting coefficients (α_i, β_i) in (35) must also satisfy

$$423 \quad \alpha_i x_i + \beta_i y_i \geq P^i(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \quad (36)$$

424 and be maximal for (36). This is sufficient to prove the result since restricting $(x_i, y_i) = (0, 0)$ instead of
 425 $(1, 0)$ does not change the restricted set and, therefore, the seed inequality is still a face of same dimension.
 426 Let $(0, y_i^*)$ satisfy (35) at equality. Such a point exists since lifting is assumed to be maximal. Then,

$$427 \quad 0 \geq \alpha_i = \beta_i y_i^* \geq P^i(a_i y_i^*) \geq 0,$$

428 where the first inequality follows from (35) by setting $(x_i, y_i) = (0, 0)$, the equality holds since $(0, y_i^*)$ satisfies
 429 (35) at equality, the second inequality is satisfied from (35) with $(x_i, y_i) = (1, y_i^*)$ and the last inequality
 430 is verified since $a_i y_i^* \geq 0$. Therefore, equality holds throughout and, in particular, $\alpha_i = 0$. It follows that
 431 $\alpha_i(x_i - 1) + \beta_i y_i = \alpha_i x_i + \beta_i y_i$ and, consequently, (α_i, β_i) is valid and maximal to (36).

432 Now, we fix (x_i, y_i) at $(0, 1)$. Then, we show that any (α_i, β_i) that is valid and maximal to

$$433 \quad \alpha_i x_i + \beta_i(y_i - 1) \geq P^i(a_i x_i y_i) \quad (37)$$

434 is also valid and maximal to (36). Let $y_i^* = \min\{y_i \in [0, 1] \mid \alpha_i + \beta_i(y_i - 1) = P^i(a_i y_i)\}$, *i.e.*, $(1, y_i^*)$ satisfies
 435 (37) at equality. It follows that

$$436 \quad 0 \leq \beta_i(y_i^* - 1) = P^i(a_i y_i^*) - \alpha_i \leq P^i(a_i y_i^*) - P^i(a_i) \leq 0,$$

437 where the first inequality follows from (37) by substituting $(x_i, y_i) = (0, y_i^*)$, the equality is satisfied since
 438 $(1, y_i^*)$ satisfies (37) at equality, the second inequality is verified by substituting $(1, 1)$ in (37), and the last
 439 inequality holds since $P^i(\cdot)$ is non-decreasing and $a_i y_i^* \leq a_i$. Therefore, the equality holds throughout and,
 440 in particular, $\beta_i(y_i^* - 1) = 0$. It follows that either $\beta_i = 0$ or $y_i^* = 1$. We show that $\beta_i = 0$ in the latter case
 441 as well. If $y_i^* = 1$, because lifting is assumed to be maximal and because of the definition of y_i^* , there is a
 442 $y_i' \in [0, 1)$ such that $(0, y_i')$ satisfies (37) at equality. Therefore, $\beta_i(y_i' - 1) = 0$ and so $\beta_i = 0$. It follows that
 443 $\alpha_i x_i + \beta_i(y_i - 1) = \alpha_i x_i + \beta_i y_i$ and, consequently, (α_i, β_i) is valid and maximal for (36). \square

3.2 Subadditivity of lifting functions

In this section, we provide a general result that helps in proving subadditivity of functions. For specific functions, it has been observed (see Proposition 3.10 in [23], Theorem 7 in [11], Proposition 4.2 in [8], and Lemma 21 in [25]) that subadditivity of a function over \mathbb{R}^n can often be established by checking it at a small subset of points. The corresponding proofs are often detailed and are the key step in proving subadditivity. In Theorem 3.4, we identify a fairly large class of functions for which a similar result holds. We use this result to prove subadditivity of functions that arise during lifting of inequalities for mixed-integer 0–1 bilinear covering set. The scope of applications of Theorem 3.4 is, however, much larger and we provide this general result with the hope that it may be useful in other applications.

Theorem 3.4. *For $i \in N = \{1, \dots, n\}$, let $b_i \in \mathbb{R}^m$, $f_i \in \mathbb{R}$, and $h_i(x) : \mathbb{R}^m \mapsto \mathbb{R}$ be subadditive functions. Let $h(x) : \mathbb{R}^m \mapsto \mathbb{R}$ be a function with $h(\mathbf{0}) = 0$ that majorizes $h_i(x)$ for all i . Let $B_i = \{x \in \mathbb{R}^m \mid h_i(x - b_i) = h(x - b_i)\}$. Define $f(x) : \mathbb{R}^m \mapsto \mathbb{R}$ as $f(x) = \min_{i=1}^n \{f_i + h_i(x - b_i)\}$. Then, $f(x) + h(y - x) \geq f(y)$ for each $(x, y) \in \mathbb{R}^m \times \mathbb{R}^m$. For $x \in \mathbb{R}^m$, let $i(x) \in N$ be such that $f(x) = f_{i(x)} + h_{i(x)}(x - b_{i(x)})$. If $x \in B_{i(x)}$ then $f(b_{i(x)}) - f(b_{i(x)} + y) \leq f(x) - f(x + y)$. Further, if $y \in B_{i(y)}$ as well, then*

$$f(b_{i(x)}) + f(b_{i(y)}) - f(b_{i(x)} + b_{i(y)}) \leq f(x) + f(y) - f(x + y).$$

Proof. Let $(x, y) \in \mathbb{R}^m \times \mathbb{R}^m$. Then,

$$f(x) + h(y - x) \geq f(x) + h_{i(x)}(y - x) = f_{i(x)} + h_{i(x)}(x - b_{i(x)}) + h_{i(x)}(y - x) \geq f_{i(x)} + h_{i(x)}(y - b_{i(x)}) \geq f(y), \quad (38)$$

where the first inequality follows since $h(\cdot) \geq h_{i(x)}(\cdot)$, the second inequality holds since $h_{i(x)}$ is a subadditive function, and the third since the minimum defining $f(y)$ includes a term that equals $f_{i(x)} + h_{i(x)}(y - b_{i(x)})$. Note that $0 \leq h_{i(x)}(\mathbf{0}) \leq h(\mathbf{0}) = 0$ where the first inequality follows from subadditivity of $h_{i(x)}$. Therefore, $h_{i(x)}(\mathbf{0}) = 0$. Further, if $x \in B_{i(x)}$:

$$f(x) = f_{i(x)} + h_{i(x)}(x - b_{i(x)}) \geq f(b_{i(x)}) + h_{i(x)}(x - b_{i(x)}) = f(b_{i(x)}) + h(x - b_{i(x)}) \geq f(x),$$

where the first inequality holds since $h_{i(x)}(\mathbf{0}) = 0$ implies $f(b_{i(x)}) \leq f_{i(x)}$ as $f_{i(x)}$ is one of the terms in minimum defining $f(b_{i(x)})$, the second equality since $x \in B_{i(x)}$ and the last inequality by (38). Therefore, equality holds throughout and, in particular, $f_{i(x)} = f(b_{i(x)})$. Now, consider $(x, y) \in \mathbb{R}^m \times \mathbb{R}^m$ with $x \in B_{i(x)}$. Then,

$$f(b_{i(x)}) - f(b_{i(x)} + y) = f(x) - h(x - b_{i(x)}) - f(b_{i(x)} + y) \leq f(x) - f(x + y), \quad (39)$$

where the equality follows from the definition of $i(x)$, $x \in B_{i(x)}$, and $f_{i(x)} = f(b_{i(x)})$, and the inequality follows from (38) since $f(b_{i(x)} + y) + h(x - b_{i(x)}) \geq f(x + y)$. Further, if $y \in B_{i(y)}$,

$$f(b_{i(x)}) + f(b_{i(y)}) - f(b_{i(x)} + b_{i(y)}) \leq f(b_{i(x)}) + f(y) - f(b_{i(x)} + y) \leq f(x) + f(y) - f(x + y),$$

where each of the inequalities follows from (39). \square

Any function, say $g(x)$, that is subadditive and satisfies $g(\mathbf{0}) = 0$ can be expressed as $f(x)$, defined in Theorem 3.4, by setting $n = 1$, $b_1 = \mathbf{0}$, $f_1 = 0$, and $h(x) = h_1(x) = g(x)$. Observe that lifting functions derived from seed inequalities that are tight on the restricted set always satisfy the condition $g(\mathbf{0}) = 0$. In other words, Theorem 3.4 can be interpreted as a recursive tool for proving subadditivity of such lifting functions, $f(x)$, that exploits the subadditivity of the constituent simpler functions, $h_i(\cdot)$. For example, Theorem 3.4 shows that it suffices to check the subadditivity of $f(x)$ at a small subset of points if $f(x)$ can be expressed as a minimum of finitely many translates of a subadditive function. Since positively-homogenous convex functions belong to the class of subadditive functions, see Theorem 4.7 in [26], they can be used as building blocks in the application of Theorem 3.4. In Corollaries 3.5 and 3.6, we apply Theorem 3.4 to prove subadditivity of two functions that will be used in Sections 3.3 and 3.4 to derive inequalities for the 0–1 mixed-integer bilinear covering set. In both cases, functions h_i are univariate positively-homogenous convex functions.

487 **Corollary 3.5.** Let ν and D_i for $i = 0, 1, \dots, r$ be nonnegative integers that satisfy $\nu > 0$, $D_0 = 0$, and
 488 $D_i \geq D_{i-1} + \nu$ for $i = 1, \dots, r$. Then the function

$$489 \quad g(w) := \begin{cases} 0 & \text{if } w < D_0 \\ w - i\nu & \text{if } D_i \leq w < D_{i+1} - \nu, \quad i = 0, \dots, r-1, \\ D_i - i\nu & \text{if } D_i - \nu \leq w < D_i, \quad i = 1, \dots, r-1, \\ D_r - r\nu & \text{if } D_r - \nu \leq w \end{cases}$$

490 is subadditive over \mathbb{R} if and only if $D_i + D_j \geq D_{i+j}$ for $0 \leq i \leq j \leq r$ with $i + j \leq r$.

491 *Proof.* First note that $g(w) = \min_{i=0}^r \{D_i - i\nu + h_i(w - D_i)\}$, where $h_i(w) = \max\{0, w\}$ for $i = 0, \dots, r-1$ and
 492 $h_r(w) = 0$. We observe that $i(x) = 0$ for $x < D_1 - \nu$, $i(x) = i$ for $x \in [D_i - \nu, D_{i+1} - \nu)$ and $i = 1, \dots, r-1$,
 493 and $i(x) = r$ for $x \geq D_r - \nu$. Let $h(w) = \max\{0, w\}$. We also see that $B_0 = B_1 = \dots = B_{r-1} = \mathbb{R}$ and
 494 $B_r = \mathbb{R}_-$.

495 Assume that $D_i + D_j \geq D_{i+j}$ for $0 \leq i \leq j \leq r$ with $i + j \leq r$. Consider $x, y \in \mathbb{R}$ with $x \leq y$. We
 496 argue next that $g(x) + g(y) \geq g(x + y)$. Define $i = i(x)$ and $j = i(y)$. Clearly, $i \leq j$. We consider two cases.
 497 Assume first that $j = r$. Then, $g(y) = D_r - r\nu \geq g(x + y)$ and therefore $g(x) + g(y) \geq g(x + y)$ as $g(x) \geq 0$.
 498 Assume next that $j \leq r - 1$. Since $x \leq y < D_r - \nu$, it follows that $x \in B_i$ and $y \in B_j$ and, therefore, from
 499 Theorem 3.4 that

$$500 \quad g(D_i) + g(D_j) - g(D_i + D_j) \leq g(x) + g(y) - g(x + y)$$

501 We next argue that the left-hand-side of the above expression is nonnegative, which proves the result. Let
 502 $t = \min\{j, r - i\}$. Then,

$$503 \quad g(D_i + D_j) = g(D_i + D_t) = g(D_{i+t} + D_i + D_t - D_{i+t}) \leq g(D_{i+t}) + D_i + D_t - D_{i+t} = g(D_i) + g(D_t) \leq g(D_i) + g(D_j),$$

504 where the first equality holds since $t = r - i$ implies $D_i + D_j \geq D_i + D_t \geq D_r$, the first inequality follows
 505 from (38) and $D_{i+t} \leq D_i + D_t$, the second equality since $g(D_i) = D_i - i\nu$, and the last inequality from (38)
 506 since $D_t \leq D_j$.

507 We now prove the reverse implication. For $w > 0$,

$$508 \quad g(D_k - w) \geq g(D_k - \nu) - \max\{0, w - \nu\} = g(D_k) - \max\{0, w - \nu\} > g(D_k) - w, \quad (40)$$

509 where the first inequality follows from (38) and the last inequality since $\nu > 0$ and $w > 0$. If $i + j \leq r$ and
 510 $D_i + D_j < D_{i+j}$ then

$$511 \quad g(D_i) + g(D_j) - g(D_i + D_j) < g(D_i) + g(D_j) - g(D_{i+j}) - D_i - D_j + D_{i+j} = 0,$$

512 yields a contradiction to subadditivity of g , where the strict inequality follows from (40) where $k = i + j$ and
 513 $w = D_{i+j} - D_i - D_j$ since $D_i + D_j < D_{i+j}$ and the equality holds since $g(D_k) = D_k - k\nu$ for $k \in \{i, j, i + j\}$. \square

514 Corollary 3.5 equivalently shows the superadditivity of $w - g(w)$, generalizing prior similar results in the
 515 literature. In particular, see Lemmas 6 and 7 in [3] and Definition 4 in [19].

516 **Corollary 3.6.** Let λ and C_i for $i = 0, 1, \dots, s$ be nonnegative integers that satisfy $\lambda > 0$, $C_0 = 0$ and
 517 $C_{i-1} + \lambda \leq C_i$ for $i = 1, \dots, s$. Then the function

$$518 \quad g(w) = \begin{cases} 0 & \text{if } w < C_0 \\ i + \frac{w - C_i}{\lambda} & \text{if } C_i \leq w < C_i + \lambda, \quad i = 0, \dots, s, \\ i & \text{if } C_{i-1} + \lambda \leq w < C_i, \quad i = 1, \dots, s, \\ s + 1 & \text{if } C_s + \lambda \leq w. \end{cases}$$

519 is subadditive over \mathbb{R} if and only if $C_i + C_j \leq C_{i+j}$ for $0 \leq i \leq j \leq s$ with $i + j \leq s$.

520 *Proof.* Let $C_{s+1} = \max\{C_i + C_j \mid i + j = s + 1\}$. Note that $g(w) = \min_{i=0}^{s+1} \{i + h_i(w - C_i)\}$ where
 521 $h_i(w) = \max\{0, \frac{w}{\lambda}\}$ for $i = 0, \dots, s$ and $h_{s+1}(w) = 0$. Let $h(w) = \max\{0, \frac{w}{\lambda}\}$.

522 Assume that $C_i + C_j \leq C_{i+j}$ for $0 \leq i \leq j \leq s$ with $i + j \leq s$. Consider $x, y \in \mathbb{R}$ with $x \leq y$. We
 523 argue next that $g(x) + g(y) \geq g(x + y)$. Define $i = i(x)$ and $j = i(y)$. Assume first that $j = s + 1$. Then

524 $g(y) = s + 1 \geq g(x + y)$ and therefore, $g(x) + g(y) \geq g(x + y)$ as $g(x) \geq 0$. Next assume that $j \leq s$. Since
 525 $x \leq y < C_s + \lambda$, it follows that $x \in B_i$ and $y \in B_j$ and therefore, from Theorem 3.4 that

$$526 \quad g(C_i) + g(C_j) - g(C_i + C_j) \leq g(x) + g(y) - g(x + y).$$

527 We next argue that the left-hand-side of the above expression is nonnegative, which proves the result. Let
 528 $t = \min\{j, s + 1 - i\}$. Then,

$$529 \quad g(C_i + C_j) \leq g(C_{i+t}) = i + t \leq i + j = g(C_i) + g(C_j),$$

530 where the first inequality follows from (38) and $C_{i+t} \geq C_i + C_j$ when $t = j$ and from $h_{s+1}(w) = 0$ when
 531 $t < j$, and the last equality holds because $g(C_k) = k$ for $k \in \{i, j\}$.

532 We now prove the reverse implication. For $w > 0$ and $i \leq s$,

$$533 \quad g(C_k + w) \geq g(C_k) + \max\left\{0, \frac{w}{\lambda}\right\} > g(C_k), \quad (41)$$

534 where the first inequality follows from (38) and the strict inequality from $w > 0$ and $\lambda > 0$. If $i + j \leq s$ and
 535 $C_i + C_j > C_{i+j}$ then

$$536 \quad g(C_i) + g(C_j) - g(C_i + C_j) < g(C_i) + g(C_j) - g(C_{i+j}) = 0,$$

537 yields a contradiction to subadditivity of $g(\cdot)$ where the strict inequality follows from (41) where $k = i + j$
 538 and $w = C_i + C_j - C_{i+j}$ and the equality holds since $g(C_k) = k$ for $k \in \{i, j, i + j\}$. \square

539 3.3 Lifted inequalities by sequence-independent lifting

540 In this section, we derive strong inequalities for PB through lifting using (25) as seed inequality. To describe
 541 the general form of these inequalities, we use the notion of a cover, which is adapted from the definition of a
 542 cover for the 0–1 knapsack polytope; see Balas [4], Hammer et al. [14], and Wolsey [31].

543 **Definition 3.7.** Let $C \subseteq N$. We say that C is a cover for B if $\sum_{j \in C} a_j > d$. Further, we define the excess
 544 of the cover as $\mu = \sum_{j \in C} a_j - d > 0$.

545 We create lifted inequalities by first partitioning the set of variables N into $(C', \{l\}, M, T)$ in such a way
 546 that:

547 (A1) $C := C' \cup \{l\}$ is a cover for B with excess μ ,

548 (A2) $a_l \geq a_j, \forall j \in C'$,

549 (A3) $a_l > \mu$,

550 (A4) $\sum_{j \in C \cup T} a_j > d + a_l$, i.e., $\sum_{j \in T} a_j > a_l - \mu$.

551 Note that (A1) and (A3) might be reminiscent of conditions that make a cover minimal for the 0–1 knapsack
 552 polytope. We note however that minimal covers require $a_j > \mu$ for all $j \in C$ and not simply $a_l > \mu$. Note
 553 also that (A4) implies that $T \neq \emptyset$. To obtain lifted inequalities from $(C', \{l\}, M, T)$, we first fix the variables
 554 (x_j, y_j) for $j \in M$ to $(0, 0)$ and the variables (x_j, y_j) for $j \in C'$ to $(1, 1)$. The resulting (full-dimensional) set
 555 $B(M, C', M, C')$ is then defined by the inequality

$$556 \quad a_l x_l y_l + \sum_{j \in T} a_j x_j y_j \geq d - \sum_{j \in C'} a_j = a_l - \mu.$$

557 Since $a_l > \mu$ and $\sum_{j \in T} a_j > a_l - \mu$ from Conditions (A3) and (A4), we conclude from Proposition 2.9(i) that

$$558 \quad (a_l - \mu)x_l + \sum_{j \in T} a_j y_j \geq a_l - \mu \quad (42)$$

559 is facet-defining for $PB(M, C', M, C')$. We will create two different families of lifted inequalities for PB by
 560 reintroducing the variables (x_j, y_j) for $j \in M \cup C'$ in different orders. To derive both families, we use the
 561 lifting function

$$\begin{aligned}
 562 \quad P(w) := \quad & \max \quad (a_l - \mu) - \left\{ (a_l - \mu)x_l + \sum_{j \in T} a_j y_j \right\} \\
 563 \quad & \text{s.t.} \quad a_l x_l y_l + \sum_{j \in T} a_j x_j y_j \geq a_l - \mu - w \\
 564 \quad & x_j \in \{0, 1\}, y_j \in [0, 1] \quad \forall j \in \{l\} \cup T.
 \end{aligned} \tag{43}$$

565 We next derive a closed-form expression for $P(w)$.

Proposition 3.8.

$$566 \quad P(w) = \begin{cases} -\infty & \text{if } w < -\sum_{j \in T} a_j - \mu, \\ w + \mu & \text{if } -\sum_{j \in T} a_j - \mu \leq w < -\mu, \\ 0 & \text{if } -\mu \leq w < 0, \\ w & \text{if } 0 \leq w < a_l - \mu, \\ a_l - \mu & \text{if } a_l - \mu \leq w. \end{cases}$$

567 Further, $P(w)$ is subadditive over \mathbb{R}_- and \mathbb{R}_+ respectively.

568 *Proof.* We first derive a closed-form expression for $P(w)$. Observe that, if (43) is feasible, there exists an
 569 optimal solution (x^*, y^*) to (43) for which $x_j^* = 1$ for $j \in T$ and $y_l^* = 1$ since the coefficients of x_j for
 570 $j \in T$ and y_l in the objective are equal to 0. Defining $\bar{a} = \sum_{j \in T} a_j$ and $\bar{y} = \frac{\sum_{j \in T} a_j y_j}{\bar{a}}$, we can simplify the
 571 formulation of $P(w)$ in (43) as:

$$\begin{aligned}
 572 \quad P(w) = \quad & \max \quad (a_l - \mu) - \{(a_l - \mu)x_l + \bar{a}\bar{y}\} \\
 573 \quad & \text{s.t.} \quad a_l x_l + \bar{a}\bar{y} \geq a_l - \mu - w \\
 574 \quad & x_l \in \{0, 1\}, \bar{y} \in [0, 1].
 \end{aligned} \tag{44}$$

575 When $w < -\bar{a} - \mu$, (44) is infeasible and so $P(w) = -\infty$. When $w \geq a_l - \mu$, the optimal solution is $x_l^* = 0$
 576 and $\bar{y}^* = 0$ with $P(w) = a_l - \mu$. For $-\bar{a} - \mu \leq w < a_l - \mu$, there are two cases. When $-\bar{a} - \mu \leq w < a_l - \bar{a} - \mu$,
 577 then every feasible solution (x_l^*, \bar{y}^*) has $x_l^* = 1$. Further, the optimal solution has $\bar{y}^* = \max\{\frac{-\mu - w}{\bar{a}}, 0\}$. It
 578 follows that $P(w) = \min\{w + \mu, 0\}$. When $a_l - \bar{a} - \mu \leq w \leq a_l - \mu$, an optimal solution must be found among
 579 the solutions $(1, \frac{(-\mu - w)^+}{\bar{a}})$ and $(0, \frac{a_l - \mu - w}{\bar{a}})$. It follows that $P(w) = \max\{(w + \mu)^-, w\}$ from which we obtain
 580 the desired expression for $P(w)$ after considering both the cases where $a_l - \bar{a} < 0$ and $a_l - \bar{a} \geq 0$.

581 Subadditivity of $P(w)$ over \mathbb{R}_- and \mathbb{R}_+ follows from Karamata/Hardy-Littlewood-Polya inequality [15],
 582 concavity of $P(w)$ over these domains and $P(0) = 0$. \square

584 We note that, although $P(w)$ is subadditive over \mathbb{R}_+ and over \mathbb{R}_- , $P(w)$ is not subadditive over \mathbb{R} as
 585 $P(2a_l - \mu) + P(-a_l) = (a_l - \mu) + (-a_l + \mu) = 0 < a_l - \mu = P(a_l - \mu)$.

3.3.1 Lifted bilinear cover inequalities

587 To obtain lifted bilinear cover inequalities, we will lift first the variables (x_i, y_i) for $i \in C'$ from $(1, 1)$ and
 588 then lift the variables (x_i, y_i) for $i \in M$ from $(0, 0)$. Since $P(w)$ is subadditive over \mathbb{R}_- , we can apply
 589 sequence-independent lifting for the variables (x_i, y_i) for $i \in C'$.

590 **Proposition 3.9.** Under Conditions (A1), (A2), (A3) and (A4),

$$591 \quad \sum_{j \in C} (a_j - \mu)^+ x_j + \sum_{j \in T} a_j y_j \geq \sum_{j \in C} (a_j - \mu)^+ \tag{45}$$

592 is facet-defining for $PB(M, \emptyset, M, \emptyset)$.

593 *Proof.* The seed inequality (42) is facet-defining for the full-dimensional polytope $PB(M, C', M, C')$. Since
 594 $P(w)$ is subadditive over \mathbb{R}_- , we obtain from the remark following Proposition 3.2 that the lifting coefficients
 595 (α_i, β_i) for (x_i, y_i) for $i \in C'$ are valid if they satisfy

$$596 \quad \alpha_i(x_i - 1) + \beta_i(y_i - 1) \geq P(a_i x_i y_i - a_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{1, 1\}. \quad (46)$$

597 This condition can be also written as

$$598 \quad \beta_i \leq \inf_{0 \leq \phi < 1} \frac{-P(a_i \phi - a_i)}{1 - \phi}, \quad (47)$$

$$599 \quad \alpha_i + \sup_{0 \leq \phi \leq 1} \beta_i(1 - \phi) \leq -P(-a_i). \quad (48)$$

600 From Conditions (A2) and (A4), we know that $a_i \leq a_l < \sum_{j \in T} a_j + \mu$, $\forall i \in C'$. Therefore, in (47)
 601 $a_i \phi - a_i \in (-\sum_{j \in T} a_j - \mu, 0)$ for all $\phi \in [0, 1)$. Since $P(w) \leq 0$ for $w \leq 0$, we conclude that

$$602 \quad \frac{-P(a_i \phi - a_i)}{1 - \phi} \geq 0, \quad \forall 0 \leq \phi < 1,$$

603 and therefore choosing $\beta_i = 0$ for $i \in C'$ satisfies (47). Further, as $\beta_i = 0$, it is simple to verify that choosing
 604 $\alpha_i = -P(-a_i) = (a_i - \mu)^+$ satisfies (48). Finally, observe that (46) is satisfied at equality by the two
 605 points $(0, 0)$ and $(1, \frac{(a_i - \mu)^+}{a_i})$ that are affinely independent of $(1, 1)$. Therefore, we conclude that (45) is
 606 facet-defining for $PB(M, \emptyset, M, \emptyset)$. \square

607 Now, we lift the variables (x_j, y_j) for $j \in M$ in (45). The corresponding lifting function is

$$608 \quad P^C(w) := \max \sum_{j \in C} (a_j - \mu)^+ - \left\{ \sum_{j \in C} (a_j - \mu)^+ x_j + \sum_{j \in T} a_j y_j \right\}$$

$$609 \quad \text{s.t.} \quad \sum_{j \in C \cup T} a_j x_j y_j \geq \sum_{j \in C} a_j - \mu - w \quad (49)$$

$$610 \quad x_j \in \{0, 1\}, y_j \in [0, 1] \quad \forall j \in C \cup T.$$

611 We next derive a closed-form expression for $P^C(w)$. To this end, we assume without loss of generality that
 612 $C = \{1, \dots, p\}$ and that $a_1 \geq a_2 \geq \dots \geq a_p$. We also let $q \in C$ be such that $a_q > \mu \geq a_{q+1}$. We define
 613 $A_0 = 0$ and $A_i = \sum_{j=1}^i a_j$ for all $i \in \{1, \dots, q\}$.

614 **Proposition 3.10.** *For $w \geq 0$,*

$$615 \quad P^C(w) = \begin{cases} w - i\mu & \text{if } A_i \leq w < A_{i+1} - \mu, \quad i = 0, \dots, q-1, \\ A_i - i\mu & \text{if } A_i - \mu \leq w < A_i, \quad i = 1, \dots, q-1, \\ A_q - q\mu & \text{if } A_q - \mu \leq w. \end{cases}$$

616 *Proof.* First, observe that there exists an optimal solution (x^*, y^*) of (49) in which $x_j^* = 1$ for $j \in T$ and
 617 $y_j^* = 1$ for $j \in C$ since the corresponding objective coefficients are zero. Since $a_q > \mu \geq a_{q+1}$, we have
 618 $(a_j - \mu)^+ = 0$ for $j = q+1, \dots, p$, which similarly implies that we can assume $x_j^* = 1$ for $j = q+1, \dots, p$.
 619 Defining $\bar{a} = \sum_{j \in T} a_j$ and $\bar{y} = \frac{\sum_{j \in T} a_j y_j}{\bar{a}}$, we simplify the expression of $P^C(w)$ as

$$620 \quad P^C(w) = \max \sum_{j=1}^q (a_j - \mu) - \left\{ \sum_{j=1}^q (a_j - \mu) x_j + \bar{a} \bar{y} \right\}$$

$$621 \quad \text{s.t.} \quad \sum_{j=1}^q a_j x_j + \bar{a} \bar{y} \geq \sum_{j=1}^q a_j - \mu - w \quad (50)$$

$$622 \quad x_j \in \{0, 1\}, \quad \forall j = 1, \dots, q, \quad \bar{y} \in [0, 1].$$

623 Next, we solve (50). When $w \geq A_q - \mu$, it is clear that $x_j^* = 0$ for $j = 1, \dots, q$ and $\bar{y}^* = 0$ is an optimal
 624 solution for (50), showing that $P^C(w) = A_q - q\mu$. It is therefore sufficient to consider $w \in [0, A_q - \mu)$. We
 625 consider two cases:

626 1. Assume that $A_i - \mu \leq w < A_{i+1} - \mu$ for $i \in \{1, \dots, q-1\}$. Let $\theta = (A_{i+1} - \mu) - w$. Clearly, $0 < \theta \leq a_{i+1}$.
627 Define first the solution (x^*, \bar{y}^*) where $x_j^* = 0$ for $j = 1, \dots, i+1$, $x_j^* = 1$ for $j = i+2, \dots, q$, and $\bar{y}^* = \frac{\theta}{a}$.
628 When $\theta \leq \bar{a}$, (x^*, \bar{y}^*) is a feasible solution to (50) with objective value $z^* = A_{i+1} - (i+1)\mu - \theta = w - i\mu$.
629 Next consider the solution (x', \bar{y}') where $x_j' = 0$ for $j = 1, \dots, i$, $x_j' = 1$ for $j = i+1, \dots, q$, and $\bar{y}' = 0$.
630 Solution (x', \bar{y}') is feasible to (50) and has objective value $z' = A_i - i\mu$. It is clear that $z^* \geq z'$ when
631 $\theta \leq a_{i+1} - \mu$ and that $z' \geq z^*$ when $a_{i+1} - \mu \leq \theta \leq a_{i+1}$. Further, solution (x^*, \bar{y}^*) is feasible when
632 $\theta \leq a_{i+1} - \mu$ as $a_{i+1} - \mu \leq a_1 - \mu \leq \bar{a}$ because of Condition (A4). Therefore, we conclude that
633 $P^C(w) \geq w - i\mu$ if $A_i \leq w \leq A_{i+1} - \mu$ and $P^C(w) \geq A_i - i\mu$ if $A_i - \mu \leq w < A_i$.

634 We now prove that the proposed solutions are optimal. Pick any feasible solution (x°, \bar{y}°) to (50).
635 Define $N_1 = \{j \in \{1, \dots, q\} \mid x_j^\circ = 1\}$. Consider first the case where $|N_1| = q - i + k$ for $k \in \{0, \dots, i\}$.
636 Since $\sum_{j=1}^q a_j x_j^\circ + \bar{a} \bar{y}^\circ \geq \sum_{j=1}^q a_j x_j^\circ \geq A_q - A_{i-k}$, the objective value associated with (x°, \bar{y}°) satisfies
637 $z^\circ = \sum_{j=1}^q (a_j - \mu)(1 - x_j^\circ) - \bar{a} \bar{y}^\circ \leq A_{i-k} - (i-k)\mu = A_i - i\mu - \sum_{j=i-k+1}^i (a_j - \mu) \leq z'$. Second, consider
638 the case where $|N_1| = q - i - k$ for $k \in \{1, \dots, q - i\}$. Since $\sum_{j=1}^q a_j x_j^\circ + \bar{a} \bar{y}^\circ \geq A_q - A_{i+1} + \theta$ from
639 feasibility, the corresponding objective value is $z^\circ = \sum_{j=1}^q (a_j - \mu)(1 - x_j^\circ) - \bar{a} \bar{y}^\circ \leq A_{i+1} - \theta - (i+k)\mu \leq z^*$.
640 Since whenever the solution (x^*, \bar{y}^*) corresponding to z^* is infeasible, $z^* \leq z'$, the result is proven.

641 2. Assume that $0 \leq w < A_1 - \mu$. An argument similar to that presented above shows that the feasible
642 solution $x_1^* = 0$, $x_j^* = 1$ for $j = 2, \dots, q$, and $\bar{y}^* = \frac{A_1 - \mu - w}{a}$ is optimal for (50), which implies that
643 $P^C(w) = w$.

644 □

645 In the following result, we argue that $P^C(w)$ is subadditive. This result enables us to use Proposition 3.2
646 to perform sequence-independent lifting for the variables in M .

647 **Corollary 3.11.** *The lifting function $P^C(w)$ is subadditive over \mathbb{R}_+ .*

648 *Proof.* In Corollary 3.5, define $\nu = \mu$, $r = q$, and $D_i = A_i$. Since $a_i \geq \mu$ for $i = 1, \dots, q$, it is clear that
649 $A_i \geq A_{i-1} + \mu$. Further, since A_i is defined as the sum of the largest i coefficients in C , it is clear that
650 $A_i + A_j \geq A_{i+j}$ for $0 \leq i, j \leq q$ with $i + j \leq q$. Therefore, Corollary 3.5 shows that $P^C(w)$ is subadditive
651 over \mathbb{R}_+ . □

652 We next illustrate the results of Proposition 3.9, Proposition 3.10, and Corollary 3.11 on an example.

653 **Example 3.12.** *Consider the 0–1 mixed-integer bilinear covering set*

$$654 B = \left\{ (x, y) \in \{0, 1\}^5 \times [0, 1]^5 \mid 21x_1y_1 + 19x_2y_2 + 17x_3y_3 + 15x_4y_4 + 10x_5y_5 \geq 20 \right\}.$$

655 Let $(C', \{l\}, M, T) = (\{5\}, \{4\}, \{1, 2\}, \{3\})$. Clearly, $(C', \{l\}, M, T)$ satisfies Conditions (A1)–(A4) since
656 $C = C' \cup \{l\}$ is a cover with $\mu = 5$, $a_4 \geq a_5$, $a_4 > \mu$ and $\sum_{j \in C \cup T} a_j = 17 + 15 + 10 > 20 + 15 = d + a_l$. By
657 Proposition 3.9, the inequality

$$658 17y_3 + 10x_4 + 5x_5 \geq 15 \tag{51}$$

659 is facet-defining for $PB(M, \emptyset, M, \emptyset)$. Using Proposition 3.10, the lifting function $P^C(w)$ is given by

$$660 P^C(w) = \begin{cases} w & \text{if } 0 \leq w < 10, \\ 10 & \text{if } 10 \leq w < 15, \\ w - 5 & \text{if } 15 \leq w < 20, \\ 15 & \text{if } 20 \leq w. \end{cases}$$

661 Function $P^C(w)$ is represented in Figure 1. Corollary 3.11 shows that this function is subadditive over \mathbb{R}_+ .

662 We now compute the lifting coefficients of variables (x_i, y_i) for $i \in M$ from $P^C(w)$. It follows from
663 Proposition 3.2 that lifting coefficients (α_i, β_i) for $i \in M$ must be chosen to satisfy

$$664 \alpha_i x_i + \beta_i y_i \geq P^C(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{0, 0\}. \tag{52}$$

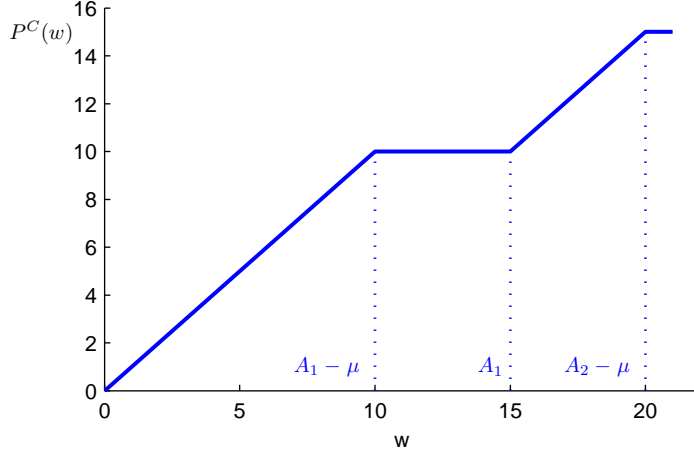


Figure 1: Lifting function $P^C(w)$ of (51)

665 For the problem described in Example 3.12, $P^C(a_i x_i y_i)$ is represented in Figure 2(a) for $i = 1$. In this figure,
 666 we observe that $P^C(a_i x_i y_i)$ is equal to zero when $x_i = 0$ and is equal to $P^C(a_i y_i)$ when $x_i = 1$. Condition (52)
 667 requires that the lifting coefficients (α_i, β_i) be chosen in such a way that the plane $\alpha_i x_i + \beta_i y_i$ (passing through
 668 the origin $(0, 0)$) overestimates the function $P^C(a_i x_i y_i)$ over $\{0, 1\} \times [0, 1]$. Possible overestimating planes
 669 are represented in Figure 2(b). A similar geometric interpretation was used in Richard and Tawarmalani [24]
 670 to obtain lifted inequalities for 0–1 mixed-integer bilinear knapsack sets. It is clear from Figure 2 that good
 671 overestimating planes $\alpha_i x_i + \beta_i y_i$ are in direct correspondence with the concave envelope $p(w)$ of $P^C(w)$ over
 672 $[0, a_i]$. This observation motivates the following result.

673 **Lemma 3.13.** For $i \in M$, define

$$674 \quad q_i := \begin{cases} 0 & \text{if } a_i \leq A_1 - \mu, \\ j & \text{if } A_j - \mu < a_i \leq A_{j+1} - \mu, \quad j = 1, \dots, q-1, \\ q & \text{if } A_q - \mu < a_i. \end{cases}$$

675 Let $Q_0^i = 0$, $Q_j^i = A_j - \mu$ for $j = 1, \dots, q_i$ and $Q_{q_i+1}^i = a_i$. Further, define $\Delta_j^i = Q_{j+1}^i - Q_j^i$ for $j = 0, \dots, q_i$.
 676 Define $p_j^i(w) = P^C(Q_j^i) + \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i}(w - Q_j^i)$ for $j = 0, \dots, q_i$. Then, the function

$$677 \quad p^i(w) := \min \left\{ p_j^i(w) \mid j \in \{0, \dots, q_i\} \right\} \quad (53)$$

678 is a concave overestimator of $P^C(w)$ over $[0, a_i]$.

679 *Proof.* Clearly, $p^i(w)$ is concave since it is defined as the minimum of affine functions. Observe that, for
 680 $j = 0, \dots, q_i$, $p_j^i(w) \geq p_{j+1}^i(w)$ when $w \geq Q_{j+1}^i$, $p_j^i(w) \leq p_{j+1}^i(w)$ when $w < Q_{j+1}^i$, and $p_j^i(w) \geq P^C(w)$
 681 when $w \in [Q_j^i, Q_{j+1}^i]$. Now, consider $j \in \{0, \dots, q_i\}$ and $k \neq j$. Then, for $w \in [Q_j^i, Q_{j+1}^i]$, $P^C(w) \leq p_j^i(w) \leq$
 682 $p_k^i(w)$. \square

683 Observe that the concave overestimator of $P^C(w)$ derived in Lemma 3.13 has $q_i + 1$ linear pieces. Also note
 684 that the definition of q_i implies that $\Delta_j^i > 0$ for all $j = 0, \dots, q_i$. Next, we compute maximal lifting coefficients
 685 for the variables (x_i, y_i) where $i \in M$ using the sequence-independent lifting result of Proposition 3.2 and
 686 Lemma 3.13.

687 **Theorem 3.14.** Under Conditions (A1), (A2), (A3) and (A4), the lifted bilinear cover inequality

$$688 \quad \sum_{j \in C} (a_j - \mu)^+ x_j + \sum_{j \in T} a_j y_j + \sum_{i \in M} \alpha_i x_i + \sum_{i \in M} \beta_i y_i \geq \sum_{j \in C} (a_j - \mu)^+ \quad (54)$$

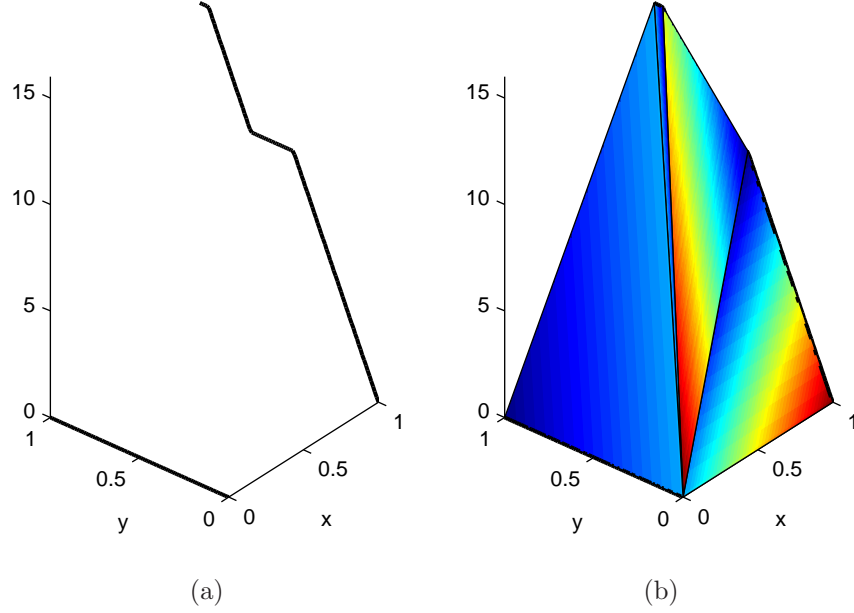


Figure 2: Deriving lifting coefficients for Example 3.15

689 is facet-defining for PB if

$$690 \quad (\alpha_i, \beta_i) \in (P^C(a_i), 0) \bigcup_{j=0}^{q_i} \left(P^C(Q_j^i) - \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} Q_j^i, \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} a_i \right)$$

691 for $i \in M$ in (54) where Q_j^i , Δ_j^i , and q_i are as defined in Lemma 3.13.

692 *Proof.* Because $P^C(w)$ is subadditive over \mathbb{R}_+ , we know that (54) is valid for PB if the lifting coefficients
693 (α_i, β_i) of (x_i, y_i) for $i \in M$ are chosen to satisfy the condition

$$694 \quad \alpha_i x_i + \beta_i y_i \geq P^C(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{0, 0\}. \quad (55)$$

695 Condition (55) can be rewritten as

$$696 \quad \beta_i \phi \geq P^C(0) \quad \text{for } 0 < \phi \leq 1, \quad (56)$$

$$697 \quad \alpha_i + \beta_i \phi \geq P^C(a_i \phi) \quad \text{for } 0 \leq \phi \leq 1. \quad (57)$$

698 To prove that (54) is facet-defining for PB , we also need to show two linearly independent points (x_i, y_i)
699 for which (55) is satisfied at equality. First, consider the case where $(\alpha_i, \beta_i) = (P^C(a_i), 0)$. Condition (56)
700 is satisfied since $\beta_i = 0$ and $P^C(0) = 0$. Condition (57) also holds because $\alpha_i = P^C(a_i)$ and $P^C(w)$ is
701 non-decreasing over \mathbb{R}_+ . Further, (55) is satisfied at equality at the two points, $(0, 1)$ and $(1, 1)$. Finally,
702 consider

$$703 \quad (\alpha_i, \beta_i) = \left(P^C(Q_j^i) - \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} Q_j^i, \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} a_i \right)$$

704 for any $j \in \{0, \dots, q_i\}$. Clearly, (α_i, β_i) satisfies (56) since $\beta_i \geq 0$ and $P^C(0) = 0$. From Lemma 3.13, we
705 have that

$$706 \quad \begin{aligned} P^C(a_i \phi) &\leq P^C(Q_j^i) + \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} (a_i \phi - Q_j^i) \\ &= \left(P^C(Q_j^i) - \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} Q_j^i \right) + \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} a_i \phi \\ &= \alpha_i + \beta_i \phi, \end{aligned}$$

707 showing that (α_i, β_i) satisfy (57) for $j = 0, \dots, q_i$. Further, (55) is satisfied at equality at the two points
 708 $\left(1, \frac{Q_j^i}{a_i}\right)$ and $\left(1, \frac{Q_{j+1}^i}{a_i}\right)$. Therefore, we conclude that (54) is facet-defining for PB . \square

709 The concave overestimator of Lemma 3.13 is in fact the concave envelope of $P^C(w)$ over $w \in [0, a_i]$. The
 710 concave envelope of $P^C(a_i xy)$ over $\{0, 1\} \times [0, 1]$ implicit in the proof of Theorem 3.14 can also be obtained
 711 using the technique for constructing envelopes of functions that satisfy pairwise complementarity described
 712 in [28]. We refer to Section 3 of [28] for definitions and, in particular, Proposition 3 therein for relevant
 713 constructions. The same construction also yields the concave envelope of $\Psi(a_i xy)$ over $\{0, 1\} \times [0, 1]$ proved
 714 later in Theorem 3.26 using the concave overestimator of $\Psi(w)$ derived in Lemma 3.25.

715 Recall that Figure 2(a) depicts $P^C(a_1 x_1 y_1)$ for inequality (51). Observe that in Figure 2(b), lifting
 716 coefficients $(0, a_1)$ define the plane passing through $(0, 0)$ and $(1, 0)$ while lifting coefficients $(P^C(a_i), 0)$
 717 define the plane passing through $(0, 0)$ and $(0, 1)$ (which is identical to the plane obtained when $j = q_i = 2$).
 718 Since there are several choices for the values of each of the pair of lifting coefficients (α_i, β_i) , the family
 719 of inequalities (54) contains an exponential number of members. Theorem 3.14 therefore provides a new
 720 illustration that sequence-independent lifting from a single seed inequality can produce exponentially large
 721 families of inequalities, a property that was discussed in a more general setting in Section 2 of [24]. We
 722 illustrate this characteristic of lifted bilinear cover inequalities in Example 3.15.

723 **Example 3.15.** *In Example 3.12, we established that (51) is facet-defining for $PB(M, \emptyset, M, \emptyset)$ and described*
 724 *the corresponding lifting function $P^C(w)$. We compute that $q_1 = 2$ (with $Q_0^1 = 0, Q_1^1 = 10, Q_2^1 = 20, Q_3^1 = 21$)*
 725 *and $q_2 = 1$ (with $Q_0^2 = 0, Q_1^2 = 10, Q_2^2 = 19$). Applying Theorem 3.14, we obtain the nine inequalities*

$$726 \quad \left\{ \begin{array}{cc} & 21y_1 \\ 5x_1 & + \frac{21}{2}y_1 \\ 15x_1 & \end{array} \right\} + \left\{ \begin{array}{cc} & 19y_2 \\ \frac{50}{9}x_2 & + \frac{76}{9}y_2 \\ 14x_2 & \end{array} \right\} + 17y_3 + 10x_4 + 5x_5 \geq 15$$

727 *which are all facet-defining for PB . The three possible choices for the lifting coefficients of (x_1, y_1) are depicted*
 728 *in Figure 2(b). The fact that there are three possible choices for (x_2, y_2) follows similarly with the exception*
 729 *that coefficient a_2 falls in the second interval $(A_1 - \mu, A_2 - \mu]$.*

730 Another look at Figure 2 also suggests that if we had fixed (x_1, y_1) at $(0, 1)$ or $(1, 0)$, we would only have
 731 been able to obtain a single lifted inequality and so fixing variables at $(0, 0)$ in this case is crucial in discovering
 732 the exponential family of lifted inequalities. This provides a graphical illustration of Proposition 3.3, which
 733 states that all interesting lifting coefficients that can be obtained from fixing variables at $(0, 1)$ or $(1, 0)$ can
 734 also be obtained from fixing variables at $(0, 0)$.

735 3.3.2 Lifted reverse bilinear cover inequalities

736 In Theorem 3.14, we derived lifted bilinear cover inequalities by first lifting the variables (x_j, y_j) for $j \in C'$
 737 and then lifting the variables (x_j, y_j) for $j \in M$. Here, we derive another family of lifted inequalities that we
 738 call *lifted reverse bilinear cover inequalities* by changing the lifting order: we start the lifting procedure with
 739 the same seed inequality (42), but we now lift the variables (x_j, y_j) for $j \in M$ before the variables (x_j, y_j)
 740 for $j \in C'$. In this case, we do not assume that $a_l \geq a_i$ for $i \in C$, *i.e.*, we do not require Condition (A2).

741 **Proposition 3.16.** *Under Conditions (A1), (A3), and (A4), the inequality*

$$742 \quad (a_l - \mu)x_l + \sum_{j \in M} \min\{a_j, a_l - \mu\}x_j + \sum_{j \in T} a_j y_j \geq a_l - \mu \quad (58)$$

743 *is facet-defining for $PB(\emptyset, C', \emptyset, C')$.*

744 *Proof.* It follows from Proposition 2.9 that

$$745 \quad (a_l - \mu)x_l + \sum_{j \in T} a_j y_j \geq a_l - \mu$$

is facet-defining for the full-dimensional polytope $PB(M, C', M, C')$. Its lifting function $P(w)$ is derived in Proposition 3.8 where it is also proven to be subadditive over \mathbb{R}_+ . Therefore, Proposition 3.2 shows that lifting coefficients (α_i, β_i) for (x_i, y_i) for $i \in M$ are valid if they satisfy the condition

$$\alpha_i x_i + \beta_i y_i \geq P(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{0, 0\}. \quad (59)$$

Condition (59) can be rewritten as

$$\beta_i \phi \geq P(0) \quad \text{for } 0 < \phi \leq 1, \quad (60)$$

$$\alpha_i + \beta_i \phi \geq P(a_i \phi) \quad \text{for } 0 \leq \phi \leq 1. \quad (61)$$

We now show that $(\alpha_i, \beta_i) = (P(a_i), 0)$ are valid lifting coefficients. Clearly, $\beta_i = 0$ satisfies (60) since $P(0) = 0$. Further, since $P(a_i \phi) = \min\{a_i \phi, a_l - \mu\}$, it is also clear that $\alpha_i = P(a_i) = \min\{a_i, a_l - \mu\} \geq \min\{a_i \phi, a_l - \mu\} = P(a_i \phi)$. To show that (58) is facet-defining for $PB(\emptyset, C', \emptyset, C')$, it suffices to verify that the two points $(0, 1)$ and $(1, 1)$ satisfy (59) at equality. \square

Proposition 3.16 also follows directly from Proposition 2.9. We provided a proof of Proposition 3.16 based on lifting techniques to emphasize that the cover and reverse-cover inequalities are obtained by reversing the order of lifting of M and C' . We remark that the above result does not require Condition (A2). Also, note that lifting coefficients $(\alpha_i, \beta_i) = (0, a_i)$ for $i \in M$ are valid for (59). These coefficients yield facet-defining inequalities for $PB(\emptyset, C', \emptyset, C')$ because (59) is satisfied at equality for $(1, 0)$ and $(1, \min\{1, \frac{a_l - \mu}{a_i}\})$. However, these variables could have been treated directly as elements of T in (42) since adding more elements to T will not violate Condition (A4).

To obtain facet-defining inequalities for PB , we lift the remaining variables (x_j, y_j) for $j \in C'$ in (58). To this end, we first compute the function

$$P^M(w) := \max (a_l - \mu) - \left\{ (a_l - \mu)x_l + \sum_{j \in M} \min\{a_j, a_l - \mu\}x_j + \sum_{j \in T} a_j y_j \right\}$$

$$s.t. \quad a_l x_l y_l + \sum_{j \in M \cup T} a_j x_j y_j \geq a_l - \mu - w \quad (62)$$

$$x_j \in \{0, 1\}, y_j \in [0, 1] \quad \forall j \in \{l\} \cup M \cup T.$$

Let $M = M_1 \cup M_2$ where $M_1 = \{i \in M \mid a_i > a_l - \mu\}$ and $M_2 = M \setminus M_1$. Assume without loss of generality that $\{l\} \cup M_1 = \{1, \dots, q\}$ and $a_1 \geq a_2 \geq \dots \geq a_q$ where $q = |M_1| + 1$. Further, define $A_0 = 0$ and $A_i = \sum_{j=1}^i a_j$ for $i = 1, \dots, q$. Observe that $a_l + \sum_{j \in M \cup T} a_j = A_q + \sum_{j \in M_2} a_j + \sum_{j \in T} a_j$. We derive a closed-form expression for $P^M(w)$ in the following proposition.

Proposition 3.17.

$$P^M(w) = \begin{cases} -\infty & \text{if } w < -\mu - \sum_{j \in M \cup T} a_j, \\ w + A_q - q(a_l - \mu) & \text{if } -\mu - \sum_{j \in M \cup T} a_j \leq w < -A_q + (a_l - \mu), \\ -i(a_l - \mu) & \text{if } -A_{i+1} + (a_l - \mu) \leq w < -A_i, & i = 0, \dots, q-1, \\ w + A_i - i(a_l - \mu) & \text{if } -A_i \leq w < -A_i + (a_l - \mu), & i = 1, \dots, q-1, \end{cases}$$

Proof. First, we observe that, if (62) has a feasible solution, then it has an optimal solution (x^*, y^*) that satisfies $x_j^* = 1$ for $j \in T$ and $y_j^* = 1$ for $j \in M \cup \{l\}$ since the objective coefficients corresponding to these variables are zero. Using the notation $\bar{a} = \sum_{j \in T} a_j$ and $\bar{y} = \frac{\sum_{j \in T} a_j y_j}{\bar{a}}$, we simplify the expression of $P^M(w)$ as

$$P^M(w) = \max (a_l - \mu) - \left\{ \sum_{j \in \{l\} \cup M_1} (a_l - \mu)x_j + \sum_{j \in M_2} a_j x_j + \bar{a}\bar{y} \right\}$$

$$s.t. \quad \sum_{j \in \{l\} \cup M_1} a_j x_j + \sum_{j \in M_2} a_j x_j + \bar{a}\bar{y} \geq a_l - \mu - w \quad (63)$$

$$x_j \in \{0, 1\} \quad \forall j \in \{l\} \cup M_1 \cup M_2, \bar{y} \in [0, 1].$$

781 After introducing $\hat{a} = \sum_{j \in M_2} a_j + \bar{a}$ and $\hat{y} = \frac{\sum_{j \in M_2} a_j x_j + \bar{a} \bar{y}}{\hat{a}}$, we claim that $P^M(w)$ can be written as

$$\begin{aligned}
782 \quad P^M(w) = \quad & \max \quad (a_l - \mu) - \left\{ \sum_{j=1}^q (a_l - \mu) x_j + \hat{a} \hat{y} \right\} \\
783 \quad & \text{s.t.} \quad \sum_{j=1}^q a_j x_j + \hat{a} \hat{y} \geq a_l - \mu - w \\
784 \quad & x_j \in \{0, 1\} \quad \forall j \in \{1, \dots, q\}, \quad \hat{y} \in [0, 1].
\end{aligned} \tag{64}$$

785 We next prove that (63) and (64) are equivalent. To do so, we show that (63) has a feasible solution
786 $(x_l^*, x_{M_1}^*, x_{M_2}^*, \bar{y}^*)$ with objective value ζ^* if and only if (64) has a feasible solution $(x_l^*, x_{M_1}^*, \hat{y}^*)$ with objective
787 value ζ^* . On the one hand, given $(x_l^*, x_{M_1}^*, x_{M_2}^*, \bar{y}^*)$, we can obtain $(x_l^*, x_{M_1}^*, \hat{y}^*)$ directly from the definition of
788 \hat{y} . The objective values of these two solutions are identical. On the other hand, let $M_2 = \{q+1, \dots, m\}$. Define
789 $\hat{A}_0 = 0$ and $\hat{A}_i = \sum_{j=q+1}^{q+i} a_j$ for $i = 1, \dots, m - q$. Then, for a given $(x_l^*, x_{M_1}^*, \hat{y}^*)$, we build $(x_l^*, x_{M_1}^*, x_{M_2}^*, \bar{y}^*)$
790 as follows. Define $\hat{m} = \max\{i \in \{0, \dots, m - q\} \mid \hat{A}_i \leq \hat{a} \hat{y}^*\}$ and set $x_{q+j}^* = 1$ for $j \leq \hat{m}$, $x_{q+j}^* = 0$ for $j > \hat{m}$
791 and $\bar{y}^* = \frac{\hat{a} \hat{y}^* - \hat{A}_{\hat{m}}}{\bar{a}}$. We argue next that this solution is feasible. First observe that $\hat{a} \hat{y}^* - \hat{A}_{\hat{m}} \leq a_{q+\hat{m}+1}$
792 when $\hat{m} \leq m - q - 1$ and that $\hat{a} \hat{y}^* - \hat{A}_{\hat{m}} \leq \bar{a}$ when $\hat{m} = m - q$. Since $\bar{a} = \sum_{j \in T} a_j > a_l - \mu \geq a_i$ for all
793 $i \in M_2$ because of Condition (A4) and the definition of M_2 , we easily conclude that $0 \leq \frac{\hat{a} \hat{y}^* - \hat{A}_{\hat{m}}}{\bar{a}} \leq 1$. Also,
794 $\sum_{j \in \{l\} \cup M_1} a_j x_j^* + \sum_{j \in M_2} a_j x_j^* + \bar{a} \bar{y}^* = \sum_{j \in \{l\} \cup M_1} a_j x_j^* + \hat{A}_{\hat{m}} + \hat{a} \hat{y}^* - \hat{A}_{\hat{m}} = \sum_{j \in \{l\} \cup M_1} a_j x_j^* + \hat{a} \hat{y}^*$. This
795 shows that the proposed solution is feasible for (63) and has the same objective value as $(x_l^*, x_{M_1}^*, \hat{y}^*)$.

796 Next, we study (64). It is clear that this problem is infeasible if and only if $w < a_l - \mu - A_q - \hat{a} =$
797 $-\mu - \sum_{j \in M \cup T} a_j$. Therefore assume that $w \geq -\mu - \sum_{j \in M \cup T} a_j$. Consider now any optimal solution (x^*, \hat{y}^*)
798 for which $x_i^* < x_t^*$ and $i < t$ for some $i, t \in \{1, \dots, q\}$. Then the solution (\bar{x}, \hat{y}^*) where $\bar{x}_k = x_k^*$ if $k \neq i$
799 and $k \neq t$, $\bar{x}_i = x_t^*$, and $\bar{x}_t = x_i^*$ is also feasible for (64) since $a_i \geq a_t$ and has the same objective value as
800 (x^*, \hat{y}^*) . It follows that (64) has an optimal solution that satisfies $x_1^* = \dots = x_i^* = 1$ and $x_{i+1}^* = \dots = x_q^* = 0$
801 for some $i \in \{1, \dots, q\}$. Consider such a solution further. On the one hand, if $\sum_{j=1}^i a_j \geq a_l - \mu - w$, then
802 $\sum_{j=1}^{i-1} a_j < a_l - \mu - w$ and $\hat{y}^* = 0$ otherwise the solution $x_j^\circ = 1$ for $j = 1, \dots, i-1$, $x_j^\circ = 0$ for $j = i, \dots, q$ and
803 $\hat{y}^\circ = 0$ would be feasible and would have a better objective value. On the other hand if $\sum_{j=1}^i a_j < a_l - \mu - w$
804 for $i \leq q - 1$ then $\sum_{j=1}^{i+1} a_j \geq a_l - \mu - w$. Otherwise the solution $x_j^\circ = 1$ for $j = 1, \dots, i+1$, $x_j^\circ = 0$ for
805 $j = i+2, \dots, q$ and $\hat{y}^\circ = \hat{y}^* - \frac{a_{i+1}}{\hat{a}}$ would be feasible and would have an objective value $a_{i+1} - (a_l - \mu)$ larger
806 than that of (x^*, \hat{y}^*) . This is a contradiction since $a_{i+1} > a_l - \mu$.

807 We consider two situations. First, assume $-A_q + (a_l - \mu) - \hat{a} \leq w < -A_q + (a_l - \mu)$, it follows from
808 the above discussion that there is an optimal solution (x^*, \hat{y}^*) with $x^* = \mathbf{1}$. Then $\hat{y}^* = \frac{a_l - \mu - w - A_q}{\hat{a}}$. Clearly,
809 $\hat{y}^* \in [0, 1]$ and so $P^M(w) = w + A_q - q(a_l - \mu)$. Second, assume $-A_{i+1} + (a_l - \mu) \leq w < -A_i + (a_l - \mu)$ for
810 some $i \in \{0, \dots, q-1\}$, it follows from the above discussion that one of the following two solutions

$$\begin{aligned}
811 \quad x_1^\wedge = x_2^\wedge = \dots = x_{i+1}^\wedge = 1, \quad x_{i+2}^\wedge = \dots = x_q^\wedge = 0, \quad \hat{y}^\wedge = 0, \quad \text{and} \\
812 \quad x_1^\ddagger = x_2^\ddagger = \dots = x_i^\ddagger = 1, \quad x_{i+1}^\ddagger = \dots = x_q^\ddagger = 0, \quad \hat{y}^\ddagger = \frac{a_l - \mu - w - A_i}{\hat{a}}
\end{aligned}$$

813 with objective values $z^\wedge = -i(a_l - \mu)$ and $z^\ddagger = -i(a_l - \mu) + (w + A_i)$ is optimal for (64) since $a_l - \mu - w \in$
814 $(A_i, A_{i+1}]$. Note that the second solution is feasible only when $a_l - \mu - w - A_i \leq \hat{a}$. We now consider two
815 cases. When $w \leq -A_i$ then $z^\wedge \geq z^\ddagger$ and so $P^M(w) = -i(a_l - \mu)$. When $w > -A_i$, then $z^\ddagger > z^\wedge$. Further,
816 solution $(x^\ddagger, \hat{y}^\ddagger)$ is feasible since $a_l - \mu - w - A_i < a_l - \mu \leq \hat{a}$ because of Condition (A4). It follows that
817 $P^M(w) = -i(a_l - \mu) + (w + A_i)$. \square

818 To perform sequence-independent lifting for the variables (x_j, y_j) for $j \in C'$, we verify that the function
819 $P^M(w)$ is subadditive over \mathbb{R}^- .

820 **Proposition 3.18.** *The lifting function $P^M(w)$ is subadditive over \mathbb{R}_- .*

821 *Proof.* First, note that $P^M(w)$ is subadditive over \mathbb{R}_- if it is subadditive over $I = [-\mu - \sum_{j \in M \cup T} a_j, 0]$.
822 Consider Corollary 3.5 and define $D_i = A_i$, $\nu = a_l - \mu$, $r = q$, and notice that $P^M(w) = g(-w) + w$. Clearly,
823 $A_i + A_j \geq A_{i+j}$ for $0 \leq i \leq j \leq q$ with $i + j \leq q$ since A_i is the sum of the largest i coefficients in $M_1 \cup \{l\}$.
824 It then follows from Corollary 3.5 that $P^M(w)$ is subadditive over I , proving the result. \square

825 We next illustrate the results of Propositions 3.16, 3.17, and 3.18 via an example.

826 **Example 3.19.** For the set B of Example 3.12, consider the partition $(C', \{l\}, M, T) = (\{3\}, \{4\}, \{5\}, \{1, 2\})$.
827 This partition satisfies Conditions (A1), (A3), and (A4) since C is a cover with $\mu = 12$, $a_4 > \mu$, and
828 $\sum_{j \in C \cup T} a_j = 21 + 19 + 17 + 15 > 20 + 15 = d + a_l$. We obtain from Proposition 3.16 that

$$829 \quad 3x_4 + 3x_5 + 21y_1 + 19y_2 \geq 3 \quad (65)$$

830 is facet-defining for $PB(\emptyset, C', \emptyset, C')$. Further, the lifting function $P^M(w)$ over \mathbb{R}_- is given by

$$831 \quad P^M(w) = \begin{cases} -\infty & \text{if } w < -62 \\ w + 19 & \text{if } -62 \leq w < -22 \\ -3 & \text{if } -22 \leq w < -15 \\ w + 12 & \text{if } -15 \leq w < -12 \\ 0 & \text{if } -12 \leq w \leq 0, \end{cases}$$

832 as described in Proposition 3.17 since $q = 2$, $A_0 = 0$, $A_1 = 15$ and $A_2 = 25$.

833 Similar to Theorem 3.14, we compute the lifting coefficients for the variables (x_i, y_i) for $i \in C'$ using
834 sequence-independent lifting; refer to the discussion following Proposition 3.2.

835 **Theorem 3.20.** Suppose that Conditions (A1), (A3), and (A4) hold. Then, the lifted reverse bilinear cover
836 inequality

$$837 \quad (a_l - \mu)x_l - \sum_{j \in C'} P^M(-a_j)x_j + \sum_{j \in M} \min\{a_j, a_l - \mu\}x_j + \sum_{j \in T} a_j y_j \geq (a_l - \mu) - \sum_{j \in C'} P^M(-a_j) \quad (66)$$

838 is facet-defining for PB .

839 *Proof.* Since $P^M(w)$ is subadditive over \mathbb{R}_- , the lifting coefficients (α_i, β_i) of the variables (x_i, y_i) for $i \in C'$
840 are valid if they are chosen to satisfy

$$841 \quad \alpha_i(x_i - 1) + \beta_i(y_i - 1) \geq P^M(a_i x_i y_i - a_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{1, 1\}. \quad (67)$$

842 Condition (67) can be rewritten as

$$843 \quad \beta_i \leq \inf_{0 \leq \phi < 1} \frac{-P^M(a_i \phi - a_i)}{1 - \phi}, \quad (68)$$

$$844 \quad \alpha_i + \sup_{0 \leq \phi \leq 1} \beta_i(1 - \phi) \leq -P^M(-a_i). \quad (69)$$

845 Because of Assumption 2, we know that $a_i \leq \sum_{j \in N} a_j - d = \sum_{j \in C \cup M \cup T} a_j - (\sum_{j \in C} a_j - \mu) = \mu + \sum_{j \in M \cup T} a_j$
846 for all $i \in C' \subseteq N$ and so $P^M(a_i \phi - a_i) > -\infty$ for all $\phi \in [0, 1)$. Choosing $\beta_i = 0$ satisfies (68) since
847 $P^M(a_i \phi - a_i) \leq 0$ for all $\phi \in [0, 1)$. Moreover, as $\beta_i = 0$, it is easily verified that choosing $\alpha_i = -P^M(-a_i)$
848 satisfies (69). Finally, note that (67) is tight at the points $(0, 0)$ and $\left(1, \frac{(a_i - A_i + a_l - \mu)^+}{a_i}\right)$, which proves that
849 (66) is facet-defining for PB . \square

850 Note that the lifted reverse bilinear cover inequality (66) we obtained through lifting is unique. This is a
851 significant difference from lifted bilinear cover inequalities (54). We next illustrate in an example the reason
852 that we obtain a single lifted reverse bilinear cover inequality and show that not all lifted reverse bilinear
853 cover inequalities (66) can be derived as lifted bilinear cover inequalities (54).

854 **Example 3.21.** For the partition $(C', \{l\}, M, T) = (\{3\}, \{4\}, \{5\}, \{1, 2\})$, we established in Example 3.19
 855 that (65) is facet-defining for $PB(\emptyset, C', \emptyset, C')$. Applying Theorem 3.20, we obtain the following lifted reverse
 856 bilinear cover inequality

$$857 \quad 3x_3 + 3x_4 + 3x_5 + 21y_1 + 19y_2 \geq 6, \quad (70)$$

858 which is facet-defining for PB . We represent in Figure 3(a), the function $P^M(a_3x_3y_3 - a_3)$ that was overes-
 859 timated to construct valid lifting coefficients. We represent in Figure 3(b) the only choice of coefficients that
 860 yields an overestimating plane to $P^M(a_3x_3y_3 - a_3)$ over $(x_3, y_3) \in \{0, 1\} \times [0, 1]$ and is tight at $(1, 1)$. Further,
 861 Inequality (70) cannot be obtained as a lifted bilinear cover inequality (54). In fact, if (70) was of the form
 862 (54), it should be that $C \subseteq \{3, 4, 5\}$. However, none of the four possible covers $C_1 = \{3, 4\}$, $C_2 = \{3, 5\}$,
 863 $C_3 = \{4, 5\}$ and $C_4 = \{3, 4, 5\}$ yield (70).

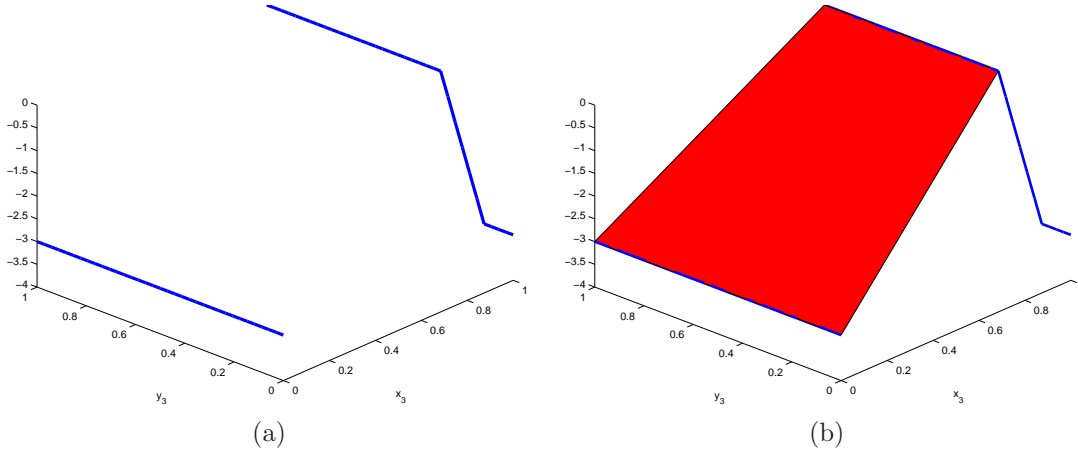


Figure 3: Deriving lifting coefficients for Example 3.21

864 3.4 Inequalities through approximate lifting

865 We now derive another family of lifted inequalities from the seed inequality (27) developed in Proposition 2.11.
 866 To this end, we first identify a partition (K, \overline{M}) of the set of variables N that satisfies the following conditions

$$867 \quad (C1) \quad \sum_{j \in K} a_j - a_k \geq d \text{ for all } k \in K,$$

$$868 \quad (C2) \quad \sum_{j \in K} a_j - a_k - a_m < d \text{ for all } k \neq m \in K, \text{ i.e., } a_k + a_m > \mu \text{ for all } k \neq m \in K,$$

869 where $\mu = \sum_{j \in K} a_j - d$ is the excess of K . Note that Condition (C1) implies that K is a cover. Further,
 870 Condition (C1) requires that $K \setminus \{k\}$ is also a cover for all $k \in K$ and so $a_k \leq \mu$ for all $k \in K$. It also follows
 871 from Condition (C1) that $|K| \geq 2$. We refer to a set K satisfying Conditions (C1) and (C2) as a *clique*. After
 872 fixing the variables (x_i, y_i) for $i \in \overline{M}$ to $(0, 0)$, it follows from Proposition 2.11 that the clique inequality

$$873 \quad \sum_{j \in K} x_j \geq |K| - 1 \quad (71)$$

874 is facet-defining for $PB(\overline{M}, \emptyset, \overline{M}, \emptyset)$.

875 We now lift the remaining variables (x_i, y_i) for $i \in \overline{M}$ in two steps. We assume without loss of generality
 876 that $K = \{1, \dots, r\}$ and that $a_1 \leq a_2 \leq \dots \leq a_r$. We define $\mu' = a_1 + a_2 - \mu$. We assume that $a_{r+1} \leq \dots \leq a_n$
 877 and define p such that $\sum_{i=r+1}^p a_i < \mu' \leq \sum_{i=r+1}^{p+1} a_i$. (More generally, \overline{M} can be taken to be any subset of
 878 \overline{M} such that $\sum_{i \in \widehat{M}} a_i < \mu'$ without altering the form of the derived inequality.) Let $\widehat{M} = \{a_{r+1}, \dots, a_p\}$.
 879 We show that (71) is facet-defining for $PB(\overline{M} \setminus \widehat{M}, \emptyset, \overline{M} \setminus \widehat{M}, \emptyset)$. First, we show by contradiction that the

inequality is valid. Let (x, y) be such that $\sum_{j \in K} x_j < r - 1$. Then,

$$\sum_{j=1}^p a_j x_j y_j \leq \sum_{j=3}^p a_j = d - \mu' + \sum_{j=r+1}^p a_j < d,$$

where the first inequality holds since $a_1 \leq \dots \leq a_r$ and $\sum_{j \in K} x_j < r - 1$ and the last inequality follows since $\sum_{j=r+1}^p a_j < \mu'$. This inequality implies that $(x, y) \notin B(\overline{M} \setminus \widehat{M}, \emptyset, \overline{M} \setminus \widehat{M}, \emptyset)$, the desired contradiction. By Proposition 2.10, it suffices to show that (71) is facet-defining for $PB(\overline{M} \setminus \widehat{M}, \emptyset, \overline{M} \setminus \widehat{M}, K \cup \widehat{M})$. Define χ such that $\chi_j = 1$ for $j \in K$ and $\chi_j = 0$ for $j \in \widehat{M}$. Then, by (C1), $p^k = \chi - e_k$ for $k \in K$ and $q^k = \chi - e_1 + e_k$ for $k \in \widehat{M}$, are feasible. Since these points are linearly independent, (71) is facet-defining.

We now lift variables (x_i, y_i) for $i \in M = \overline{M} \setminus \widehat{M}$. The lifting function corresponding to (71) is defined as

$$\begin{aligned} \Phi(w) := \max \quad & (|K| - 1) - \sum_{j \in K} x_j \\ \text{s.t.} \quad & \sum_{j \in K} a_j x_j y_j + \sum_{j \in \widehat{M}} a_j x_j y_j \geq d - w \\ & x_j \in \{0, 1\}, y_j \in [0, 1] \quad \forall j \in K. \end{aligned} \tag{72}$$

We define $a' = \sum_{j \in \widehat{M}} a_j$, $\bar{\mu} = \mu' - a'$, $B_0 = 0$, and $B_i = \sum_{j=1}^i a_{j+2} - a'$ for $i = 1, \dots, r-2$. It follows from the definition of \widehat{M} that $\bar{\mu} > 0$. Observe that $B_0 \leq B_1$ because $a_3 - a' \geq a_3 - \mu' = a_3 - a_1 - a_2 + \mu \geq -a_2 + \mu \geq 0$, where the last inequality follows from (C1). Also, observe that $B_{r-2} + \bar{\mu} = d - a'$ and, for all $i \in M$, $a_i \geq a_{p+1} \geq \mu' - a' = \bar{\mu}$, where the last inequality follows from the definition of \widehat{M} .

Proposition 3.22. For $w \geq 0$,

$$\Phi(w) = \begin{cases} 0 & \text{if } 0 \leq w < \bar{\mu}, \\ i & \text{if } B_{i-1} + \bar{\mu} \leq w < B_i + \bar{\mu}, \quad i = 1, \dots, r-2, \\ r-1 & \text{if } B_{r-2} + \bar{\mu} \leq w. \end{cases}$$

Proof. Problem (72) is feasible for $w \geq 0$ and has an optimal solution (x^*, y^*) that satisfies $(x_j^*, y_j^*) = 1$ for $j \in \widehat{M}$ and $y_j^* = 1$ for $j \in K$ since the objective coefficients of these variables are zero. Hence, $\Phi(w)$ can be rewritten as

$$\begin{aligned} \Phi(w) = \max \quad & (|K| - 1) - \sum_{j \in K} x_j \\ \text{s.t.} \quad & \sum_{j \in K} a_j x_j \geq d - a' - w \\ & x_j \in \{0, 1\} \quad \forall j \in K. \end{aligned} \tag{73}$$

Further, we claim that there exists an optimal solution x^* to (73) in which

$$x_1^* \leq x_2^* \leq \dots \leq x_r^*. \tag{74}$$

This is because, given any solution x^* to (73) with $x_i^* > x_j^*$ for $i < j$, the solution \bar{x} defined as $\bar{x}_k = x_k^*$ if $k \neq i$ and $k \neq j$, $\bar{x}_i = x_j^*$, and $\bar{x}_j = x_i^*$, is feasible and has the same objective value. It follows from (74) that, given $w \in [0, d - a']$, the solution

$$x_j^* = \begin{cases} 0 & \text{if } j = 1, \dots, t(w) - 1, \\ 1 & \text{if } j = t(w), \dots, r, \end{cases}$$

where $t(w) = \max\{i \mid \sum_{j=i}^r a_j \geq d - a' - w = B_{r-2} + \bar{\mu} - w\}$ is optimal for (73) and has an objective value of $t(w) - 2$. When $w \in [0, \bar{\mu})$, $d - a' - w \in (B_{r-2}, B_{r-2} + \bar{\mu}]$ showing that $t(w) = 2$ and $\Phi(w) = 0$. When $w \in [B_{i-1} + \bar{\mu}, B_i + \bar{\mu})$ for $i = 1, \dots, r-2$, $d - a' - w \in (\sum_{j=i+3}^r a_j, \sum_{j=i+2}^r a_j]$ showing that $t(w) = i + 2$ and $\Phi(w) = i$. Finally, when $w \geq d - a'$, it is clear that $\Phi(w) = |K| - 1 = r - 1$. \square

913 In Section 3.3.1, all lifting functions were subadditive over appropriate ranges. As a result, strong valid
914 inequalities for PB were easily obtained using sequence-independent lifting. The lifting function $\Phi(w)$ derived
915 in Proposition 3.22, however, is not subadditive. To circumvent the difficulties associated with sequence-
916 dependent lifting in such a situation, Gu et al. [13] proposed to use approximate lifting. Following their
917 approach, we say that $\Psi(w)$ is a *valid subadditive approximation* of $\Phi(w)$ if $\Psi(w) \geq \Phi(w)$ for all $w \in \mathbb{R}_+$ and
918 $\Psi(w)$ is subadditive. We say that a valid subadditive approximation $\Psi(w)$ is *nondominated* if there is no
919 other valid subadditive approximation $\Psi'(w)$ of $\Phi(w)$ with $\Psi'(w) \leq \Psi(w)$ for all $w \in \mathbb{R}_+$ and $\Psi'(w') < \Psi(w')$
920 for some $w' \in \mathbb{R}_+$. We also define the notion of maximal set $E = \{w \in \mathbb{R}_+ \mid \Phi^i(w) = \Phi(w) \forall i \in$
921 $M, \text{ for all coefficients } a_i \in \mathbb{R}_+ \text{ and for all lifting orders}\}$. A valid subadditive approximation $\Psi(w)$ of $\Phi(w)$ is
922 called *maximal* if $\Psi(w) = \Phi(w)$ for all $w \in E$. It is clear that a maximal nondominated approximation of Φ
923 leads to strong inequalities that can be obtained efficiently for PB . The approximation of $\Phi(w)$ we use has
924 the form presented in Corollary 3.6.

925 We next describe in Proposition 3.23 a subadditive, nondominated and maximal approximation of $\Phi(w)$
926 over \mathbb{R}_+ .

927 **Proposition 3.23.** *The function*

$$928 \quad \Psi(w) := \begin{cases} i + \frac{w-B_i}{\bar{\mu}} & \text{if } B_i \leq w < B_i + \bar{\mu}, \quad i = 0, \dots, r-2, \\ i & \text{if } B_{i-1} + \bar{\mu} \leq w < B_i, \quad i = 1, \dots, r-2, \\ r-1 & \text{if } B_{r-2} + \bar{\mu} \leq w, \end{cases}$$

929 *is a valid subadditive approximation of $\Phi(w)$ that is nondominated and maximal over \mathbb{R}_+ .*

930 *Proof.* Note that $\Psi(w) = \Phi(w)$ when $w \in [B_{i-1} + \bar{\mu}, B_i]$ for $i \in \{1, \dots, r-2\}$ and when $w \geq B_{r-2} + \bar{\mu}$.
931 Further,

$$932 \quad \Psi(w) = \Phi(w) + \frac{w - B_i}{\bar{\mu}} \geq \Phi(w)$$

933 when $w \in (B_i, B_i + \bar{\mu})$ for $i \in \{0, \dots, r-2\}$. Next, we show that $\Psi(w)$ is subadditive over \mathbb{R}_+ . In Corollary 3.6,
934 let $s = r-2$, $C_i = B_i$ and $\lambda = \bar{\mu}$. Since B_i is the sum of the smallest i coefficients in $K \setminus \{1, 2\}$, it is clear
935 that $B_i + B_j \leq B_{i+j}$ for $0 \leq i \leq j \leq r-2$ with $i+j \leq r-2$. Therefore, $\Psi(w)$ is subadditive over \mathbb{R}_+ . We
936 now argue nondominance and maximality over \mathbb{R}_+ . To this end, we first observe that for all $w' \in \mathbb{R}_+$ there
937 exists $w'' \in \mathbb{R}_+$ such that

$$938 \quad \Psi(w') + \Psi(w'') = \Phi(w' + w''). \quad (75)$$

939 In particular, w'' can be chosen to be $B_i + \bar{\mu} - w'$ when $w' \in (B_i, B_i + \bar{\mu})$ and w'' can be chosen to be 0 otherwise.
940 If Ψ' dominates Ψ strictly at w' then $\Psi'(w' + w'') \leq \Psi'(w') + \Psi'(w'') < \Psi(w') + \Psi(w'') = \Phi(w' + w'')$ yielding
941 a contradiction to the assumption that Ψ' is an overestimator of Φ . Similarly, if $\Phi(w') < \Psi(w')$ then (75)
942 implies that $\Phi(w' + w'') - \Phi(w') > \Psi(w'') \geq \Phi(w'')$. Therefore, $\Phi(w'')$ does not yield a valid lifting coefficient
943 for the sequential perturbation of w'' after w' . \square

944 **Example 3.24.** *For the bilinear set B studied in Example 3.12, consider $K = \{3, 4, 5\}$. Set K satisfies*
945 *Conditions (C1) and (C2) with $\mu = 22$. It follows from Proposition 2.11 that*

$$946 \quad x_3 + x_4 + x_5 \geq 2 \quad (76)$$

947 *is facet-defining for $B(\{1, 2\}, \emptyset, \{1, 2\}, \emptyset)$. Let $\widehat{M} = \emptyset$. The lifting function of (76) obtained using Proposi-*
948 *tion 3.22 and its valid subadditive approximation $\Psi(w)$ obtained in Proposition 3.23 are given by*

$$949 \quad \Phi(w) = \begin{cases} 0 & \text{if } 0 \leq w < 3, \\ 1 & \text{if } 3 \leq w < 20, \\ 2 & \text{if } 20 \leq w \end{cases} \quad \text{and} \quad \Psi(w) = \begin{cases} \frac{w}{3} & \text{if } 0 \leq w < 3, \\ 1 & \text{if } 3 \leq w < 17, \\ 1 + \frac{w-17}{3} & \text{if } 17 \leq w < 20, \\ 2 & \text{if } 20 \leq w \end{cases}$$

950 *as $r = 3$, $\bar{\mu} = 3$, $B_0 = 0$, and $B_1 = 17$.*

951 In Figure 4, we present the lifting function $\Phi(w)$ of the clique inequality derived in Proposition 3.22 and
 952 its valid subadditive approximation $\Psi(w)$ obtained in Proposition 3.23 for the particular case of inequality
 953 (76) discussed in Example 3.24. The function $\Phi(w)$ is depicted with a dotted line while $\Psi(w)$ is represented
 954 using a solid line. Observe that, for $0 < w \leq \bar{\mu} = 3$, the approximation is exact only when $w = \bar{\mu} = 3$, *i.e.*,
 955 $\Psi(\bar{\mu}) = \Phi(\bar{\mu})$. For $w \geq \bar{\mu} = 3$, the approximation is exact when $3 = \bar{\mu} \leq w \leq B_1 = 17$ and $w \geq B_1 + \bar{\mu} = 20$.
 956 Next, we obtain a concave overestimator of $\Psi(w)$ in Lemma 3.25 that we will use in Theorem 3.26 to compute
 957 lifting coefficients for the variables in M .

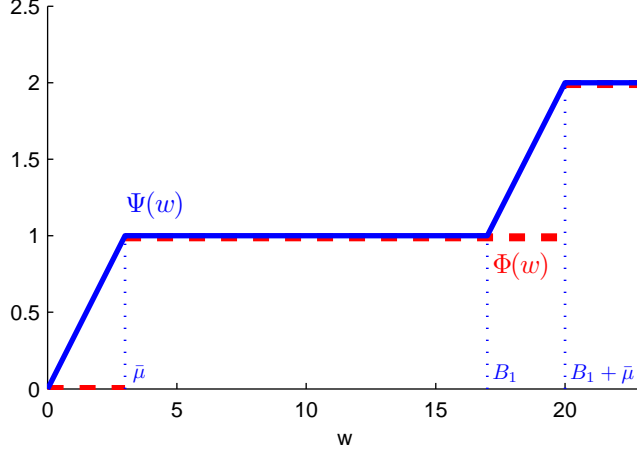


Figure 4: A valid subadditive approximation $\Psi(w)$ of $\Phi(w)$ for Example 3.24.

958 **Lemma 3.25.** For $i \in M$, define

$$959 \quad q_i := \begin{cases} 0 & \text{if } a_i \leq \bar{\mu}, \\ j+1 & \text{if } B_j + \bar{\mu} < a_i \leq B_{j+1} + \bar{\mu}, \quad j = 0, \dots, r-3, \\ r-1 & \text{if } B_{r-2} + \bar{\mu} < a_i. \end{cases}$$

960 Let $W_0^i = 0$, $W_j^i = B_{j-1} + \bar{\mu}$ for $j = 1, \dots, q_i$ and $W_{q_i+1}^i = a_i$. Define $\Delta_j^i = W_{j+1}^i - W_j^i$ for $j = 0, \dots, q_i$.

961 Define also $\psi_j^i(w) = \Psi(W_j^i) + \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i}(w - W_j^i)$ for $j = 0, \dots, q_i$. Then, the function

$$962 \quad \psi^i(w) := \min \left\{ \psi_j^i(w) \mid j \in \{0, \dots, q_i\} \right\} \quad (77)$$

963 is a concave overestimator of $\Psi(w)$ over $[0, a_i]$.

964 *Proof.* First, $\psi^i(w)$ is concave since it is obtained as the minimum of affine functions. Observe that, for
 965 $j = 0, \dots, q_i$, $\psi_j^i(w) \geq \psi_{j+1}^i(w)$ when $w \geq W_{j+1}^i$, $\psi_j^i(w) \leq \psi_{j+1}^i(w)$ when $w < W_{j+1}^i$, and $\psi_j^i(w) \geq \Psi(w)$
 966 when $w \in [W_j^i, W_{j+1}^i]$. Now, consider $j \in \{0, \dots, q_i\}$ and $k \neq j$. Then, for $w \in [W_j^i, W_{j+1}^i]$, $\Psi(w) \leq \psi_j^i(w) \leq$
 967 $\psi_k^i(w)$. \square

968 The concave overestimator $\psi^i(w)$ of Lemma 3.25 can be used to obtain lifting coefficients in a manner
 969 similar to that of Theorem 3.14. Because of the way the concave overestimator is built, it can be observed that
 970 all of its affine pieces (except possibly $\psi_{q_i}^i$) touch the original lifting function Φ at two points and therefore
 971 can be used to generate strong lifting coefficients. To describe whether $\psi_{q_i}^i$ touches Φ in two points, we define
 972 $I(a_i)$ to be the function that returns 0 if $\Phi(a_i) = \Psi(a_i)$ and returns 1 otherwise, *i.e.*,

$$973 \quad I(a_i) := \begin{cases} 0 & \text{if } B_{q_i-1} + \bar{\mu} < a_i \leq B_{q_i} \text{ or } a_i > B_{r-2} + \bar{\mu}, \\ 1 & \text{if } B_{q_i} < a_i \leq B_{q_i} + \bar{\mu}. \end{cases}$$

974 We observe that, when $I(a_i) = 0$, it is possible to derive maximal lifting coefficients (with respect to Φ) from
 975 all affine pieces of ψ^i . When $I(a_i) = 1$, however, we can only guarantee the derivation of maximal lifting

976 coefficients (with respect to Φ) from ψ^i for $i = 0, \dots, q_i - 1$. This intuitive observation is formally proven in
 977 the following theorem.

978 **Theorem 3.26.** *Under Conditions (C1) and (C2),*

$$979 \quad \sum_{j \in K} x_j + \sum_{i \in M} \alpha_i^{j_i} x_i + \sum_{i \in M} \beta_i^{j_i} y_i \geq |K| - 1 \quad (78)$$

980 defines a face of PB of dimension at least $(2n - 1) - \sum_{i \in M} I(a_i) \times \mathbf{1}_{\{j_i \geq q_i\}}$ for all $j_i \in \{0, \dots, q_i + 1\}$ and
 981 for all $i \in M$ where

$$982 \quad (\alpha_i^j, \beta_i^j) = \left(\Psi(W_j^i) - \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} W_j^i, \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} a_i \right) \text{ for } j = 0, \dots, q_i \quad (79)$$

$$983 \quad (\alpha_i^{q_i+1}, \beta_i^{q_i+1}) = (\Psi(a_i), 0)$$

984 and $\bar{\mu}$, W_j^i , Δ_j^i and q_i are as defined in Lemma 3.25. For a given inequality of the form (78), let $L = \{i \in$
 985 $M \mid j_i \geq q_i, I(a_i) = 1\}$. In particular, (78) is facet-defining for PB if one of the following conditions holds:

- 986 1. $L = \emptyset$.
- 987 2. $\exists \bar{i} \in M$ such that $j_{\bar{i}} = 0$.

988 *Proof.* It follows from Proposition 3.23 that $\Psi(w)$ is a valid subadditive approximation of $\Phi(w)$ for $w \geq 0$.
 989 Hence, lifting coefficients (α_i, β_i) of (x_i, y_i) for $i \in M$ are valid if they satisfy the condition

$$990 \quad \alpha_i x_i + \beta_i y_i \geq \Psi(a_i x_i y_i) \quad \text{for } (x_i, y_i) \in \{0, 1\} \times [0, 1] \setminus \{0, 0\}. \quad (80)$$

991 Condition (80) can be restated as

$$992 \quad \beta_i \phi \geq \Psi(0) \quad \text{for } 0 < \phi \leq 1, \quad (81)$$

$$993 \quad \alpha_i + \beta_i \phi \geq \Psi(a_i \phi) \quad \text{for } 0 \leq \phi \leq 1. \quad (82)$$

994 To prove that (78) defines a face of PB of dimension at least $(2n - 1) - \sum_{i \in M} I(a_i) \times \mathbf{1}_{\{j_i \geq q_i\}}$ when lifting
 995 coefficients are chosen according to (79), we will show that, for each $i \in M$,

$$996 \quad \alpha_i x_i + \beta_i y_i = \Phi(a_i x_i y_i) \quad (83)$$

997 is satisfied at equality by at least $2 - I(a_i) \times \mathbf{1}_{\{j_i \geq q_i\}}$ independent points.

998 First, consider the case where $(\alpha_i, \beta_i) = (\Psi(a_i), 0)$. Observe that (81) is satisfied since $\beta_i = 0$ and
 999 $\Psi(0) = 0$. Further, (82) holds as $\alpha_i = \alpha_i + \beta_i \phi = \Psi(a_i) \geq \Psi(a_i \phi)$ since Ψ is a nondecreasing function. It is
 1000 easily verified that (83) is satisfied at equality at the point $(0, 1)$. Further, when $I(a_i) = 0$, then (83) is also
 1001 satisfied at equality at the point $(1, 1)$.

1002 Second, consider the case where

$$1003 \quad (\alpha_i, \beta_i) = \left(\Psi(W_j^i) - \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} W_j^i, \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} a_i \right).$$

1004 Clearly, (α_i, β_i) satisfies (81) since $\beta_i \geq 0$. From Lemma 3.25, we have that

$$1005 \quad \begin{aligned} \Phi(a_i \phi) \leq \Psi(a_i \phi) &\leq \Psi(W_j^i) + \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} (a_i \phi - W_j^i) \\ &= \left(\Psi(W_j^i) - \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} W_j^i \right) + \frac{\Psi(W_{j+1}^i) - \Psi(W_j^i)}{\Delta_j^i} a_i \phi \\ &= \alpha_i + \beta_i \phi. \end{aligned}$$

1006 We now present points that satisfy (83) at equality. Observe first that, for $j = 0, \dots, q_i$, the point $(x_i^*, y_i^*) =$
 1007 $\left(1, \frac{W_j^i}{a_i} \right)$ satisfies (83) at equality since $\Psi(a_i x_i^* y_i^*) = \Psi(W_j^i) = \Psi(B_{j-1} + \bar{\mu}) = \Phi(B_{j-1} + \bar{\mu})$. Similarly, for

1008 $j = 0, \dots, q_i - 1$, the point $(x_i^*, y_i^*) = \left(1, \frac{W_{j+1}^i}{a_i}\right)$ satisfies (83) at equality. For $j = q_i$, the point $\left(1, \frac{W_{j+1}^i}{a_i}\right)$
1009 reduces to $(1, 1)$ which satisfies (83) at equality when $\Psi(a_i) = \Phi(a_i)$, i.e., when $I(a_i) = 0$. Therefore, we
1010 conclude that (78) defines a face of PB of dimension at least $(2n - 1) - \sum_{i \in M} I(a_i) \times \mathbf{1}_{\{j_i \geq q_i\}}$.

1011 We conclude from the above derivation that when, for all $i \in M$, either $j_i < q_i$ or $I(a_i) = 0$, then the
1012 face of PB that (78) defines has dimension $2n - 1$ showing that (78) is facet-defining for PB and proving 1.
1013 Now, we show that (78) is also facet-defining if $j_{\bar{i}} = 0$ for some $\bar{i} \in M$. We first lift $(x_{\bar{i}}, y_{\bar{i}})$. Since $a_{\bar{i}} \geq \bar{\mu}$
1014 (see discussion preceding Proposition 3.22) it follows that $(\alpha_{\bar{i}}^0, \beta_{\bar{i}}^0) = (0, \frac{a_{\bar{i}}}{\bar{\mu}})$. Then, we lift the variables
1015 in $M \setminus \{L \cup \{\bar{i}\}\}$ and choose any $j_i \leq q_i + 1$ for these variables. The above proof shows that the resulting
1016 inequality is facet-defining for $PB(L \setminus \{\bar{i}\}, \emptyset, L \setminus \{\bar{i}\}, \emptyset)$. Since $PB(L \setminus \{\bar{i}\}, \emptyset, L \setminus \{\bar{i}\}, \emptyset) \subseteq PB$, all the points tight
1017 for (78) are feasible to PB . Now, we lift a variable $i' \in L \setminus \{\bar{i}\}$. Let

$$F(w, a) = \left\{ (x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{i \in K} a_i x_i y_i \geq d - a' - w \text{ and } \sum_{i \in K} x_i = |K| - 1 - a \right\}.$$

1018 We show that there exists $p \in F(B_{q_{i'}} + \mu, q_{i'} + 1)$ which is feasible to PB and tight on (78). First note that
1019 $F(B_{q_{i'}} + \mu, q_{i'} + 1) \neq \emptyset$ because $\Phi(B_{q_{i'}} + \mu) = q_{i'} + 1$. Let $p = (x', y')$. By the definition of $F(w, a)$, we are
1020 free to redefine (x'_i, y'_i) for $i \notin K$. Let $x'_i = y'_i = 0$ for $i \in M \setminus \{L \cup \{\bar{i}\}\}$ and let $x'_i = y'_i = 1$ for $i \in \widehat{M}$. Let
1021 $x'_{\bar{i}} = 1$ and $y'_{\bar{i}} = \frac{B_{q_{i'}} + \bar{\mu} - a_{i'}}{a_{\bar{i}}}$. Since $a_{\bar{i}} \geq \bar{\mu}$ and $B_{q_{i'}} < a_{i'} \leq B_{q_{i'}} + \bar{\mu}$ it follows that $0 < y_{\bar{i}} \leq 1$. Finally, we set
1022 $(x'_{i'}, y'_{i'}) = (1, 1)$. Note that $a_{\bar{i}} x'_{\bar{i}} y'_{\bar{i}} + a_{i'} x'_{i'} y'_{i'} = B_{q_{i'}} + \bar{\mu}$ and

$$\alpha_{\bar{i}}^0 x'_{\bar{i}} + \beta_{\bar{i}}^0 y'_{\bar{i}} + \alpha_{i'}^{j_{i'}} x'_{i'} + \beta_{i'}^{j_{i'}} y'_{i'} = \frac{B_{q_{i'}} + \bar{\mu} - a_{i'}}{\bar{\mu}} + q_{i'} + \frac{a_{i'} - B_{q_{i'}}}{\bar{\mu}} = q_{i'} + 1 = \Psi(B_{q_{i'}} + \mu) = \Phi(B_{q_{i'}} + \mu),$$

1023 where the first equality holds since $(\alpha_{\bar{i}}^0, \beta_{\bar{i}}^0) = (0, \frac{a_{\bar{i}}}{\bar{\mu}})$, $(\alpha_{i'}^{j_{i'}}, \beta_{i'}^{j_{i'}}) = \left(q_{i'} - \theta \frac{B_{q_{i'}-1} + \bar{\mu}}{a_{i'}}, \theta\right)$ when $j_{i'} = q_{i'}$ and
1024 $(\alpha_{i'}^{j_{i'}}, \beta_{i'}^{j_{i'}}) = (\Psi(a_{i'}), 0)$ when $j_{i'} = q_{i'} + 1$ where $\theta = \frac{(\Psi(a_{i'}) - q_{i'}) a_{i'}}{a_{i'} - B_{q_{i'}-1} - \bar{\mu}}$ and $\Psi(a_{i'}) = q_{i'} + \frac{a_{i'} - B_{q_{i'}}}{\bar{\mu}}$. Therefore,
1025 $p \in PB$ and is tight for (78). For $j_{i'} = q_{i'}$, we have already demonstrated that there exists a point of PB
1026 tight for (78) that sets $(x_{i'}, y_{i'}) = \left(1, \frac{W_{j_{i'}}^{i'}}{a_{i'}}\right)$ and for $j_{i'} = q_{i'} + 1$, there is a point of PB tight for (78) such that
1027 $(x_{i'}, y_{i'}) = (0, 1)$. For $j_{i'} = q_{i'}$, affine independence follows since $a_{i'} > W_{j_{i'}}^{i'}$ implies that $(0, 0)$, $(1, 1)$, and
1028 $\left(1, \frac{W_{j_{i'}}^{i'}}{a_{i'}}\right)$ are affinely independent. For $j_{i'} = q_{i'} + 1$, affine independence follows from the affine independence
1029 of $(0, 0)$, $(1, 1)$, and $(0, 1)$. \square

1030 Inequalities (78) can be facet-defining depending on the value of the coefficients a_i and the choice of lifting
1031 coefficients (α_i, β_i) for $i \in M$. As mentioned before, \widehat{M} may be chosen to be any subset of \overline{M} that satisfies
1032 $\sum_{i \in \widehat{M}} a_i < \mu'$. In this case, (78) will be facet-defining if $\max\{a_i \mid i \in M, j_i = 0\} \geq \bar{\mu}$ but it may not be
1033 facet-defining otherwise. The next example illustrates the use of (78) in deriving facets of PB .

1034 **Example 3.27.** Consider the clique inequality (76) of Example 3.24 and its corresponding approximate lifting
1035 function. We have $q_1 = 2$ and $q_2 = 1$ with $W_0^1 = 0$, $W_1^1 = 3$, $W_2^1 = 20$, $W_3^1 = 21$, and $W_0^2 = 0$, $W_1^2 = 3$,
1036 $W_2^2 = 19$. Applying Theorem 3.26, we obtain the following nine inequalities

$$\left\{ \begin{array}{l} \frac{14}{17}x_1 + \frac{21}{17}y_1 \\ 2x_1 \end{array} \right\} + \left\{ \begin{array}{l} \frac{21}{24}x_2 + \frac{19}{24}y_2 \\ \frac{5}{3}x_2 \end{array} \right\} + x_3 + x_4 + x_5 \geq 2,$$

1040 which define faces of PB of dimension at least 8 since $I(a_1) = 0$ and $I(a_2) = 1$. It follows from the first
1041 condition of Theorem 3.26 that following three inequalities

$$\left\{ \begin{array}{l} \frac{14}{17}x_1 + \frac{21}{17}y_1 \\ 2x_1 \end{array} \right\} + \frac{19}{3}y_2 + x_3 + x_4 + x_5 \geq 2$$

1042 are facet-defining for PB since $j_2 < q_2$. The two inequalities

$$\frac{21}{3}y_1 + \left\{ \begin{array}{l} \frac{21}{24}x_2 + \frac{19}{24}y_2 \\ \frac{5}{3}x_2 \end{array} \right\} + x_3 + x_4 + x_5 \geq 2$$

1045 are also facet-defining for PB since they satisfy the second condition for facet-defining inequalities in Theo-
 1046 rem 3.26 as $j_1 = 0$.

1047 4 Relations to fixed-charge single-node flow model without inflows

1048 In Section 3, we derived strong valid inequalities for the bilinear set B using lifting. In this section, we show
 1049 that many of these lifted inequalities are also facet-defining for the convex hull of the fixed-charge single-node
 1050 flow model without inflows

$$1051 \quad F = \left\{ (x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{j=1}^n a_j y_j \geq d, x_j \geq y_j \forall j \in N \right\}.$$

1052 In the following lemma, we show that $F \subseteq B$.

1053 **Lemma 4.1.** *The bilinear covering set B is a relaxation of the flow set F .*

1054 *Proof.* We prove that $F \subseteq B$. Let $(x, y) \in \{0, 1\}^n \times [0, 1]^n$ be an arbitrary point of F . It suffices to show
 1055 that $\sum_{j=1}^n a_j x_j y_j \geq d$. Let $N_0 = \{j \in N \mid x_j = 0\}$ and $N_1 = \{j \in N \mid x_j = 1\}$. Since $(x, y) \in F$, $y_j = 0$ for
 1056 all $j \in N_0$. Then,

$$1057 \quad \sum_{j \in N} a_j x_j y_j = \sum_{j \in N_1} a_j y_j = \sum_{j \in N} a_j y_j \geq d,$$

1058 where the last inequality holds because $(x, y) \in F$. □

1059 Fixed-charge single-node flow sets are important in practice since they can be used as a source of cutting
 1060 planes for 0–1 mixed-integer programs. Further, they naturally arise in the formulation of fixed-charge
 1061 network problems; see [2, 12, 19, 20, 22]. The fixed-charge single-node flow set F without inflows was first
 1062 studied by Padberg et al. [22] under the assumptions that (i) $a_i \leq d$ and (ii) $\sum_{j=1}^n a_j > d + a_i$ for all $i \in N$.
 1063 In the following, we relate the facets of PF to those of PB without assuming that the sets are full-dimensional.

1064 **Lemma 4.2** (Adapted from Proposition 8 in Padberg et al. [22]). *Every facet-defining inequality of PF that*
 1065 *is not a multiple of $y_i \leq x_i$ can be expressed as $\alpha x + \beta y \geq \delta$, where $\beta \geq 0$.*

1066 *Proof.* If for some i , $\beta_i < 0$ then the only points tight on this inequality are such that $y_i = x_i$. If F satisfies
 1067 this equality then we may rewrite the facet-defining inequality as $\alpha x + \beta_i x_i + \beta y - \beta_i y_i \geq \delta$. □

1068 In the following, we refer to the facet-defining inequalities of PF that are not a multiple of $y_i \leq x_i$ as
 1069 *non-trivial* facet-defining inequalities.

1070 **Lemma 4.3.** $\text{aff}(F) = \text{aff}(B)$.

1071 *Proof.* Clearly, $\text{aff}(F) \subseteq \text{aff}(B)$ since $F \subseteq B$ by Lemma 4.1. It therefore remains to prove that $\text{aff}(B) \subseteq$
 1072 $\text{aff}(F)$. Consider any point $(x, y) \in B$. If $(x, y) \in F$, then clearly $(x, y) \in \text{aff}(F)$. We may therefore
 1073 assume that $(x, y) \in B \setminus F$. Define $p = (x', y')$ where $(x'_i, y'_i) = (x_i, x_i y_i)$ for $i \in N$. It is easy to see that
 1074 $\sum_{i \in N} a_i y'_i = \sum_{i \in N} a_i x_i y_i \geq d$ and $y'_i \leq x'_i$ for $i \in N$ and so $p \in F$. Let $I' = \{i \in N \mid y_i > x_i\}$. We show next
 1075 that for each $i \in I'$, $p^i = p + (0, e_i) \in \text{aff}(F)$. To this end, observe that $x'_i = 0$ for each $i \in I'$. It follows easily
 1076 that $q^i = p + (e_i, 0)$ and $r^i = p + (e_i, e_i)$ belong to F . Therefore, $p^i = p + (r^i - q^i) \in \text{aff}(F)$. Now, observe
 1077 that $(x, y) = p + \sum_{i \in I'} y_i (p^i - p) \in \text{aff}(F)$. It follows that $B \subseteq \text{aff}(F)$ and therefore $\text{aff}(B) \subseteq \text{aff}(F)$. □

1078 **Proposition 4.4.** *Assume that*

$$1079 \quad \alpha x + \beta y \geq \delta \tag{84}$$

1080 *is valid for PF and, for each $i \in N$, either $\alpha_i \leq 0$ or $\beta_i \geq 0$. Then, (84) is valid for PB . Further, for every*
 1081 *non-trivial facet (84) of PF with $\beta \geq 0$, (84) is facet-defining inequality for PB .*

1082 *Proof.* We first show that (84) is valid for B . Consider $(x, y) \in B$. Let $I = \{i \in N \mid \alpha_i \leq 0\}$. Define (x', y')
 1083 such that $(x'_i, y'_i) = (1, y_i)$ for $i \in I$ and $(x'_i, y'_i) = (x_i, x_i y_i)$ for $i \in N \setminus I$. Then,

$$1084 \quad \sum_{i \in N} a_i y'_i = \sum_{i \in I} a_i y_i + \sum_{i \in N \setminus I} a_i x_i y_i \geq \sum_{i \in N} a_i x_i y_i \geq d,$$

1085 where the last inequality holds because $(x, y) \in B$. Further, since $y'_i \leq x'_i$, it follows that $(x', y') \in F$. Then,

$$1086 \quad \delta \leq \alpha x' + \beta y' \leq \alpha x + \beta y,$$

1087 where the first inequality holds because $(x', y') \in F$ and the second inequality is satisfied since, by construc-
 1088 tion, $\alpha(x' - x) + \beta(y' - y) \leq 0$. It follows that (84) is valid for PB .

1089 Consider a non-trivial facet-defining inequality $\alpha'x + \beta'y \geq \delta'$ of PF with $\beta' \geq 0$. Then, by the first part
 1090 of this result, it follows that $\alpha'x + \beta'y \geq \delta'$ is valid for PB . Since, by Lemmas 4.1 and 4.3 respectively, $B \supseteq F$
 1091 and $\dim(B) = \dim(F)$, it follows that $\alpha'x + \beta'y \geq \delta'$ defines a facet of B . \square

1092 In Proposition 4.4, the assumption that $\beta \geq 0$ for a facet-defining inequality is without loss of generality
 1093 because of Lemma 4.2. As an immediate consequence of Proposition 4.4, it can be shown that lifting functions
 1094 associated with inequalities $\alpha x + \beta y \geq \delta$, such that for each i either $\alpha_i \leq 0$ or $\beta_i \geq 0$ are identical when they are
 1095 computed over B or over F . Since the inequalities derived in Section 3 as well as the seed inequalities satisfy
 1096 these assumptions, many of our results in Section 3 extend to the study of F . We record this observation in
 1097 the following corollary.

1098 **Corollary 4.5.** *Let $(\alpha, \beta) \in \mathbb{R}^{2n}$ and, for each $i \in N$, assume that either $\alpha_i \leq 0$ or $\beta_i \geq 0$. Let $B(w) =$
 1099 $\{(x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{i \in N} a_i x_i y_i \geq d - w\}$, and $F(w) = \{(x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{i \in N} a_i y_i \geq$
 1100 $d - w \text{ and } y_i \leq x_i \text{ for all } i \in N\}$, where $a_i \geq 0$ for all $i \in N$. Let $z_B(w) = \min\{\alpha x + \beta y \mid (x, y) \in B(w)\}$
 1101 and $z_F(w) = \min\{\alpha x + \beta y \mid (x, y) \in F(w)\}$. Then, $z_B(w) = z_F(w)$.*

1102 *Proof.* By Lemma 4.1, $B(w) \supseteq F(w)$. It follows that $z_B(w) \leq z_F(w)$. We now argue that $z_B(w) \geq z_F(w)$. By
 1103 the definition of $z_F(w)$, $\alpha x + \beta y \geq z_F(w)$ is valid for $F(w)$, which is a flow-set. Let $(x', y') \in \operatorname{argmin}\{\alpha x + \beta y \mid$
 1104 $(x, y) \in B(w)\}$. Then, $z_B(w) = \alpha x' + \beta y' \geq z_F(w)$, where the inequality follows from Proposition 4.4. We
 1105 conclude that $z_B(w) = z_F(w)$. \square

1106 Now, we illustrate Proposition 4.4 via an example.

1107 **Example 4.6.** *Consider the fixed-charge single-node flow set without inflows*

$$1108 \quad F = \left\{ (x, y) \in \{0, 1\}^4 \times [0, 1]^4 \mid 19y_1 + 17y_2 + 15y_3 + 10y_4 \geq 20, x_j \geq y_j, \forall j = 1, \dots, 4 \right\},$$

1109 *corresponding to the bilinear covering set B discussed in Example 2.6. We obtained a complete linear de-*
 1110 *scription of PF using PORTA; see Christof and Löbel [6]. This linear description is given in the Appendix.*
 1111 *We observe that inequalities (10), (11), (17), and (18) are facets for both PB and PF . However, it can be*
 1112 *verified that inequalities (12), (13), (15), and (16) are facet-defining for PB but not for PF . We mention*
 1113 *that the inequalities of PF described in the Appendix have been numbered according to their counterparts in*
 1114 *PB .*

1115 Proposition 4.4 is surprising in light of Lemma 4.1 because on the one hand $F \subsetneq B$ and on the other hand
 1116 the nontrivial facets of PF are facets of PB . In other words, a polyhedral description of PF can be derived
 1117 from that of PB by adding the trivial facets of PF . The converse, however, is not true. As an illustration,
 1118 inequality (20) in the Appendix is a non-trivial facet-defining inequality of PB that is not facet-defining for
 1119 PF . Surprisingly, a partial converse to Proposition 4.4 does hold.

1120 We will show that an inequality description of PB can be obtained given the facet-defining inequalities
 1121 for PF . The key to this construction is the result of Lemma 4.7 which shows that F and B can be viewed as
 1122 projections of the same set onto different subspaces. Let

$$1123 \quad S = \left\{ (x, y, z) \in \{0, 1\}^n \times [0, 1]^n \times \mathbb{R}^n \mid \sum_{j=1}^n a_j z_j \geq d, z_j = x_j y_j, \forall j \in N \right\}.$$

1124 **Lemma 4.7.** *The projection of S onto the (x, z) space is F while the projection of S onto (x, y) space is B .*
 1125 *Consequently, $\text{proj}_{(x,z)} \text{conv}(S) = PF$ and $\text{proj}_{(x,y)} \text{conv}(S) = PB$.*

1126 *Proof.* First, we show that $\text{proj}_{(x,z)} S = F$. If $(x, y, z) \in S$, it is clear that $(x, z) \in F$ since $0 \leq z_j \leq x_j$ and
 1127 $\sum_{j=1}^n a_j z_j \geq d$. If $(x, z) \in F$, then $0 \leq z_j \leq x_j$ and $x_j \in \{0, 1\}$ imply that $z_j = x_j z_j$. Therefore, $(x, z, z) \in S$.
 1128 Second, we show that $\text{proj}_{(x,y)} S = B$. This follows by substituting $x_j y_j$ for z_j in $\sum_{j=1}^n a_j z_j \geq d$. The last
 1129 statement follows since $\text{conv}(AS) = A \text{conv}(S)$ for any linear transformation A . \square

1130 Surprisingly, $\text{conv}(S)$ can be described using facet-defining inequalities for PF . We write that $(\alpha, \beta, \gamma) \in$
 1131 $\mathcal{F}(PF)$ if $\alpha x + \beta y \geq \gamma$ is a facet-defining inequality of PF that is not a multiple of $y_j \leq x_j$. Define

$$1132 \quad G = \{(x, y, z) \mid \alpha x + \beta z \geq \gamma \forall (\alpha, \beta, \gamma) \in \mathcal{F}(PF), z \leq \min\{x, y\}, y \leq z + \mathbf{1} - x\}.$$

1133 **Theorem 4.8.** $G = \text{conv}(S)$.

1134 *Proof.* (\supseteq) To show that $\text{conv}(S) \subseteq G$, it suffices to show that $S \subseteq G$ because G is convex. Consider
 1135 $(x, y, z) \in S$. Then, by Lemma 4.7, $(x, z) \in F$ and, therefore, $\alpha x + \beta z \geq \gamma$ for all (α, β, γ) in $\mathcal{F}(PF)$. Further,
 1136 by McCormick inequalities, $x_j + y_j - 1 \leq x_j y_j \leq \min\{x_j, y_j\}$. Therefore, (x, y, z) satisfies the defining
 1137 inequalities of G .

1138 (\subseteq) Now, we show that $G \subseteq \text{conv}(S)$. If $(x, y, z) \in G$, then $(x, z) \in PF$ and $z \leq y \leq z + \mathbf{1} - x$. Therefore,
 1139 there exists I such that $(x, z) = \sum_{i \in I} \lambda_i (x^i, z^i)$ where $(x^i, z^i) \in F$, $\lambda_i \geq 0$ for $i \in I$, and $\sum_{i \in I} \lambda_i = 1$. We
 1140 define $f_j = \frac{y_j - z_j}{1 - x_j}$ if $x_j < 1$ and 0 otherwise. Let $I_j^1 = \{i \in I \mid x_j^i = 1\}$. Now, consider (x^i, y^i, z^i) where
 1141 $y_j^i = z_j^i$ if $i \in I_j^1$ and $y_j^i = f_j$ if $i \in I \setminus I_j^1$. Then, $z_j^i \leq x_j^i$ and $x_j^i \in \{0, 1\}$ imply that $z_j^i = x_j^i y_j^i$. Further,

$$1142 \quad \sum_{i \in I} \lambda_i y_j^i = \sum_{i \in I_j^1} \lambda_i z_j^i + \sum_{i \in I \setminus I_j^1} \lambda_i f_j = z_j + (1 - x_j) f_j = y_j,$$

1143 where the second equality follows since $z_j = \sum_{i \in I} \lambda_i z_j^i = \sum_{i \in I_j^1} \lambda_i z_j^i$, $\sum_{i \in I_j^1} \lambda_i = x_j$, and $\sum_{i \in I} \lambda_i = 1$, and
 1144 the last equality since $x_j = 1$ implies that $z_j = y_j$. Therefore, $(x, y, z) = \sum_{i \in I} \lambda_i (x^i, y^i, z^i) \in \text{conv}(S)$. \square

1145 Finally, we show that the projections of G to (x, z) space and (x, y) space are not altered even if G is
 1146 relaxed in a certain way. Let

$$1147 \quad R = \{(x, y, z) \mid \alpha x + \beta z \geq \gamma \forall (\alpha, \beta, \gamma) \in \mathcal{F}(PF), z \leq \min\{x, y\}, y \leq \mathbf{1}\}.$$

1148 **Corollary 4.9.** $PF = \text{proj}_{(x,z)} R$ and $PB = \text{proj}_{(x,y)} R$.

1149 *Proof.* We will show that $\text{proj}_{(x,z)} R = \text{proj}_{(x,z)} G$ and $\text{proj}_{(x,y)} R = \text{proj}_{(x,y)} G$. Then, the result follows
 1150 from Lemma 4.7 and Theorem 4.8. Since $z + \mathbf{1} - x \leq \mathbf{1}$, it follows that $R \supseteq G$. First, we show that
 1151 $\text{proj}_{(x,z)} R \subseteq \text{proj}_{(x,z)} G$. Assume that $(x, y, z) \in R$. Then, define $y' = z + \mathbf{1} - x$. Since $z + \mathbf{1} - x \geq z$ it
 1152 follows that $(x, y', z) \in G$. Second, we show that $\text{proj}_{(x,y)} R \subseteq \text{proj}_{(x,y)} G$. Assume that $(x, y, z) \in R$. Then,
 1153 let $z' = \max\{z, x + y - \mathbf{1}\}$. By Lemma 4.2, for all $(\alpha, \beta, \gamma) \in \mathcal{F}(PF)$, $\beta \geq 0$. Therefore, $\alpha x + \beta z' \geq \alpha x + \beta z \geq \gamma$.
 1154 Further, $z' = \max\{z, x + y - \mathbf{1}\} \leq \min\{x, y\}$ since $z \leq \min\{x, y\}$ and $x, y \in [0, 1]^2$. Finally, by construction,
 1155 $y \leq z' + \mathbf{1} - x$. Therefore, $(x, y, z') \in G$. \square

1156 Corollary 4.9 implies every non-trivial facet of PB arises as a conic combination of a single non-trivial
 1157 facet of PF and (possibly multiple) trivial facet-defining inequalities $y_j \leq x_j$.

1158 **Corollary 4.10.** *Let $\alpha x + \beta y \geq \gamma$ be a facet-defining inequality for PB where $\beta \geq 0$. Then, $\alpha x + \beta y \geq \gamma$
 1159 defines a non-empty face of F . Further, there exists (α', β') and $\lambda \geq 0$ such that $(\alpha, \beta) = (\alpha' + \lambda, \beta' - \lambda)$,
 1160 where $\alpha' x + \beta' y \geq \gamma$ is facet-defining for PF and for $j = 1, \dots, n$, $\lambda_j \beta_j = 0$.*

1161 *Proof.* Let $\delta = \min\{\alpha x + \beta y \mid (x, y) \in PF\}$. Since, by Lemma 4.1, $F \subseteq B$, it follows that $\delta \geq \gamma$. By
 1162 Proposition 4.4, $\alpha x + \beta y \geq \delta$ is valid for PB . Therefore, $\delta \leq \gamma$. In other words, $\delta = \gamma$ and $\alpha x + \beta y \geq \gamma$
 1163 defines a non-empty face of F . By Corollary 4.9 and Fourier-Motzkin elimination of z from R it follows that,

$$1164 \quad PB = \{(x, y) \mid \alpha' x + \beta'^J x + \beta'^{N \setminus J} y \geq \gamma' \forall (\alpha', \beta', \gamma') \in \mathcal{F}(PF) \text{ and } J \subseteq N, y \leq \mathbf{1}\},$$

1165 where $\beta'^J = \beta'_j$ if $j \in J$ and $\beta'^J = 0$ otherwise. Since (α, β, γ) is not a multiple of $y_j \leq 1$, it follows that
 1166 there exists $J \subseteq N$ and $(\alpha', \beta', \gamma') \in \mathcal{F}(PF)$ such that $(\alpha, \beta) = (\alpha' + \beta'^J, \beta' - \beta'^J)$. \square

1167 **Example 4.11.** Consider the inequality $126x_1+90x_3+45x_4+153y_2 \geq 135$ that is facet-defining for the bilinear
1168 covering set of Example 2.6 but not facet-defining for the corresponding flow set presented in Example 4.6; see
1169 Appendix for a complete description of facet-defining inequalities of PB where this inequality is numbered (20).
1170 Then, as Corollary 4.10 proves, this inequality can be expressed as a sum of $50x_1+90x_3+45x_4+76y_1+153y_2 \geq$
1171 135 and $76x_1 - 76y_1 \geq 0$, which are the facet-defining inequalities of the corresponding flow-set numbered (1)
1172 and (f1) in the Appendix.

1173 Proposition 4.4 and Corollary 4.9 show that a polyhedral description of either PF or PB can be derived
1174 explicitly given the facet-defining inequalities of the other. In fact, Proposition 4.4 also shows that an affine
1175 function over either B or F can be optimized if we have an oracle for optimizing an affine function over the
1176 other set. We discuss the reduction below. Let $l(x, y) = \alpha x + \beta y - \gamma$ and define $I = \{i \in N \mid \alpha_i > 0, \beta_i < 0\}$.
1177 Let $z_B(l) = \min\{l(x, y) \mid (x, y) \in B\}$ and $z_F(l) = \min\{l(x, y) \mid (x, y) \in F\}$. Define $l'(x, y) = \alpha x +$
1178 $\sum_{i \in N \setminus I} \beta_i y_i + \sum_{i \in I} \beta_i - \gamma$. While minimizing $l(x, y)$ over B , y_i can be set to 1 whenever $\beta_i \leq 0$. Therefore,
1179 it follows that $z_B(l) = z_B(l')$. However, by Proposition 4.4, $z_F(l') = z_B(l')$. Therefore, $z_B(l) = z_F(l')$. If
1180 (x, y) is the optimal solution to $z_F(l')$, the optimal solution to $z_B(l)$ is (x, y') where $y'_i = 1$ if $i \in I$ and
1181 $y'_i = y_i$ otherwise. Now, we consider the converse. Define $l''(x, y) = \alpha x + \sum_{i \in I} \beta_i x_i + \sum_{i \in N \setminus I} \beta_i y_i - \gamma$. While
1182 minimizing $l(x, y)$ over F , y_i can be set to x_i whenever $\beta_i \leq 0$. Therefore, $z_F(l) = z_F(l'')$. But, then by
1183 Proposition 4.4, $z_F(l'') = z_B(l'')$. Therefore, $z_F(l) = z_B(l'')$. The optimal solution can be obtained as in the
1184 proof of Proposition 4.4.

1185 Given the relationships between the polyhedra PB and PF , it is reasonable to expect that the inequalities
1186 we developed in Section 3 reveal facets of PF . We now provide a detailed discussion of which inequalities are
1187 facet-defining for PF . For the remainder of this section, we assume, as we did for PB , that

1188 **Assumption 3.** $\sum_{j=1}^n a_j \geq d + a_i$ for all $i \in N$.

1189 Under Assumption 3, it follows from Lemma 4.3 that PF is a full-dimensional polytope.

1190 **Theorem 4.12.** A lifted bilinear cover inequality (54) is facet-defining for PF if and only if

$$1191 (\alpha_i, \beta_i) \in \bigcup_{j=0}^{q_i} \left(P^C(Q_j^i) - \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} Q_j^i, \frac{P^C(Q_{j+1}^i) - P^C(Q_j^i)}{\Delta_j^i} a_i \right) \quad (85)$$

1192 for all $i \in M$.

1193 *Proof.* The proof of Proposition 2.9 already shows that (42) is facet-defining for $PF(M, C', M, C')$ since all
1194 the points considered are feasible to the flow set.

1195 Now, it suffices to show that sufficiently many of the tight points added when lifting variables (x_i, y_i) for
1196 $i \in M \cup C'$ also belong to PF . When we lifted variables (x_i, y_i) for $i \in C'$ in the proof of Proposition 3.9, we
1197 added the two affinely independent points $(0, 0)$ and $\left(1, \frac{(a_i - \mu)^+}{a_i}\right)$ that both correspond to feasible solutions
1198 of F ; see (42) and Corollary 4.5. When lifting the variables (x_i, y_i) for $i \in M$ in Theorem 3.14, we added the
1199 two points $\left(1, \frac{Q_j^i}{a_i}\right)$ and $\left(1, \frac{Q_{j+1}^i}{a_i}\right)$ that both correspond to feasible solutions of F ; see (45) and Corollary 4.5.

1200 Next, we show that if (54) is facet-defining for PF , then (α_i, β_i) must be chosen as in (85). It suffices to
1201 show that if $(\alpha_i, \beta_i) = (P^C(a_i), 0)$ for some $i \in M$ and if at least one of the coefficients pair $(P^C(a_i), 0)$ does
1202 not reduce to coefficients studied before (which happens when $P^C(a_v) \neq P^C(Q_{q_v}^v)$ for some v), then (54) is
1203 not facet-defining for PF . We will do so by showing that in such a case, (54) can be obtained by combining
1204 a different (facet-defining) inequality of the form (54) for PF with trivial facets $y_i \leq x_i$ of PF . Let $V \subseteq M$
1205 be the set of lifting coefficients (α_v, β_v) chosen to be $(P^C(a_v), 0)$. Inequality (54) then reduces to

$$1206 \sum_{v \in V} P^C(a_v) x_v + \sum_{i \in C} (a_i - \mu)^+ x_i + \sum_{j \in T} a_j y_j + \sum_{i \in M \setminus V} \alpha_i x_i + \sum_{i \in M \setminus V} \beta_i y_i \geq \sum_{i \in C} (a_i - \mu)^+. \quad (86)$$

1207 Using the first part of this proof, we know that choosing lifting coefficients

$$1208 \left(\left(P^C(Q_{q_v}^v) - \frac{P^C(Q_{q_v+1}^v) - P^C(Q_{q_v}^v)}{\Delta_{q_v}^v} Q_{q_v}^v \right), \left(\frac{P^C(Q_{q_v+1}^v) - P^C(Q_{q_v}^v)}{\Delta_{q_v}^v} a_v \right) \right)$$

1209 for $v \in V$ yields the following facet-defining inequality

$$1210 \quad \sum_{v \in V} \left(P^C(Q_{q_v}^v) - \frac{P^C(Q_{q_{v+1}}^v) - P^C(Q_{q_v}^v)}{\Delta_{q_v}^v} Q_{q_v}^v \right) x_v + \left(\frac{P^C(Q_{q_{v+1}}^v) - P^C(Q_{q_v}^v)}{\Delta_{q_v}^v} a_v \right) y_v \quad (87)$$

$$1211 \quad + \sum_{i \in C} (a_i - \mu)^+ x_i + \sum_{j \in T} a_j y_j + \sum_{i \in M \setminus V} \alpha_i x_i + \sum_{i \in M \setminus V} \beta_i y_i \geq \sum_{i \in C} (a_i - \mu)^+$$

1212 for PF . Summing (87) with

$$1213 \quad \left(\frac{P^C(Q_{q_{v+1}}^v) - P^C(Q_{q_v}^v)}{\Delta_{q_v}^v} a_v \right) (x_v - y_v) \geq 0, \forall v \in V \quad (88)$$

1214 we obtain (86) since $Q_{q_{v+1}}^v = a_v$ and $\Delta_{q_v}^v = a_v - Q_{q_v}^v$. Since we assumed that $P^C(a_v) - P^C(Q_{q_v}^v) > 0$ for
 1215 some $v \in V$ and because it is easy to see that (88) does not define the same face of PF that (87) defines, we
 1216 conclude that (86) is not facet-defining for PF . \square

1217 We remark that in the proof of Theorem 4.12, we proved that a few inequalities of the type (54) are facet-
 1218 defining for PB but not for PF . This was shown by expressing these inequalities using another non-trivial
 1219 facet of PF and the inequalities $y_j \leq x_j$. We have already shown in Corollary 4.10 that this construction
 1220 can be used to describe all facet-defining inequalities of PB that are not facet-defining for PF . We will use
 1221 similar constructions later in the section. As a consequence of Theorem 4.12, we obtain the following result
 1222 initially obtained by Padberg et al. [22].

1223 **Corollary 4.13.** (Adapted from Proposition 12 in Padberg et al. [22]) Assume that (i) \mathcal{C} is a cover with
 1224 excess $\bar{\mu} = \sum_{j \in \mathcal{C}} a_j - d$ such that $\bar{a} = \max_{j \in \mathcal{C}} a_j > \bar{\mu}$ and (ii) $\mathcal{L} \subseteq N \setminus \mathcal{C}$ is chosen so that $0 < \bar{a} - \bar{\mu} < a_k \leq \bar{a}$
 1225 for all $k \in \mathcal{L}$ and $\sum_{j \in N \setminus \mathcal{L}} a_j > d + \bar{a}$. Then

$$1226 \quad \sum_{j \in \mathcal{C}} (a_j - \bar{\mu})^+ x_j + \sum_{j \in \mathcal{L}} (\bar{a} - \bar{\mu}) x_j + \sum_{j \in N \setminus (\mathcal{C} \cup \mathcal{L})} a_j y_j \geq \sum_{j \in \mathcal{C}} (a_j - \bar{\mu})^+ \quad (89)$$

1227 is facet-defining for PF .

1228 *Proof.* Let \mathcal{C} and $\mathcal{L} \subseteq N \setminus \mathcal{C}$ be given that satisfy conditions (i) and (ii) of Corollary 4.13. Select $l \in$
 1229 $\operatorname{argmax}\{a_j \mid j \in \mathcal{C}\}$. Define $C' = \mathcal{C} \setminus \{l\}$, $M = \mathcal{L}$, and $T = N \setminus (\mathcal{C} \cup \mathcal{L})$. Clearly, $\mu = \bar{\mu}$. Observe further that
 1230 $a_l = \bar{a} > \mu$ and that $\sum_{j \in T} a_j > a_l - \bar{\mu}$ since $\sum_{j \in N \setminus \mathcal{L}} a_j > d + \bar{a}$. It follows that $(C', \{l\}, M, T)$ is a partition
 1231 of N that satisfies Conditions (A1), (A2), (A3), and (A4) of Theorem 3.14. We obtain from Assumption (ii)
 1232 that $A_1 - \mu < a_i \leq A_1 < A_2 - \mu$ for $i \in M$, which implies that $q_i = 1$ for all $i \in M$ in Lemma 3.13. Further,
 1233 since $Q_1^i = A_1 - \mu$ and $Q_2^i = a_i$ for $i \in M$, we can select (α_i, β_i) as $(A_1 - \mu, 0)$ in (54), yielding

$$1234 \quad \sum_{j \in \mathcal{C}} (a_j - \mu)^+ x_j + \sum_{j \in T} a_j y_j + \sum_{j \in M} (A_1 - \mu) x_j \geq \sum_{j \in \mathcal{C}} (a_j - \mu)^+,$$

1235 which is exactly (89) after performing the substitutions $C = \mathcal{C}$, $T = N \setminus (\mathcal{C} \cup \mathcal{L})$, $M = \mathcal{L}$, $A_1 = \bar{a}$ and
 1236 $\mu = \bar{\mu}$. \square

1237 Observe that in (89), for each $j \in N$, either the coefficient of x_j or that of y_j is zero, whereas this is not
 1238 the case for (54). Therefore, the facet-defining inequalities obtained via (89) are strictly contained in the
 1239 facet-defining inequalities obtained via (54). In Padberg et al. [22], the authors did not explicitly impose the
 1240 condition $\sum_{j \in N \setminus \mathcal{L}} a_j > d + \bar{a}$. However, in its absence, the inequalities are not necessarily facet-defining as
 1241 we show in Example 4.14. The authors' proof implicitly made use of this assumption during an induction
 1242 step. The next example illustrates that without this assumption (89) may not define a facet of the flow set.

Example 4.14. Consider the flow set defined by

$$F = \left\{ (x, y) \in \{0, 1\}^4 \times [0, 1]^4 \mid 7y_1 + 6y_2 + 5y_3 + 4y_4 \geq 10, x_j \geq y_j \forall j = 1, \dots, 4 \right\}.$$

1243 Define $\mathcal{C} = \{1, 3\}$ and $\mathcal{L} = \{2\}$ where $\bar{a} = 7$ and $\bar{\mu} = 2$. Clearly $\bar{a} - \bar{\mu} < a_2 \leq \bar{a}$. However, the assumption
 1244 that $\sum_{j \in N \setminus \mathcal{L}} a_j > d + \bar{a}$ does not hold. Inequality (89) takes the form

$$1245 \quad 5x_1 + 5x_2 + 3x_3 + 4y_4 \geq 8. \quad (90)$$

1246 Observe that whenever (90) is satisfied at equality by a point of F , the inequality $x_1 + x_2 \geq 1$ is also tight.
 1247 Since $x_1 + x_2 \geq 1$ is clearly valid for F , it follows that (90) is not facet-defining for PF .

1248 We next show that the family of lifted bilinear cover inequalities that are proven to be facet-defining for
 1249 PF in Theorem 4.12 is larger than the family given by (89).

1250 **Example 4.15.** As established in Example 4.6, (10) and (11) are facet-defining lifted bilinear cover in-
 1251 equalities (54) for both PB and PF . They are obtained by choosing $(C', l, M, T) = (\{4\}, \{3\}, \{1\}, \{2\})$ and
 1252 $(C', l, M, T) = (\{4\}, \{2\}, \{1\}, \{3\})$ respectively in Theorem 3.14. However, as mentioned above, (10) and (11)
 1253 cannot be obtained using Corollary 4.13.

1254 Next, we show that the lifted reverse bilinear cover inequalities (66) that were shown to define facets of
 1255 PB in Section 3 also define facets of PF .

1256 **Theorem 4.16.** Lifted bilinear reverse cover inequalities (66) are facet-defining for PF if and only if $a_i >$
 1257 $a_l - \mu$ for all $i \in M$.

1258 *Proof.* Assume first that $a_i > a_l - \mu$ for all $i \in M$. It is clear that (66) is valid for F since $F \subseteq B$. Recall that
 1259 (66) is obtained in Section 3 by lifting the seed inequality (42) which is facet-defining for $PB(M, C', M, C')$.
 1260 We have shown in the proof of Theorem 4.12 that (42) is facet-defining for $PF(M, C', M, C')$. Now, we show
 1261 that the tight points added when lifting variables (x_i, y_i) for $i \in M \cup C'$ also belong to PF . When we lifted
 1262 the variables (x_i, y_i) for $i \in M$ in the proof of Proposition 3.16, we added the two linearly independent
 1263 points $(0, 1)$ and $(1, 1)$. The first of these points is not feasible for F and cannot be used for the present
 1264 derivation. However, when $a_i > a_l - \mu$, the solution $(1, 1 - \epsilon)$ for ϵ sufficiently small is feasible for F and
 1265 satisfies (59) at equality. Therefore, $(1, 1)$ and $(1, 1 - \epsilon)$ provide the desired two tight independent feasible
 1266 solutions of F ; see (42) and Corollary 4.5. When lifting the variables (x_i, y_i) for $i \in C'$ in Theorem 3.20, we
 1267 added the two points $(0, 0)$ and $\left(1, \frac{(a_i - A_1 + a_l - \mu)^+}{a_i}\right)$ which are affinely independent of $(1, 1)$ and correspond
 1268 to feasible solutions of F ; see (58) and Corollary 4.5. This proves that (66) is facet-defining for PF .

1269 Assume now that $a_i \leq a_l - \mu$ for some $i \in M$. Define $M_2 = \{i \in M \mid a_i \leq a_l - \mu\} \neq \emptyset$ and $M_1 = M \setminus M_2$.
 1270 Inequality (66) can be written as

$$1271 \quad (a_l - \mu)x_l - \sum_{j \in C'} P^M(-a_j)x_j + \sum_{j \in M_1} (a_l - \mu)x_j + \sum_{j \in M_2} a_j x_j + \sum_{j \in T} a_j y_j \geq (a_l - \mu) - \sum_{j \in C'} P^M(-a_j). \quad (91)$$

1272 Observe next that partition $(C', \{l\}, M_1, T \cup M_2)$ satisfies Conditions (A1), (A2) and (A4) since $(C', \{l\}, M, T)$
 1273 does. Further, since $a_i > a_l - \mu$ for $i \in M_1$, it follows from the first part of this proof that the lifted reverse
 1274 bilinear cover inequality

$$1275 \quad (a_l - \mu)x_l - \sum_{j \in C'} P^{M_1}(-a_j)x_j + \sum_{j \in M_1} (a_l - \mu)x_j + \sum_{j \in T \cup M_2} a_j y_j \geq (a_l - \mu) - \sum_{j \in C'} P^{M_1}(-a_j) \quad (92)$$

1276 is facet-defining for PF , where it is easy to verify that $P^{M_1}(w) = P^M(w)$ for $w \leq 0$. Now, observe that (91)
 1277 can be obtained by summing (92) and inequalities

$$1278 \quad a_j(x_j - y_j) \geq 0 \quad (93)$$

1279 for $j \in M_2$. Since (92) and (93) define different facets of the full-dimensional polyhedron PF , we conclude
 1280 that (91) is not facet-defining for PF . \square

1281 The inequalities of Theorem 4.16 are known to be valid for PF , as first shown in Gu et al. [12].

1282 **Corollary 4.17** (Adapted from Theorem 12 in Gu et al. [12]). Assume that (i) $\mathcal{C} \subseteq N$ is a generalized cover
1283 for F such that $\sum_{j \in \mathcal{C}} a_j = d - \lambda$ with $\lambda > 0$ and (ii) $\mathcal{L} \neq \emptyset$ and $\sum_{j \in N \setminus \mathcal{L}} a_j > d$ where $\mathcal{L} = \{j \in N \setminus \mathcal{C} \mid a_j > \lambda\}$.
1284 Assume also that $\mathcal{L} = \{j_1, j_2, \dots, j_r\}$ with $a_{j_i} \geq a_{j_{i+1}}$ for $i = 1, \dots, r - 1$. Let $r = |\mathcal{L}|$, $A_0 = 0$, and
1285 $A_i = \sum_{k=1}^i a_{j_k}$ for $i = 1, \dots, r$. Further, let $d' = \sum_{j \in N \setminus \mathcal{C}} a_j - \lambda$. Define

$$1286 \quad f(z) = \begin{cases} i\lambda & \text{if } A_i \leq z \leq A_{i+1} - \lambda, \quad i = 0, \dots, r - 1, \\ z - A_i + i\lambda & \text{if } A_i - \lambda \leq z \leq A_i, \quad i = 1, \dots, r - 1, \\ z - A_r + r\lambda & \text{if } A_r - \lambda \leq z \leq d'. \end{cases} \quad (94)$$

1287 Then, the lifted simple generalized flow cover inequality (LSGFICI)

$$1288 \quad \sum_{j \in \mathcal{L}} \lambda x_j + \sum_{j \in \mathcal{C}} f(a_j) x_j + \sum_{j \in N \setminus (\mathcal{C} \cup \mathcal{L})} a_j y_j \geq \lambda + \sum_{j \in \mathcal{C}} f(a_j) \quad (95)$$

1289 is facet-defining for PF .

1290 *Proof.* For a given generalized cover \mathcal{C} of F , we define $C = \mathcal{C} \cup \{l\}$ where $l \in \mathcal{L} \neq \emptyset$. Set C is a cover
1291 since $a_j > \lambda$ for all $j \in \mathcal{L}$. Further, $\sum_{j \in C} a_j = d + a_l - \lambda > d$ and so $\mu = a_l - \lambda > 0$, i.e. C satisfies
1292 Conditions (A1) and (A3) in Theorem 3.20. Now set $M = \mathcal{L} \setminus \{l\}$ in (66). Condition (A4) in Theorem 3.20
1293 also holds since $\sum_{j \in N \setminus (\mathcal{L} \setminus \{l\})} a_j - d = \sum_{j \in N \setminus \mathcal{L}} a_j + a_l - d > 0$. Next, we observe that $C \cup M = \mathcal{C} \cup \mathcal{L}$ and
1294 that $\min\{a_i, a_l - \mu\} = \min\{a_i, \lambda\} = \lambda = a_l - \mu$ for all $i \in M$. Substituting $a_l - \mu = \lambda$ in Proposition 3.17,
1295 we obtain that $f(w) = -P^M(-w)$ since $M \cup \{l\} = \mathcal{L}$. Therefore, we conclude that (95) is a lifted reverse
1296 bilinear cover inequality (66). \square

1297 Because in Gu et al. [12] the fixed-charge single-node flow set studied is more general than F , the authors
1298 focused mainly on the derivation of valid inequalities and discussed only indirectly whether the resulting
1299 inequalities are facet-defining. The result of Corollary 4.17 is therefore different from that of Theorem 12 in
1300 Gu et al. [12] in two ways. First we added the condition $\sum_{j \in N \setminus \mathcal{L}} a_j > d$. This condition guarantees that
1301 the simple generalized flow cover inequality (SGFCFI) that is used as seed inequality for lifting procedures in
1302 Gu et al. [12] is facet-defining for the problem restriction. Second, we replaced the statement that inequality
1303 (95) is valid for PF with the stronger statement that it is facet-defining for PF .

1304 We conclude this section by presenting conditions under which the lifted clique inequalities (78) are
1305 facet-defining for the flow set PF .

1306 **Theorem 4.18.** A lifted clique inequality (78) is facet-defining for PF if (i) $\sum_{j \in K} a_j - a_k > d$ for all $k \in K$
1307 and (ii) lifting coefficients are chosen according to (79) and (iii) one of the following conditions holds:

- 1308 1. $L = \emptyset$.
- 1309 2. $\exists \bar{i} \in M$ such that $j_{\bar{i}} = 0$ and, for all $i \in L \setminus \{\bar{i}\}$, $j_i = q_i$.

1310 *Proof.* Using a proof technique similar to that used in Theorems 4.12 and 4.16, we show that seed inequality
1311 (71) is facet-defining for $PF(\overline{M} \setminus \widehat{M}, \emptyset, \overline{M} \setminus \widehat{M}, \emptyset)$ and that lifting (x_i, y_i) for $i \in M$ adds two tight independent
1312 points in (78) that belong to F . Let $K = \{1, \dots, l\}$ and $\widehat{M} = \{l + 1, \dots, h\}$. Define χ such that $\chi_j = 1$ for
1313 $j \leq l$ and 0 for $l + 1 \leq j \leq k$. Consider $p^i = (\chi - e_i, \chi - e_i)$ for $i = 1, \dots, l$, $q^i = (\chi - e_i, \chi - e_i - \epsilon e_{i+1})$ for
1314 $i = 1, \dots, l - 1$, $q^l = (\chi - e_l, \chi - e_l - \epsilon e_1)$ where ϵ is positive, and, for $j = l + 1, \dots, h$, $r^j = (\chi - e_1 + e_j, \chi - e_1)$
1315 and $s^j = (\chi - e_1 + e_j, \chi - e_1 + e_j)$. These points satisfy (71) at equality, are affinely independent and,
1316 because of Assumption (i), belong to F when ϵ is sufficiently small. This shows that (71) is facet-defining
1317 for $PF(M, \emptyset, M, \emptyset)$. Assume first that $L = \emptyset$ and consider now the lifting of variables (x_i, y_i) for $i \in M$ in the
1318 proof of Theorem 3.26. For $j_i \in \{0, \dots, q_i\}$, lifting adds the two independent points $\left(1, \frac{W_j^i}{a_i}\right)$ and $\left(1, \frac{W_{j+1}^i}{a_i}\right)$
1319 that both correspond to feasible solutions of F because of (71) and Corollary 4.5, proving the result. Then
1320 it follows from the first part of this proof that the inequality obtained after lifting the variables in $M \setminus L$ is
1321 facet-defining for $PF(L \setminus \{\bar{i}\}, \emptyset, L \setminus \{\bar{i}\}, \emptyset)$. Consider now the lifting of variables (x_i, y_i) for $i \in L \setminus \{\bar{i}\}$. When
1322 $j_i = q_i$, we derived in the proof of Theorem 3.26 that lifting adds the two independent points $(1, 1)$ and
1323 $\left(1, \frac{W_{j_i}^i}{a_i}\right)$ that both correspond to feasible solutions of F because the first point sets $(x_{\bar{i}}, y_{\bar{i}}) = \left(1, \frac{B_{q_i} + \bar{\mu} - a_i}{a_{\bar{i}}}\right)$,
1324 and the structure of (71) satisfies the assumptions of Corollary 4.5. \square

1325 To the best of our knowledge, Theorem 4.18 presents a new family of facet-defining inequalities for fixed-
1326 charge single-node flow models without inflows.

5 Discussion and Conclusion

Many of the results presented in this paper extend to 0–1 mixed integer sets defined by constraints of the form $\sum_{i=1}^k (a_i x_i y_i + b_i x_i) + \sum_{i=k+1}^n a_i y_i \geq d$, *i.e.* bilinear covering sets where a linear term has been added to the left-hand-side. The primary reason the inequalities extend without significant changes is summarized in the next simple observation.

Proposition 5.1. *Consider an inequality $\alpha x + \beta y \geq \gamma$ such that for each i , $\alpha_i \beta_i = 0$. Let $H^+ = \{(x, y) \in \mathbb{R}^{2n} \mid \alpha x + \beta y \geq \gamma\}$. Let $I = \{i \in N \mid \alpha_i = 0\}$ and $I^c = N \setminus I$. Let $a_i \in \mathbb{R}$ and $(b_i, c_i) \in \mathbb{R}_+^2$ be such that, for each $i \in N$, $a_i + \min\{b_i, c_i\} \geq 0$. Consider the sets $A(w) = \{(x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{i=1}^n (a_i x_i y_i + b_i x_i + c_i y_i) \geq d - w\}$ and $B(w) = \{(x, y) \in \{0, 1\}^n \times [0, 1]^n \mid \sum_{i \in I} (a_i + c_i) x_i y_i + \sum_{i \in I^c} (a_i + b_i) x_i y_i \geq d - \sum_{i \in I} b_i - \sum_{i \in I^c} c_i - w\}$. Then, $\min\{\alpha x + \beta y \mid (x, y) \in A(w)\} = \min\{\alpha x + \beta y \mid (x, y) \in B(w)\}$. Further, $H^+ \supseteq A(w)$ if and only if $H^+ \supseteq B(w)$.*

Proof. Consider a point (x', y') with $x'_i = 1$ for $i \in I$ and $y'_i = 1$ for $i \in I^c$. Then, for $i \in I$, $a_i x'_i y'_i + b_i x'_i + c_i y'_i = (a_i + c_i) x'_i y'_i + b_i$. Similarly, for $i \in I^c$, $a_i x'_i y'_i + b_i x'_i + c_i y'_i = (a_i + b_i) x'_i y'_i + c_i$. In other words, $\sum_{i \in I} (a_i + c_i) x'_i y'_i + \sum_{i \in I^c} (a_i + b_i) x'_i y'_i - \sum_{i \in I} b_i - \sum_{i \in I^c} c_i = \sum_{i=1}^n (a_i x'_i y'_i + b_i x'_i + c_i y'_i)$. Therefore, $(x', y') \in A(w)$ if and only if $(x', y') \in B(w)$. Now, it is easy to see that

$$\begin{aligned} z_{A(w)} &= \min\{\alpha x + \beta y \mid (x, y) \in A(w)\} = \min\{\alpha x + \beta y \mid (x, y) \in A(w), x_i = 1 \forall i \in I, y_i = 1 \forall i \in I^c\} \\ &= \min\{\alpha x + \beta y \mid (x, y) \in B(w), x_i = 1 \forall i \in I, y_i = 1 \forall i \in I^c\} = \min\{\alpha x + \beta y \mid (x, y) \in B(w)\} = z_{B(w)}, \end{aligned}$$

where the second and the second last equality follow from the assumptions which imply that $a_i x_i y_i + b_i x_i + c_i y_i \leq \min\{a_i x_i + b_i x_i + c_i, a_i y_i + b_i + c_i y_i\}$, $(a_i + c_i) x_i y_i \leq (a_i + c_i) y_i$, and $(a_i + b_i) x_i y_i \leq (a_i + b_i) x_i$. Since $z_{A(w)} = z_{B(w)}$ and $H^+ \supseteq A(w)$ (resp. $H^+ \supseteq B(w)$) if and only if $z_{A(w)} \geq \gamma$ (resp. $z_{B(w)} \geq \gamma$), it follows that $H^+ \supseteq A(w)$ if and only if $H^+ \supseteq B(w)$. \square

Note that the seed inequalities and the intermediate inequalities we derive during lifting satisfy the condition $\alpha_i \beta_i = 0$ for all i . Then, Proposition 5.1 essentially shows that the lifting functions derived for the problem with only bilinear terms on the left-hand-side also carry over to problems containing a linear term. For detailed derivations of facet-defining inequalities for bilinear covering sets with linear terms, we refer the reader to [7].

In this paper, we study the polyhedral structure of the 0–1 mixed-integer bilinear covering set. We give a complete linear description of its convex hull when $n = 2$. We also show that, for a fairly large class of functions, it is sufficient to check that subadditivity holds on a subset of points of the domain to show that the function is subadditive over \mathbb{R}^n . This result enables short subadditivity proofs for many practically useful functions. In particular, we use this result to derive three families of strong inequalities for PB that can be obtained using sequence-independent lifting. Among them, two families have an exponential number of members. We study relations between 0–1 mixed-integer bilinear covering sets and fixed-charge single-node flow sets without inflows. We show that valid inequalities for bilinear sets are also valid for flow sets and prove that all nontrivial facets of PF can be obtained through the study of facets of PB . We then show that the inequalities we derive generalize two classical families of lifted flow cover inequalities for PF and provide a new family for PF . Future research will focus on evaluating the computational benefits of using these lifted cuts in branch-and-bound frameworks for both linear and nonlinear mixed integer programming.

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1431 Appendix

1432 Linear descriptions of $\text{conv}(B)$ and $\text{conv}(F)$

1433 The linear description of the convex hulls of the bilinear set B and the flow set F are obtained by PORTA
1434 as the following:

$$1435 B = \left\{ (x, y) \in \{0, 1\}^4 \times [0, 1]^4 \mid 19x_1y_1 + 17x_2y_2 + 15x_3y_3 + 10x_4y_4 \geq 20 \right\}$$

1436	(1)	$50x_1$		$+90x_3$	$+45x_4$	$+76y_1$	$+153y_2$		≥ 135
1437	(2)	$70x_1$	$+90x_2$		$+27x_4$	$+38y_1$		$+135y_3$	≥ 117
1438	(3)	$25x_1$		$+65x_3$	$+45x_4$	$+76y_1$	$+153y_2$		≥ 110
1439	(4)		$+50x_2$	$+70x_3$	$+35x_4$	$+133y_1$	$+34y_2$		≥ 105
1440	(5)		$+25x_2$	$+45x_3$	$+35x_4$	$+133y_1$	$+34y_2$		≥ 80
1441	(6)	$21x_1$	$+41x_2$		$+27x_4$	$+38y_1$		$+135y_3$	≥ 68
1442	(7)	$30x_1$	$+35x_2$	$+21x_3$		$+19y_1$		$+70y_4$	≥ 56
1443	(8)	$18x_1$	$+23x_2$	$+21x_3$		$+19y_1$		$+70y_4$	≥ 44
1444	(9)	$19x_1$	$+17x_2$					$+15y_3$	$+10y_4 \geq 20$
1445	(10)	$19x_1$		$+15x_3$			$+17y_2$		$+10y_4 \geq 20$
1446	(11)	$19x_1$			$+10x_4$		$+17y_2$	$+15y_3$	≥ 20
1447	(12)	$19x_1$					$+17y_2$	$+15y_3$	$+10y_4 \geq 20$
1448	(13)		$+17x_2$	$+15x_3$		$+19y_1$			$+10y_4 \geq 20$
1449	(14)		$+17x_2$		$+10x_4$	$+19y_1$		$+15y_3$	≥ 20
1450	(15)		$+17x_2$			$+19y_1$		$+15y_3$	$+10y_4 \geq 20$
1451	(16)			$+15x_3$	$+10x_4$	$+19y_1$	$+17y_2$		≥ 20
1452	(17)			$+15x_3$		$+19y_1$	$+17y_2$		$+10y_4 \geq 20$
1453	(18)				$+10x_4$	$+19y_1$	$+17y_2$	$+15y_3$	≥ 20
1454	(19)					$+19y_1$	$+17y_2$	$+15y_3$	$+10y_4 \geq 20$
1455	(20)	$14x_1$		$+10x_3$	$+5x_4$		$+17y_2$		≥ 15

1456	(21)		+12x ₂	+10x ₃	+5x ₄	+19y ₁				≥ 15
1457	(22)			+10x ₃	+5x ₄	+19y ₁	+17y ₂			≥ 15
1458	(23)	12x ₁	+10x ₂		+3x ₄			+15y ₃		≥ 13
1459	(24)		+10x ₂	+10x ₃	+3x ₄	+19y ₁				≥ 13
1460	(25)		+10x ₂		+3x ₄	+19y ₁		+15y ₃		≥ 13
1461	(26)	10x ₁	+10x ₂		+x ₄			+15y ₃		≥ 11
1462	(27)	10x ₁		+10x ₃	+x ₄		+17y ₂			≥ 11
1463	(28)	10x ₁			+x ₄		+17y ₂	+15y ₃		≥ 11
1464	(29)	7x ₁	+5x ₂	+3x ₃				+10y ₄		≥ 8
1465	(30)		+5x ₂	+3x ₃		+19y ₁		+10y ₄		≥ 8
1466	(31)		+5x ₂	+3x ₃	+5x ₄	+19y ₁				≥ 8
1467	(32)		+3x ₂	+3x ₃	+3x ₄	+19y ₁				≥ 6
1468	(33)	5x ₁		+x ₃			+17y ₂	+10y ₄		≥ 6
1469	(34)	5x ₁	+5x ₂	+x ₃				+10y ₄		≥ 6
1470	(35)	5x ₁		+x ₃	+5x ₄		+17y ₂			≥ 6
1471	(36)	3x ₁	+x ₂					+15y ₃	+10y ₄	≥ 4
1472	(37)	3x ₁	+x ₂	+3x ₃				+10y ₄		≥ 4
1473	(38)	3x ₁	+x ₂		+3x ₄			+15y ₃		≥ 4
1474	(39)	x ₁	+x ₂	+x ₃				+10y ₄		≥ 2
1475	(40)	x ₁	+x ₂		+x ₄			+15y ₃		≥ 2
1476	(41)	x ₁		+x ₃	+x ₄		+17y ₂			≥ 2
1477	(42)	x ₁	+x ₂	+x ₃	+x ₄					≥ 2
1478	(43)	x ₁								≥ 0
1479	(44)		x ₂							≥ 0
1480	(45)			x ₃						≥ 0
1481	(46)				x ₄					≥ 0
1482	(47)					y ₁				≥ 0
1483	(48)						y ₂			≥ 0
1484	(49)							y ₃		≥ 0
1485	(50)								y ₄	≥ 0
1486	(51)								y ₄	≤ 1
1487	(52)							y ₃		≤ 1
1488	(53)						y ₂			≤ 1
1489	(54)					y ₁				≤ 1
1490	(55)				x ₄					≤ 1
1491	(56)			x ₃						≤ 1
1492	(57)		x ₂							≤ 1
1493	(58)	x ₁								≤ 1

1494 $F = \left\{ (x, y) \in \{0, 1\}^4 \times [0, 1]^4 \mid 19y_1 + 17y_2 + 15y_3 + 10y_4 \geq 20, x_j \geq y_j \forall j = 1, \dots, 4 \right\}$

1495	(1)	50x ₁		+90x ₃	+45x ₄	+76y ₁	+153y ₂			≥ 135
1496	(2)	70x ₁	+90x ₂		+27x ₄	+38y ₁		+135y ₃		≥ 117
1497	(3)	25x ₁		+65x ₃	+45x ₄	+76y ₁	+153y ₂			≥ 110
1498	(4)		+50x ₂	+70x ₃	+35x ₄	+133y ₁	+34y ₂			≥ 105
1499	(5)		+25x ₂	+45x ₃	+35x ₄	+133y ₁	+34y ₂			≥ 80
1500	(6)	21x ₁	+41x ₂		+27x ₄	+38y ₁		+135y ₃		≥ 68
1501	(7)	30x ₁	+35x ₂	+21x ₃		+19y ₁			+70y ₄	≥ 56
1502	(8)	18x ₁	+23x ₂	+21x ₃		+19y ₁			+70y ₄	≥ 44
1503	(19)					+19y ₁	+17y ₂	+15y ₃	+10y ₄	≥ 20
1504	(22)			+10x ₃	+5x ₄	+19y ₁	+17y ₂			≥ 15
1505	(24)		+10x ₂	+10x ₃	+3x ₄	+19y ₁				≥ 13

1506	(25)		+10x ₂		+3x ₄	+19y ₁		+15y ₃	≥ 13
1507	(26)	10x ₁	+10x ₂		+x ₄			+15y ₃	≥ 11
1508	(27)	10x ₁		+10x ₃	+x ₄		+17y ₂		≥ 11
1509	(28)	10x ₁			+x ₄		+17y ₂	+15y ₃	≥ 11
1510	(30)		+5x ₂	+3x ₃		+19y ₁		+10y ₄	≥ 8
1511	(31)		+5x ₂	+3x ₃	+5x ₄	+19y ₁			≥ 8
1512	(32)		+3x ₂	+3x ₃	+3x ₄	+19y ₁			≥ 6
1513	(33)	5x ₁		+x ₃			+17y ₂	+10y ₄	≥ 6
1514	(34)	5x ₁	+5x ₂	+x ₃				+10y ₄	≥ 6
1515	(35)	5x ₁		+x ₃	+5x ₄		+17y ₂		≥ 6
1516	(36)	3x ₁	+x ₂					+15y ₃ +10y ₄	≥ 4
1517	(37)	3x ₁	+x ₂	+3x ₃				+10y ₄	≥ 4
1518	(38)	3x ₁	+x ₂		+3x ₄			+15y ₃	≥ 4
1519	(39)	x ₁	+x ₂	+x ₃				+10y ₄	≥ 2
1520	(40)	x ₁	+x ₂		+x ₄			+15y ₃	≥ 2
1521	(41)	x ₁		+x ₃	+x ₄	+17y ₂			≥ 2
1522	(42)	x ₁	+x ₂	+x ₃	+x ₄				≥ 2
1523	(47)					y ₁			≥ 0
1524	(48)						y ₂		≥ 0
1525	(49)							y ₃	≥ 0
1526	(50)								y ₄ ≥ 0
1527	(55)				x ₄				≤ 1
1528	(56)			x ₃					≤ 1
1529	(57)		x ₂						≤ 1
1530	(58)	x ₁							≤ 1
1531	(f1)	x ₁				-y ₁			≥ 0
1532	(f2)		x ₂				-y ₂		≥ 0
1533	(f3)			x ₃				-y ₃	≥ 0
1534	(f4)				x ₄				-y ₄ ≥ 0