

A universal and structured way to derive dual optimization problem formulations

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December 19, 2019

Abstract

The dual problem of a convex optimization problem can be obtained in a relatively simple and structural way by using a well-known result in convex analysis, namely Fenchel's duality theorem. This alternative way of forming a strong dual problem is the subject in this paper. We recall some standard results from convex analysis and then discuss how the dual problem can be written in terms of the conjugates of the objective function and the constraint functions. This is a didactically valuable method to explicitly write the dual problem. We demonstrate the method by deriving dual problems for several classical problems and also for a practical model for radiotherapy treatment planning for which deriving the dual problem using other methods is a more tedious task. Additional material is presented in appendices, including useful tables for finding conjugate functions of many functions.

Key words. convex optimization, duality theory, conjugate functions, support functions, Fenchel duality

AMS subject classifications 46N10, 65K99, 90C25, 90C46

1 Introduction

Nowadays optimization is one of the most widely taught and applied subjects of mathematics. A mathematical optimization problem consists of an objective function that has to be minimized (or maximized) subject to some constraints. Every maximization problem can be written as a minimization problem by

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21 replacing the objective function by its opposite. In this paper we always assume that the given problem is
22 a minimization problem. To every such problem one can assign another problem, called the *dual problem*
23 of the original problem; the original problem is called the *primal problem*. In many cases a dual problem
24 exists that has the same optimal value as the primal problem. If this happens then we say that we have
25 strong duality, otherwise weak duality. It is well-known that if the primal problem is convex then there
26 exists a strong dual problem under some mild conditions. In this paper we consider only convex problems
27 with a strong dual. Strong duality and obtaining explicit formulations of the dual problem are of great
28 importance for at least four reasons:

- 29 1. Sometimes the dual problem is easier to solve than the primal problem.
- 30 2. If we have feasible solutions of the primal problem and the dual problem, respectively, and their
31 objective values are equal, then we may conclude that both solutions are optimal. In that case the
32 primal solution is a certificate for the optimality of the dual solution, and vice versa.
- 33 3. If the two objective values are not equal then the dual objective value is a lower bound for the
34 optimal value of the primal problem and the primal objective value is an upper bound for the
35 optimal value of the dual problem. This provides information on the ‘quality’ of the two solutions,
36 i.e., on how much their values may differ from the optimal value.
- 37 4. Some optimization problems benefit from being solved with the dual decomposition method. This
38 method requires repeated evaluation of the dual objective value of two (smaller) problems for fixed
39 dual variables. Our explicit formulation of the dual can potentially speed up the evaluation of the
40 dual objective function.

41 In textbooks, different ways for deriving dual optimization problems are discussed. One option is to
42 write the optimization problem as a conic optimization problem, and then state the dual optimization
43 problem in terms of the dual cone. However, rewriting general convex constraints as conic constraints
44 may be awkward or even not possible. Another option is to derive the dual via Lagrange duality, which
45 often leads to an unnecessarily long and tedious derivation process. We propose a universal and more
46 structured way to derive dual problems, using basic results from convex analysis. The duals derived this
47 way are ordinary Lagrange duals. The value of this paper is therefore not to present new duality theory,
48 but rather to provide a unifying, pedagogically relevant framework of deriving the dual.

49 Let the primal optimization problem be given. Then by using Fenchel’s duality theorem one can derive a
50 dual problem in terms of so-called perspectives of the conjugates of the objective function and constraint

51 functions. The main task is therefore to derive expressions for the conjugates of each of these functions
52 separately (taking the perspective functions is easy). At first this may seem to be a difficult task –
53 and this may be the reason that in textbooks this is not further elaborated on. However, the key issue
54 underlying this paper is that it is not necessary to derive explicit expressions for the conjugates of these
55 functions: if one can derive an expression for the conjugate as an “inf over some variables”, then this is
56 sufficient to state the dual problem. And, indeed, by using standard composition rules for conjugates –
57 well-known results in convex analysis – we arrive at such expressions with an inf.

58 Therefore, the aim of this paper is to show that the dual problem can be derived by using the composition
59 rules for conjugates and by using the well-known results for the conjugates of basic univariate and
60 multivariate functions. Thus, the proposed method is similar in nature to the way Fourier and Laplace
61 transforms are taught, where a table with known transforms of functions is combined with a table with
62 identities for transformations of functions. We consider this the most structured way to teach deriving
63 dual optimization problems. Using this structured way, one might even computerize the derivation of the
64 dual optimization problem. We must emphasize that the ingredients in this approach are not new. About
65 the same dual is presented in [8, p. 256], where it is obtained via Lagrange duality, and in [15, p.45]. We
66 also mention [24] where three strategies for deriving a dual problem are discussed, and Fenchel’s duality
67 theorem is used to construct the dual problem of a specific second order cone problem. As far as we
68 know, no comprehensive survey as presented in this paper exists.

69 We should also mention that in special cases where one would like to derive a partial Lagrange dual
70 problem, e.g., for relaxation purposes, Fenchel’s duality formulation cannot be used. On the other hand,
71 the use of Fenchel’s dual has found application in recent developments in the field of Robust Optimization,
72 see, e.g., [6, 13].

73 The outline of the paper is as follows. In the next section we recall some fundamental notions from convex
74 analysis, as defined in [19]. In Section 3 we present the general form of the primal problem that we want
75 to solve and Fenchel’s dual problem. We also discuss under which condition strong duality is guaranteed.
76 To make the paper self-supporting, we present a proof of this strong duality in the Appendices C and D.
77 In Section 4 we present a scheme that describes in four steps how to get Fenchel’s dual problem. In this
78 scheme, tables of conjugate functions and support functions in Appendix E are of crucial importance. In
79 Appendix A we demonstrate how conjugates for some functions are computed.

80 The use of the scheme in Section 4 is illustrated in Section 6, where we show how dual problems for
81 several classical convex optimization problems can be obtained. In these classical problems the functions

82 that appear in the primal problem all have a similar structure. However, also problems with a variety of
 83 constraint functions can be handled systematically. In Section 6.7 we demonstrate this by applying our
 84 approach to a model for radiotherapy treatment.

85 2 Preliminaries

86 We recall in this section some notations and terminology that are common in the literature. The *effective*
 87 *domain* of a convex function $f : \mathbf{R}^n \rightarrow [-\infty, +\infty]$ is defined by

$$88 \quad \text{dom } f = \{x \mid f(x) < \infty\}.$$

89 A convex function f is *proper* if $f(x) > -\infty$ for all $x \in \mathbf{R}^n$ and $f(x) < \infty$ for at least one $x \in \mathbf{R}^n$. The
 90 function f is *closed* if the set $\{x \mid f(x) \leq \alpha\}$ is closed for every $\alpha \in \mathbf{R}$.

91 For any function f its convex conjugate f^* is defined by

$$92 \quad f^*(y) := \sup_{x \in \text{dom } f} \{y^T x - f(x)\}. \quad (1)$$

93 Figure 1 illustrates this notion graphically for the case where $n = 1$.

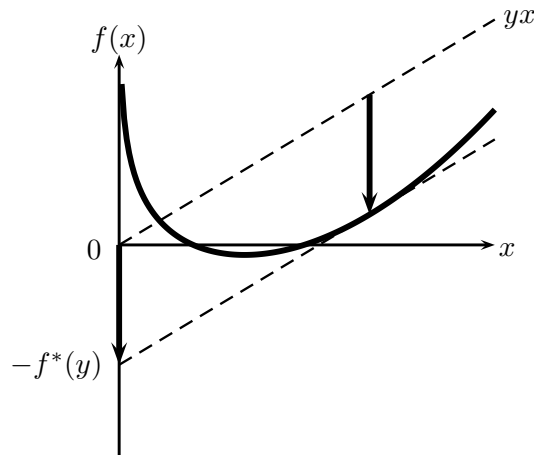


Figure 1: Value of convex conjugate at y

94 Given y , the maximal value of $yx - f(x)$ occurs where the derivative of $f(x)$ equals y , and the tangent
 95 of the graph of f where this happens intersects the y -axis in the point $(0, -f^*(y))$.

96 Since $f^*(y)$ is the pointwise maximum of linear functions (linear in y), f^* is convex. Moreover, if f is
 97 convex then f^* is closed, $(\text{cl } f)^* = f^*$ and $f^{**} = \text{cl } f$, which implies $f^{**} = f$ if f is closed convex [19,

98 Theorem 12.2].

99 The *indicator function* of a set $S \subseteq \mathbf{R}^n$ is denoted as $\delta_S(x)$ and defined by

$$100 \quad \delta_S(x) = \begin{cases} 0, & x \in S \\ \infty, & x \notin S. \end{cases}$$

101 Note that $\delta_S(x)$ is convex if and only if S is convex, and proper if S is nonempty. Obviously, we have

$$102 \quad x \in S \quad \Leftrightarrow \quad \delta_S(x) \leq 0.$$

103 This implies that any convex *set-constraint* can be replaced with a convex *functional constraint* and vice
104 versa.

105 The conjugate function of $\delta_S(x)$ is called the *support function* of S and given by

$$106 \quad \delta_S^*(y) = \sup_{x \in S} \{y^T x\}. \quad (2)$$

107 Our approach highly depends on a simple relation between the conjugate f^* of a convex function f and
108 the support function of the set $S := \{x \mid f(x) \leq 0\}$. Provided that S is nonempty, and $\text{ri } S$ is nonempty
109 if f is nonlinear, we have (cf. Lemma D.1):

$$110 \quad \delta_S^*(y) = \min_{\lambda \geq 0} \{(\lambda f)^*(y)\}. \quad (3)$$

111 Here the function λf is defined in the natural way: $(\lambda f)(x) = \lambda f(x)$. If $\lambda > 0$ then one has

$$112 \quad (\lambda f)^*(y) = \sup_x \{y^T x - \lambda f(x)\} = \sup_x \left\{ \lambda \left(x^T \frac{y}{\lambda} - f(x) \right) \right\} = \lambda f^* \left(\frac{y}{\lambda} \right), \quad (4)$$

113 and for $\lambda = 0$ we get*

$$114 \quad (0f)^*(y) = \sup_x \{y^T x\} = \begin{cases} 0, & \text{if } y = 0 \\ \infty & \text{otherwise} \end{cases} = \delta_{\{0\}}(y). \quad (5)$$

115 The function $(\lambda f)^*$ is called the *perspective function* of f^* (with λ as the new variable). Due to the
116 first equality in (4), the perspective function of f^* is convex (more precisely: jointly convex in y and λ).
117 Although $(\lambda f)^*$ is closed for each fixed value of λ , it is not necessarily closed for λ . It may therefore

*Here we adopt the calculation rules $0\infty = \infty 0 = 0$, as in [19, p. 24].

118 occur that $(\lambda f)^*(y)$ is well-defined even if λ approaches zero whereas y stays away from zero. Closing
119 the perspective function for λ can therefore result in the wrong dual. An example of this phenomenon
120 can be found in Section B.9.

121 The concave conjugate g_* of a function g is defined by

$$122 \quad g_*(y) := \inf_{x \in \text{dom } g} \{y^T x - g(x)\}.$$

123 Since $g_*(y)$ is the pointwise minimum of linear functions, g_* is concave. Putting $f = -g$, one easily
124 verifies the following relation [19, p. 308]:

$$125 \quad g_*(y) = -f^*(-y) = -(-g)^*(-y). \quad (6)$$

126 Hence, all properties of convex conjugates lead to similar properties of concave conjugates.

127 **3 Fenchel's dual problem of a convex optimization problem**

128 In this section we derive Fenchel's dual problem for a convex optimization problem by using conjugate
129 functions. We will use some properties of conjugate functions without proving them at this stage; this
130 will be the subject in the following sections and in the Appendix. So for the moment we take these
131 properties for granted.

132 Let us first consider the case of *unconstrained* minimization. So, we have a convex function f_0 and want
133 to know its minimal value z^* . Due to definition (1) of f_0^* we may write

$$134 \quad z^* = \inf_{x \in \mathbf{R}} \{f_0(x)\} = -\sup_x \{-f_0(x)\} = -f_0^*(0). \quad (7)$$

135 This already reveals the relevance of the notion of a conjugate function for this relatively simple opti-
136 mization problem: if we know the conjugate function $f_0^*(y)$ of f_0 then $-f_0^*(0)$ gives the minimal value.
137 As we will see this simple property is the keystone of Fenchel's duality theory.

138 Next we consider the case where S is a convex subset of \mathbf{R} and we want to find the minimal value of f_0
139 on the set S . So the problem we want to solve is

$$140 \quad \inf_x \{f_0(x) \mid x \in S\}.$$

141 In that case we consider the function $g(x)$ defined by

$$142 \quad g(x) := f_0(x) + \delta_S(x).$$

143 Obviously the effective domain of $g(x)$ is the set S and on this set g coincides with f_0 . So, we have
 144 transformed our *constrained* problem to an unconstrained problem, and therefore the optimal value will
 145 be $-g^*(0)$. The question arises how to find the conjugate function of g . It turns out that it can be derived
 146 easily from the conjugate functions of f_0 and δ_S . Using the formula in line 6 of Table 2 we obtain

$$147 \quad g^*(y) = \min_{y^0, y^1} \{f_0^*(y^0) + \delta_S^*(y^1) \mid y^0 + y^1 = y\}.$$

148 We need to compute $g^*(y)$ at $y = 0$. Substitution of $y = 0$ in the above expression yields $y^0 + y^1 = 0$,
 149 i.e., $y^1 = -y^0$. So, we can eliminate y^1 . Thus we get

$$150 \quad g^*(0) = \min_{y^0} \{f_0^*(y^0) + \delta_S^*(-y^0)\}.$$

151 In general the problem that we want to solve is the following (primal) convex optimization problem

$$152 \quad (P) \quad \inf_x \{f_0(x) \mid f_i(x) \leq 0, i = 1, \dots, m\},$$

153 where the functions $f_i : \mathbf{R}^n \rightarrow [-\infty, \infty]$ are proper convex functions, for $i = 0, \dots, m$. In that case the
 154 set S is the intersection of the sets

$$155 \quad S_i = \{x \mid f_i(x) \leq 0\}, \quad 1 \leq i \leq m.$$

156 Obviously the indicator function of S satisfies

$$157 \quad \delta_S(x) = \delta_{S_1}(x) + \dots + \delta_{S_m}(x)$$

158 and hence, by line 6 in Table 3,

$$159 \quad \delta_S^*(-y^0) = \min_{y^1, \dots, y^m} \{\delta_{S_1}^*(y^1) + \dots + \delta_{S_m}^*(y^m) \mid y^1 + \dots + y^m = -y^0\}.$$

160 By (3) we have

$$161 \quad \delta_{S_i}^*(y^i) = \min_{u_i \geq 0} \{(u_i f_i)^*(y^i)\}, \quad 1 \leq i \leq m.$$

162 Substitution gives

$$163 \quad \delta_S^*(-y^0) = \min_{\{y^i\}_{i=1}^m} \left\{ \sum_{i=1}^m \min_{u_i \geq 0} \{(u_i f_i)^*(y^i)\} \mid \sum_{i=0}^m y^i = 0 \right\}$$

$$164 \quad = \min_{\{y^i\}_{i=1}^m, u \succeq 0} \left\{ \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0 \right\}.$$

165 Substitution into the expression for $g^*(0)$ gives

$$166 \quad g^*(0) = \min_{(y^i)_{i=0}^m, u \succeq 0} \left\{ f^*(y^0) + \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0 \right\}.$$

167 By changing the sign at both sides this implies

$$168 \quad -g^*(0) = \max_{(y^i)_{i=0}^m, u \succeq 0} \left\{ -f^*(y^0) - \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0 \right\}.$$

169 Thus we may conclude that the optimal value of (P) is the same as that of the problem

$$170 \quad (D) \quad \sup_{\{y^i\}_{i=0}^m, u} \left\{ -f_0^*(y^0) - \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0, u \succeq 0 \right\}.$$

171 This is Fenchel's dual problem of (P) .[†]

172 Note that the objective function in (D) is concave (since the perspective functions under the sum are
173 convex) and the constraints are linear (so convex as well). Hence, when writing (D) as a minimization
174 problem it becomes a convex problem.

175 It is clear that (D) is a practical formulation of the dual if the convex conjugates f_i^* can easily be
176 computed, which turns out to be true in many cases. This important fact is the main motivation for this
177 paper.

178 In the sequel, we make the following assumption to satisfy the condition in line 6 of Table 2.

179 **Assumption 3.1** Assume that $\text{ri.dom } f_0 \cap S \neq \emptyset$ or $\bigcap_{i=0}^m \text{ri.dom } f_i \neq \emptyset$.

180 This is the same as requiring that (P) is *Slater regular* [21]. In that case the optimal values of (P) and
181 (D) are equal; for a proof we refer to Appendix D. We should add that if f_0 is linear then $x \in \text{ri.dom } f_0$

[†]As in [8], the notation $x \succeq 0$ defines a vector $x \in \mathbf{R}^n$ to be nonnegative; only if $n = 1$ we use the notation $x \geq 0$.

182 can be replaced by $x \in \text{dom } f_0$ in the Slater condition.

183 In the sequel, especially when the objective function has the same form as the constraint functions, we
184 extend the vector u with the entry $u_0 = 1$. Then (D) can be written as

$$185 \quad (D') \quad \sup_{\{y^i\}_{i=0}^m, u} \left\{ - \sum_{i=0}^m (u_i f_i)^* (y^i) \mid \sum_{i=0}^m y^i = 0, u \succeq 0, u_0 = 1 \right\}.$$

186 In the above expressions for (D) and (D') we used the sup operator to indicate that we are dealing with
187 a maximization problem. Below we use the max operator only if the maximum value is attained. In this
188 paper it is not our goal to establish when this happens; our main focus is to find a dual problem yielding
189 strong duality. Therefore, if we use the sup operator in a maximization problem below, it is not excluded
190 that the maximal value is attained, and similarly for the inf operator.

191 4 Recipe for setting up Fenchel's dual problem

192 As became clear in the previous section, Fenchel's dual gives us a straightforward method for finding the
193 dual problem of (P) . The method can be split into four steps as discussed below. We will give some
194 demonstrations of the resulting recipe in Section 6.

195 Step 1: Formulate the primal problem

196 The primal problem should be a convex optimization problem. Moreover, all set constraints should be
197 replaced by their indicator functions. Strong duality, i.e, equality of the optimal values of (P) and (D) ,
198 is guaranteed only if (P) is Slater regular.

199 Step 2: Derive tractable expressions for the conjugates

200 In our approach we need the conjugates of the functions that occur in the problem. For that purpose
201 one can use the three tables in Appendix E. Table 1 contains expressions for the conjugates of some
202 basic functions and the domains of these functions, Table 2 some formulas for conjugates of function
203 transformations, and Table 3 some support functions of common sets. Usually a conjugate function (or
204 its perspective function) occurring in (D) or (D') is written as $\inf_y \{h(y, z)\}$, where $h(y, z)$ is jointly

205 convex in y and z . A frequently used tool in our approach is the obvious equality

$$206 \quad \sup_z \left\{ - \inf_y h(y, z) \right\} = \sup_{y, z} \{ -h(y, z) \}.$$

207 In other words, when writing down (D) or (D') we can get rid of inf operators by moving the concerning
 208 variables under the sup operator. This helps to simplify the formulation of the dual problem. Deriving
 209 tractable expressions for the conjugates is further elaborated on and illustrated in Appendix A. The
 210 domain of a conjugate function is also important. For each i the vector y^i belongs to the domain of
 211 $(u_i f_i)^*$ in (D) or (D') . This not only leads to convex constraints in Fenchel's dual problem; in most
 212 cases it also enables to eliminate the variables y^i . In some cases we get a conjugate function that consists
 213 of 'branches', i.e., one obtains different formulas for $f^*(y)$ on several parts of the domain of f^* . This
 214 phenomenon is demonstrated in Appendix B.7, where it is also shown how this can be prevented.

215 **Step 3: Derive tractable convex expressions for perspectives of conjugates**

216 The conjugate functions found in step 2 are substituted into the term $(uf)^*$ in (D) or (D') . For $u > 0$
 217 this term is

$$218 \quad u f^* \left(\frac{y}{u} \right), \text{ with } \frac{y}{u} \in \text{dom } f^*,$$

219 One may be troubled by the fact that the quotient y/u is not jointly convex in y and u . Indeed,
 220 sometimes the resulting formulation of the dual problem may be not convex, even though the set $\{(y, u) :$
 221 $u > 0, y/u \in \text{dom } f^*\}$ is convex. A simple substitution yields a convex formulation. For example, when
 222 using line 8b in Table 2 one has

$$223 \quad f^*(y) = \inf_z \{ h^*(z) - b^T z \mid A^T z = y \}.$$

224 This gives rise to the term $(uf)^*(y)$ in the dual problem, i.e.,

$$225 \quad (uf)^*(y) = \inf_z \left\{ u h^*(z) - u b^T z \mid A^T z = \frac{y}{u} \right\}, \quad u > 0$$

$$226 \quad = \inf_z \{ u h^*(z) - b^T(uz) \mid A^T(uz) = y \}, \quad u > 0.$$

227 Due to the introduction of the variable u we are left with an expression that is not jointly convex in y
 228 and u , because of the occurrence of the products $u h^*(z)$ and uz . However, by the substitution $\tilde{z} = uz$ we

229 get

$$\begin{aligned} 230 \quad (uf)^*(y) &= \inf_{\tilde{z}} \left\{ uh^* \left(\frac{\tilde{z}}{u} \right) - b^T \tilde{z} \mid A^T \tilde{z} = y \right\} \\ 231 \quad &= \inf_{\tilde{z}} \left\{ (uh)^*(\tilde{z}) - b^T \tilde{z} \mid A^T \tilde{z} = y \right\}. \end{aligned}$$

232 Obviously, both the objective function and the constraint are now convex. In Section 6 one will find many
233 substitutions of this type, that are necessary to obtain a convex (and hence computationally tractable)
234 formulation of the dual problem.

235 Step 4: Formulate the dual problem

236 The results of Steps 1-3 are used in the dual formulation (D) or (D'). Often one can simplify the resulting
237 formulation by eliminating variables.

238 5 Sensitivity analysis

239 The existing theory for sensitivity analysis is based on Lagrange multipliers (e.g., [8, par 5.6.3]). We
240 therefore show that the variables u in (D) are precisely the Lagrange multipliers.

241 The classical approach to sensitivity theory investigates how the objective value of the dual problem
242 depends on the right-hand sides in the primal problem.

243 We first recall the classical approach to sensitivity theory. The Lagrange dual of (P) is given by

$$244 \quad (LD) \quad \sup_{\lambda \geq 0} \inf_x \left\{ f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) \right\}.$$

245 Let $(x^\diamond, \lambda^\diamond)$ be an optimal solution. If the i -th constraint $f_i(x) \leq 0$ is replaced by $f_i(x) + \varepsilon_i \leq 0$, with
246 $\varepsilon_i \in \mathbf{R}$, then the optimal value z^\diamond becomes

$$247 \quad z^\diamond = f_0(x^\diamond) + \sum_{i=1}^m \lambda_i^\diamond (f_i(x^\diamond) + \varepsilon_i) = f_0(x^\diamond) + \sum_{i=1}^m \lambda_i^\diamond f_i(x^\diamond) + \sum_{i=1}^m \lambda_i^\diamond \varepsilon_i.$$

248 Due to the above expression the Lagrange multiplier λ_i is called the *shadow price* for the i -th constraint.
249 This is because one usually concludes that adding a small number ε_i to $f_i(x)$ will change the objective
250 value with the amount $\varepsilon_i \lambda_i$. Though this might be true in many cases, it does not hold in general,
251 especially not if either the primal or the dual problem is degenerate. For examples we refer to [16].

252 We now show that the variables u_i in Fenchel's dual problem are also shadow prices in the above sense.

253 Recall that Fenchel's dual of (P) is given by

$$254 \quad (D) \quad \sup_{y^i, u \geq 0} \left\{ -f_0^*(y^0) - \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0 \right\}.$$

255 Now $(u_i(f_i + \varepsilon_i))^*(y^i) = (u_i f_i)^*(y^i) - u_i \varepsilon_i$, by line 1 in Table 2. Hence at optimality the part of the
 256 objective value in the above problem that is due to the perturbation in the primal problem is just
 257 $\sum_{i=0}^m u_i^\diamond \varepsilon_i$, where u^\diamond is optimal for (D) . This proves that in sensitivity analysis u takes over the role of
 258 the vector λ of Lagrange multipliers.

259 6 Some applications

260 In this section we demonstrate how Fenchel's dual can be obtained for some well-known convex optimiza-
 261 tion problems and a mathematical model for radiotherapy treatment planning.

262 6.1 Linear optimization

263 We consider the standard linear optimization problem:

$$264 \quad \inf_x \{ c^T x \mid Ax = b, x \in \mathbf{R}_+^n \},$$

265 where A is an $m \times n$ matrix and $b \in \mathbf{R}^m, c \in \mathbf{R}^n$. For the conjugate f^* of the objective function
 266 $f(x) = c^T x$ we have (cf. Table 1, line 1)

$$267 \quad f^*(y) = 0, \quad \text{dom } f^* = \{c\}.$$

268 Denoting the support function of the set $\{x \mid Ax = b\}$ as $g_1^*(y)$ we have (cf. Table 3, line 1)

$$269 \quad g_1^*(y) = \min_z \{ b^T z \mid A^T z = y \}.$$

270 Denoting the support function of the set constraint $x \in \mathbf{R}_+^n$ as $g_2^*(y)$ we have (cf. Table 3, line 2)

$$271 \quad g_2^*(y) = 0, \quad \text{dom } g_2^* = \{y \mid y \preceq 0\} = -\mathbf{R}_+^n.$$

272 It follows that Fenchel's dual problem is given by

$$273 \quad \sup \left\{ -f^*(y^0) - \sum_{i=1}^2 g_i^*(y^i) \mid \sum_{i=0}^2 y^i = 0 \right\}.$$

274 Substitution of the expressions for f^* and g_i^* ($i = 1, 2$), yields the following formulation of the dual
275 problem:

$$276 \quad \sup_{z, y^i} \{ -b^T z \mid y^0 + y^1 + y^2 = 0, y^0 = c, A^T z = y^1, y^2 \preceq 0 \}.$$

277 We can eliminate the vectors y^i , which gives

$$278 \quad \sup_z \{ -b^T z \mid c + A^T z \succeq 0 \}.$$

279 Changing the sign of z leads to the well-known duality theorem for linear optimization, namely

$$280 \quad \inf_x \{ c^T x \mid Ax = b, x \succeq 0 \} = \sup_z \{ b^T z \mid A^T z \preceq c \}.$$

281 6.2 Conic optimization

282 Let \mathcal{K} be a convex cone in \mathbf{R}^n . We consider the standard linear optimization problem over the cone \mathcal{K} :

$$283 \quad \inf_x \{ c^T x \mid Ax = b, x \in \mathcal{K} \},$$

284 with A , b and c as in Section 6.1. This problem differs from the linear optimization problem in Section
285 6.1 only in the set constraint $x \in \mathcal{K}$, so it suffices to find the support function of this constraint. Denoting
286 this function as g_2^* we have (cf. Table 3, line 3)

$$287 \quad g_2^*(y) = 0, \quad \text{dom } g_2^* = -\mathcal{K}_*.$$

288 An overview of commonly used cones and their dual cones is given in Table 4. Just as in the previous
289 section we obtain that Fenchel's dual problem is given by

$$290 \quad \sup_z \{ -b^T z \mid A^T z + c \in \mathcal{K}_* \}.$$

291 Replacing z by $-z$ we get the usual form of the duality theorem for conic optimization, namely

$$292 \quad \inf_x \{c^T x \mid Ax = b, x \in \mathcal{K}\} = \sup_z \{b^T z \mid c - A^T z \in \mathcal{K}_*\}.$$

293 This formulation is valid for any cone, including the positive semidefinite cone and the copositive cone.
 294 Assumption 3.1 simplifies to the typical Slater condition: there should be an x in the interior of \mathcal{K} that
 295 satisfies $Ax = b$.

296 6.3 Quadratically constrained quadratic optimization

297 Consider the following quadratically constrained quadratic optimization (QCQO) problem:

$$298 \quad \inf_x \left\{ \frac{1}{2}x^T P_0 x + q_0^T x + r_0 \mid \frac{1}{2}x^T P_i x + q_i^T x + r_i \leq 0, \quad i = 1, \dots, m \right\},$$

299 where P_i is positive semidefinite for $i = 0, \dots, m$. Defining $f_i(x) := g_i(x) + h_i(x)$, where $g_i(x) = \frac{1}{2}x^T P_i x$
 300 and $h_i(x) = q_i^T x + r_i$, we have (cf. Table 1, lines 1 and 13)

$$301 \quad \begin{aligned} g_i^*(y) &= \frac{1}{2}y^T P_i^\dagger y, & \text{dom } g_i^* &= \{y \mid y = P_i z, z \in \mathbf{R}^n\}, \\ h_i^*(x) &= -r_i, & \text{dom } h_i^* &= \{q_i\}. \end{aligned}$$

302 Hence, by using the sum rule in Table 2 (line 6), we get

$$303 \quad \begin{aligned} f_i^*(y) &= \min_{y^1, y^2} \{g_i^*(y^1) + h_i^*(y^2) \mid y^1 + y^2 = y\} \\ 304 \quad &= \min_{y^1, y^2} \left\{ \left\{ \frac{1}{2}y^1{}^T P_i^\dagger y^1 \mid y^1 = P_i z \right\} + \{-r_i \mid y^2 = q_i\} \mid y^1 + y^2 = y \right\} \\ 305 \quad &= \min_z \left\{ \frac{1}{2}z^T P_i P_i^\dagger P_i z - r_i \mid P_i z + q_i = y \right\} \\ 306 \quad &= \min_z \left\{ \frac{1}{2}z^T P_i z - r_i \mid P_i z + q_i = y \right\}, \end{aligned}$$

307 where we used that the generalized inverse P_i^\dagger of P_i satisfies $P_i P_i^\dagger P_i = P_i$ (see, e.g., [2]). The dual
 308 problem (D') thus becomes

$$\begin{aligned}
 309 \quad & \sup_{u \succeq 0, y^i} \left\{ -\sum_{i=0}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0, u_0 = 1 \right\} \\
 310 \quad & = \sup_{u \succeq 0, y^i} \left\{ -\sum_{i=0}^m u_i \min_{z^i} \left\{ \frac{1}{2} (z^i)^T P_i z^i - r_i \mid P_i z^i + q_i = \frac{y^i}{u_i} \right\} \mid \sum_{i=0}^m y^i = 0, u_0 = 1 \right\} \\
 311 \quad & = \sup_{u, z^i} \left\{ \sum_{i=0}^m u_i \left(r_i - \frac{1}{2} (z^i)^T P_i z^i \right) \mid \sum_{i=0}^m u_i (P_i z^i + q_i) = 0, u \succeq 0, u_0 = 1 \right\}.
 \end{aligned}$$

312 Defining $\tilde{z}^i = u_i z^i$ this becomes

$$313 \quad \sup_{u, \tilde{z}^i} \left\{ u^T r - \frac{1}{2} \sum_{i=0}^m \frac{(\tilde{z}^i)^T P_i \tilde{z}^i}{u_i} \mid \sum_{i=0}^m u_i q_i + \sum_{i=0}^m P_i \tilde{z}^i = 0, u \succeq 0, u_0 = 1 \right\}.$$

314 It may be noted that by using standard tricks, the dual problem can be written as a conic quadratic
 315 problem.

316 6.4 Second order cone optimization

317 Consider the following primal second order cone optimization (SOCO) problem:

$$318 \quad \min \left\{ \|A_0 x - b_0\|_2 - p_0^T x + q_0 \mid \|A_i x - b_i\|_2 - p_i^T x + q_i \leq 0, i = 1, \dots, m \right\}.$$

319 The objective function and the constraint functions are written as $f_i(x) = g_i(x) + h_i(x)$, for $i = 0, \dots, m$,
 320 where $g_i(x) = \|A_i x - b_i\|_2$, and $h_i(x) = -p_i^T x + q_i$. This enables us to find the conjugate of $f_i(x)$ by
 321 using the sum rule. The conjugate of $h_i(x)$ is given by

$$322 \quad h_i^*(y) = -q_i, \quad \text{dom } h_i^* = \{-p_i\}.$$

323 In order to derive the conjugate of $g_i(x)$, first note that this is a function of a linear transformation of
 324 x , i.e., $g_i(x) = \sigma_i(A_i x - b_i)$, where $\sigma_i(x) = \|x\|_2$. Due to line 6 in Table 1 we have $\sigma_i^*(y) = 0$, with
 325 $\text{dom } \sigma^* = \{y \mid \|y\|_2 \leq 1\}$. Hence, by the linear substitution rule in Table 2 (line 8b), the conjugate of
 326 $g_i(x)$ is given by

$$327 \quad g_i^*(y) = \inf_{z^i} \left\{ \sigma_i^*(z^i) + b_i^T z^i \mid A_i^T z^i = y \right\} = \inf_{z^i} \left\{ b_i^T z^i \mid \|z^i\|_2 \leq 1, A_i^T z^i = y \right\}.$$

328 Putting the above results together, we obtain:

$$\begin{aligned}
329 \quad f_i^*(y) &= \inf_{y^1, y^2} \{g_i^*(y^1) + h_i^*(y^2) \mid y^1 + y^2 = y\} \\
330 \quad &= \inf_{y^1, y^2} \left\{ \inf_{z^i} \{b_i^T z^i \mid \|z^i\|_2 \leq 1, A_i^T z^i = y^1\} + \{-q_i \mid y^2 = -p_i\} \mid y^1 + y^2 = y \right\} \\
331 \quad &= \inf_{z^i} \{b_i^T z^i - q_i \mid \|z^i\|_2 \leq 1, A_i^T z^i - p_i = y\}.
\end{aligned}$$

332 When plugging the above expression for $f_i^*(y)$ into (D') we obtain the following dual for the SOCO
333 problem:

$$334 \quad \sup_{y^i, u \succeq 0} \left\{ -\sum_{i=0}^m u_i \inf_{z^i} \left\{ b_i^T z^i - q_i \mid \|z^i\|_2 \leq 1, A_i^T z^i - p_i = \frac{y^i}{u_i} \right\} \mid \sum_{i=0}^m y^i = 0, u_0 = 1 \right\}.$$

335 We eliminate the variables y^i and omit the inf-operator, which gives

$$336 \quad \sup_{u, z^i} \left\{ \sum_{i=0}^m u_i (q_i - b_i^T z^i) \mid \|z^i\|_2 \leq 1, \forall i, \sum_{i=0}^m u_i (A_i^T z^i - p_i) = 0, u \succeq 0, u_0 = 1 \right\}.$$

337 Introducing $\tilde{z}^i = u_i z^i$ we get

$$338 \quad \sup_{u, \tilde{z}^i} \left\{ u^T q - \sum_{i=0}^m b_i^T \tilde{z}^i \mid \|\tilde{z}^i\|_2 \leq u_i, \forall i, -\sum_{i=0}^m u_i p_i + \sum_{i=0}^m A_i^T \tilde{z}^i = 0, u \succeq 0, u_0 = 1 \right\}.$$

339 Finally, defining matrices A and P according to

$$340 \quad A^T = \begin{bmatrix} A_0^T & \dots & A_m^T \end{bmatrix}, \quad P = \begin{bmatrix} p_0 & \dots & p_m \end{bmatrix},$$

341 and b and \tilde{z} as the vectors that arise by concatenating the vectors b_i and \tilde{z}^i , respectively, the dual problem
342 gets the form

$$343 \quad \sup_{u, \tilde{z}^i} \{u^T q - b^T \tilde{z} \mid \|\tilde{z}^i\|_2 \leq u_i, \forall i, A^T \tilde{z} = Pu, u \succeq 0, u_0 = 1\}.$$

344 This is the known dual for the SOCO problem as written in, e.g., [3].

345 6.5 Geometric optimization

346 Consider the geometric optimization problem (cf. [8, p. 254]):[‡]

$$347 \min_x \left\{ \log \sum_{k=1}^{K_0} e^{a_{0k}^T x + b_{0k}} \mid \log \sum_{k=1}^{K_i} e^{a_{ik}^T x + b_{ik}} \leq 0, \quad i = 1, \dots, m \right\}, \quad (8)$$

348 where $K_i \geq 1$, $a_{ik} \in \mathbf{R}^n$, $b_{ik} \in \mathbf{R}$ for $i = 0, \dots, m$. We define A_i as the $n \times K_i$ matrix with columns
349 a_{i1}, \dots, a_{iK_i} and b_i as the vector with entries b_{i1}, \dots, b_{iK_i} , for $i = 0, \dots, m$. In order to compute the
350 conjugate functions of the objective function and the constraint functions we define, for $i = 0, \dots, m$,

$$351 h_i(x) := \log \sum_{k=1}^{K_i} e^{x_k}$$

$$352 f_i(x) := h_i(A_i^T x + b_i) = \log \sum_{k=1}^{K_i} e^{a_{ik}^T x + b_{ik}}.$$

353 Using the linear substitution rule (Table 2, line 8b) and the formula for the conjugate of $h_i(x)$ (Table 1,
354 line 5), the conjugate of $f_i(x)$ can be written as:

$$355 f_i^*(y) = \inf_z \{ h_i^*(z) - b_i^T z \mid A_i^T z = y \}$$

$$356 = \inf_z \left\{ \sum_{k=1}^{K_i} z_k \log z_k - b_i^T z \mid z \succeq 0, \mathbf{1}^T z = 1, A_i^T z = y \right\}.$$

357 To simplify notation, let p^* be the optimal value of (8). Using the dual problem as given by (D') we
358 obtain

$$359 p^* = \sup_{y^i, u \succeq 0} \left\{ - \sum_{i=0}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0, u_0 = 1 \right\}$$

$$360 = \sup_{y^i, u \succeq 0} \left\{ - \sum_{i=0}^m u_i \inf_{z^i \succeq 0} \left\{ \sum_{k=1}^{K_i} z_k^i \log z_k^i - b_i^T z^i \right\} \right.$$

$$361 \left. \text{s.t. } \sum_{i=0}^m y^i = 0, \mathbf{1}^T z^i = 1, A_i^T z^i = \frac{y^i}{u_i}, \forall i, u_0 = 1 \right\}.$$

[‡]We use e to denote the base of the natural logarithm, according to ISO Standard 80000-2:2009.

362 This is equivalent to

$$\begin{aligned}
363 \quad p^* &= \sup_{y^i, u_{\geq 0}, z^i_{\geq 0}} \left\{ \sum_{i=0}^m u_i \left[b_i^T z^i - \sum_{k=1}^{K_i} z_k^i \log z_k^i \right] \right. \\
364 \quad &\quad \left. \text{s.t. } \sum_{i=0}^m y^i = 0, \mathbf{1}^T z^i = 1, A_i^T z^i = \frac{y^i}{u_i}, \forall i, u_0 = 1 \right\}, \\
365 \quad &= \sup_{u_{\geq 0}, z^i_{\geq 0}} \left\{ \sum_{i=0}^m u_i \left[b_i^T z^i - \sum_{k=1}^{K_i} z_k^i \log z_k^i \right] \mid \sum_{i=0}^m u_i A_i^T z^i = 0, \mathbf{1}^T z^i = 1, \forall i, u_0 = 1 \right\}.
\end{aligned}$$

366 Defining $\tilde{z}^i = u_i z^i$, $i = 0, \dots, m$, we obtain a convex formulation of Fenchel's dual problem:

$$\begin{aligned}
367 \quad p^* &= \sup_{u_{\geq 0}, \tilde{z}^i_{\geq 0}} \left\{ \sum_{i=0}^m \left[b_i^T \tilde{z}^i - \sum_{k=1}^{K_i} \tilde{z}_k^i \log \frac{\tilde{z}_k^i}{u_i} \right] \mid \sum_{i=0}^m A_i^T \tilde{z}^i = 0, \mathbf{1}^T \tilde{z}^i = u_i, \forall i, u_0 = 1 \right\} \\
368 \quad &= \sup_{\tilde{z}^i_{\geq 0}} \left\{ \sum_{i=0}^m \left[b_i^T \tilde{z}^i - \sum_{k=1}^{K_i} \tilde{z}_k^i \log \frac{\tilde{z}_k^i}{\mathbf{1}^T \tilde{z}^i} \right] \mid \sum_{i=0}^m A_i^T \tilde{z}^i = 0, \mathbf{1}^T \tilde{z}^0 = 1 \right\}.
\end{aligned}$$

369 For $m = 0$ this is the same dual problem as the one given in [8, p. 254]. For earlier versions of this dual
370 problem we refer to [9, 10, 11].

371 6.6 ℓ_p -norm optimization

372 We consider the following ℓ_p -norm optimization problem [9, 23]:

$$373 \quad \max_x \left\{ \eta^T x \mid \sum_{i \in I_k} \frac{1}{p_i} |a_i^T x - c_i|^{p_i} + b_k^T x - d_k \leq 0, k = 1, \dots, m \right\}, \quad (9)$$

374 where I_1, \dots, I_m denotes a partition of $\{1, \dots, r\}$, with $1 \leq m \leq r$. Moreover, $b_k \in \mathbf{R}^n$ and $d_k \in \mathbf{R}$ for
375 each k and $a_i \in \mathbf{R}^n$, $c_i \in \mathbf{R}$ and $p_i \geq 1$ for $1 \leq i \leq r$. In order or to obtain Fenchel's dual problem we
376 consider the minimization problem

$$377 \quad \min_x \left\{ -\eta^T x \mid \sum_{i \in I_k} \frac{1}{p_i} |a_i^T x - c_i|^{p_i} + b_k^T x - d_k \leq 0, k = 1, \dots, m \right\}, \quad (10)$$

378 whose optimal value is opposite to the optimal value of (9). We define

$$379 \quad f_0(x) := -\eta^T x$$

$$380 \quad f_k(x) := g_k(x) + h_k(x), \quad k = 1, \dots, m,$$

381 where

$$382 \quad g_k(x) := \sum_{i \in I_k} \frac{1}{p_i} |a_i^T x - c_i|^{p_i}, \quad h_k(x) := b_k^T x - d_k.$$

383 To derive the dual problem of (10), we need the conjugates of f_0, \dots, f_m . The functions f_0 and h_k are
384 linear, so one has (Table 1, line 1)

$$385 \quad f_0^*(y) = 0, \quad \text{dom } f_0^* = \{-\eta\}$$

$$386 \quad h_k^*(y) = d_k, \quad \text{dom } h_k^* = \{b_k\}.$$

387 The function $g_k(x)$ can be written as $g_k(x) = \sum_{i \in I_k} \sigma_i(a_i^T x - c_i)$, where $\sigma_i(x) = |x|^{p_i}/p_i$. Using the sum
388 rule (Table 2, line 6), the linear substitution rule (Table 2, line 8b) and the convex conjugate of $\sigma_i(x)$
389 (Table 1, line 10), we obtain the following conjugate function of $g_k(x)$:

$$390 \quad g_k^*(y) = \inf_{y^i} \left\{ \sum_{i \in I_k} \inf_{z_i \in \mathbf{R}} \left\{ \frac{1}{q_i} |z_i|^{q_i} + c_i z_i \mid a_i z_i = y^i \right\} \mid \sum_{i \in I_k} y^i = y \right\}$$

$$391 \quad = \inf_{z_i} \left\{ \sum_{i \in I_k} \left(\frac{1}{q_i} |z_i|^{q_i} + c_i z_i \right) \mid \sum_{i \in I_k} a_i z_i = y \right\},$$

392 with q_i such that $1/p_i + 1/q_i = 1$. Now the convex conjugate of $f_k(x)$ can be written as:

$$393 \quad f_k^*(y) = \inf_{y^1, y^2} \{ g_k^*(y^1) + h_k^*(y^2) \mid y^1 + y^2 = y \}$$

$$394 \quad = \inf_{z_i} \left\{ \left(\sum_{i \in I_k} \frac{1}{q_i} |z_i|^{q_i} + c_i z_i \right) + d_k \mid \sum_{i \in I_k} a_i z_i + b_k = y \right\}.$$

395 Fenchel's dual of (10) is now given by

$$396 \quad \max_{y^k, u \geq 0} \left\{ -f_0^*(y^0) - \sum_{k=1}^m (u_k f_k)^*(y^k) \mid \sum_{k=0}^m y^k = 0 \right\}.$$

397 Using the formulas that we derived above, we get the problem

$$398 \quad \max_{y^k, u \geq 0} \left\{ - \sum_{k=1}^m u_k \inf_{z_i} \left\{ \sum_{i \in I_k} \left(\frac{1}{q_i} |z_i|^{q_i} + c_i z_i \right) + d_k \mid y^0 = -\eta, \sum_{i \in I_k} a_i z_i + b_k = \frac{y^k}{u_k} \right\} \right\},$$

399 subject to the condition $\sum_{k=0}^m y^k = 0$. This can be simplified to

$$400 \quad \max_{u \geq 0, z_i} \left\{ - \sum_{k=1}^m u_k \left[\sum_{i \in I_k} \left(\frac{1}{q_i} |z_i|^{q_i} + c_i z_i \right) + d_k \right] \mid \sum_{k=1}^m u_k \left(\sum_{i \in I_k} a_i z_i + b_k \right) = \eta \right\}.$$

401 To further simplify this formulation we define $\tilde{z}_i = u_k z_i$ for each i . Then we get

$$402 \quad \max_{u \geq 0, \tilde{z}_i} \left\{ - \sum_{k=1}^m u_k \left[\sum_{i \in I_k} \left(\frac{1}{q_i} \left| \frac{\tilde{z}_i}{u_k} \right|^{q_i} + c_i \frac{\tilde{z}_i}{u_k} \right) + d_k \right] \mid \sum_{k=1}^m u_k \left(\sum_{i \in I_k} a_i \frac{\tilde{z}_i}{u_k} + b_k \right) = \eta \right\}.$$

403 By changing the sign of the optimal value we arrive at

$$404 \quad \min_{u \geq 0, \tilde{z}_i} \left\{ c^T \tilde{z} + u^T d + \sum_{k=1}^m u_k \sum_{i \in I_k} \frac{1}{q_i} \left| \frac{\tilde{z}_i}{u_k} \right|^{q_i} \mid \sum_{k=1}^m \sum_{i \in I_k} a_i \tilde{z}_i + \sum_{k=1}^m u_k b_k = \eta \right\},$$

405 which is the dual problem of (9) that we are looking for. Finally, defining matrices A and B according to

$$406 \quad A = \begin{bmatrix} a_1 & \dots & a_n \end{bmatrix}, \quad B = \begin{bmatrix} b_1 & \dots & b_r \end{bmatrix},$$

407 the last problem gets the form

$$408 \quad \min_{u \geq 0, \tilde{z}_i} \left\{ u^T d + c^T \tilde{z} + \sum_{k=1}^m u_k \sum_{i \in I_k} \frac{1}{q_i} \left| \frac{\tilde{z}_i}{u_k} \right|^{q_i} \mid A \tilde{z} + B u = \eta \right\},$$

409 which is the same dual as was obtained in [23].

410 6.7 Radiotherapy treatment planning

411 Optimization plays an important role in radiotherapy treatment planning. The resulting optimization
 412 problems contain many different nonlinear objective and constraint functions. To the best of our knowl-
 413 edge the corresponding dual problems have not been derived. These dual problems might be useful for
 414 reasons already stated in the Introduction. In this section we show that with the framework proposed in
 415 this paper, the dual problem can be derived in a structured and straightforward way.

416 Treatment planning for external beam radiotherapy aims at finding beam intensities such that the tumor
 417 is controlled, while limiting the dose to the surrounding organs at risk (OARs). For planning purposes,
 418 the structures of interest (e.g. tumor, lung, heart) are virtually divided into a large number of small cubic
 419 volumes, so-called voxels. An important input parameter to all treatment planning models is the matrix
 420 D , whose (i, j) -element is the dose per unit intensity (dose rate) from *beamlet* j to voxel i (each beam is
 421 subdivided into beamlets). The dose from beamlet j to voxel i is equal to $D_{ij} x_j$ with x_j the intensity
 422 of beamlet j . The total dose to voxel i can be written as $d_i = D_i x$, where D_i is the i^{th} row of D and x
 423 the vector of beamlet intensities. Treatment plans are evaluated based on the dose to each of the voxels.

424 Treatment planning models employ constraints and objectives such as the minimum, mean or maximum
 425 dose over all voxels within a structure, dose-volume metrics such as the minimum dose to the hottest
 426 $p\%$ of the voxels, or biological metrics such as the tumor control probability (TCP). Treatment planning
 427 models may thus contain different types of constraints and do usually not fit a generic optimization format
 428 (as, e.g., linear or second order cone optimization). A readily given dual problem is thus generally not
 429 available. Here, we give an example of such a treatment planning model and derive its dual using the
 430 method proposed in this paper.

431 The objective of our model is to maximize the TCP [22]:

$$432 \quad TCP(x) = \prod_{i \in \mathcal{T}} \exp(-N_0 v_i e^{-\alpha D_i x}),$$

433 where \mathcal{T} is the set of voxels in the target volume, N_0 is the number of clonogenic cells, v_i is the relative
 434 volume of voxel i and α describes the radio resistance of cells in the target volume. This function is not
 435 concave. Taking the log of TCP gives the concave function LTCP:

$$436 \quad LTCP(x) = - \sum_{i \in \mathcal{T}} N_0 v_i e^{-\alpha D_i x},$$

437 that yields the same optimal solution for x as $\log(x)$ is increasing [20]. Therefore, LTCP is optimized
 438 instead.

439 Our planning model contains a constraint on the dose to the OARs. For simplicity, we assume there is
 440 only one OAR, denoted as S . First, we consider the generalized equivalent uniform dose (gEUD), which
 441 is the generalized mean of the dose to a structure S [18]:

$$442 \quad gEUD(x) = \left(\sum_{i \in S} v_i (D_i x)^p \right)^{\frac{1}{p}},$$

443 where S denotes the set of voxels in structure S , and $p \geq 1$ is a structure-dependent parameter for OARs
 444 that indicates if it is a serial or a parallel organ. We limit the normal tissue complication probability
 445 (NTCP) to this OAR [1]:

$$446 \quad NTCP(x) = 1 - \exp \left(- \left(\frac{gEUD(x)}{\Delta} \right)^p \right),$$

447 where Δ is a structure-dependent parameter that indicates the uniform dose level for which there is a
 448 63% probability on complications. Furthermore, $p > 1$ since we restrict ourselves to a non-parallel organ.
 449 We would like to constrain NTCP from above, however, NTCP is not a convex function. Therefore,

450 following [20], we use instead

$$451 \quad -\ln(1 - NTCP(x)) = \frac{1}{\Delta^p} \sum_{i \in \mathcal{S}} v_i (D_i x)^p,$$

452 which is convex. As was the case with the transformation from TCP to LTCP, this does not change the
 453 optimal solution. Note that we need to apply the same transformation to the upper bound on NTCP.

454 We can now formulate our treatment plan optimization model:

$$455 \quad \min_x \left\{ N_0 \sum_{i \in \mathcal{T}} v_i e^{-\alpha D_i x} \mid \begin{cases} \frac{1}{\Delta^p} \sum_{i \in \mathcal{S}} v_i (D_i x)^p - \gamma \leq 0 \\ x \succeq 0 \end{cases} \right\}.$$

456 where γ is a predetermined upper bound on $-\ln(1 - NTCP)$. The non-negativity constraint is included
 457 since a negative beam intensity is physically impossible.

458 Before giving the dual of this problem, we derive the conjugates of the objective function and the con-
 459 straint function.

460 Objective function

461 The objective function $f_0(x)$ is a constant multiplied by the sum of $v_i h_i(x)$, where $h_i(x) = g(-\alpha D_i x)$,
 462 $g(u) = e^u$. Using Table 1 (line 2) and the linear substitution rule (Table 2, line 8b), we obtain the convex
 463 conjugate of $h_i(x)$:

$$464 \quad h_i^*(y) = \inf_{z_i} \{ z_i \log z_i - z_i \mid -\alpha z_i D_i^T = y, z_i \geq 0 \}.$$

465 Application of the sum rule for conjugates (Table 2, line 6) now enables us to write the conjugate of $f_0(x)$

466 as follows:

$$\begin{aligned}
467 \quad f_0^*(y) &= N_0 \left(\sum_{i \in \mathcal{T}} v_i h_i \right)^* \left(\frac{y}{N_0} \right) && \text{(Product with a scalar)} \\
468 \quad &= N_0 \inf_{\{y^i\}_{i \in \mathcal{T}}} \left\{ \sum_{i \in \mathcal{T}} (v_i h_i)^*(y^i) \mid \sum_{i \in \mathcal{T}} y^i = \frac{y}{N_0} \right\} && \text{(Sum rule)} \\
469 \quad &= N_0 \inf_{\{y^i\}_{i \in \mathcal{T}}} \left\{ \sum_{i \in \mathcal{T}} v_i h_i^* \left(\frac{y^i}{v_i} \right) \mid \sum_{i \in \mathcal{T}} y^i = \frac{y}{N_0} \right\} && \text{(Product with a scalar)} \\
470 \quad &= N_0 \inf_{\{y^i\}_{i \in \mathcal{T}}} \left\{ \sum_{i \in \mathcal{T}} v_i \inf_{z_i} \left\{ z_i \log z_i - z_i \mid -\alpha z_i D_i^T = \frac{y^i}{v_i}, z_i \geq 0 \right\} \mid \sum_{i \in \mathcal{T}} y^i = \frac{y}{N_0} \right\} \\
471 \quad &= N_0 \inf_w \left\{ \sum_{i \in \mathcal{T}} w_i \log \frac{w_i}{v_i} - w_i \mid \sum_{i \in \mathcal{T}} -\alpha w_i D_i^T = \frac{y}{N_0}, w \succeq 0 \right\} && (w_i := v_i z_i) \\
472 \quad &= N_0 \inf_w \left\{ \sum_{i \in \mathcal{T}} w_i \log \frac{w_i}{v_i} - w_i \mid -\alpha N_0 D_{\mathcal{T}}^T w = y, w \succeq 0 \right\},
\end{aligned}$$

473 where $D_{\mathcal{T}}$ consists of the rows in D related to the voxels in \mathcal{T} .

474 NTCP constraint

475 The *NTCP* constraint function $f_1(x)$ is given by

$$476 \quad f_1(x) = \frac{1}{\Delta^p} \sum_{i \in \mathcal{S}} v_i (D_i x)^p - \gamma = \frac{1}{\Delta^p} \sum_{i \in \mathcal{S}} v_i g_i(x) - \gamma,$$

477 where $g_i(x) = h_i(D_i x)$, $h_i(u) = u^p$ ($u \geq 0$), with $p > 1$. The convex conjugates of $h_i(u)$ and $g_i(x)$ are
478 (using Table 1, line 8 and Table 2, lines 4 and 8b),

$$\begin{aligned}
479 \quad h_i^*(y) &= \inf_{t_i} \left\{ \frac{p}{q} t_i^q \mid t_i \geq 0, t_i \geq \frac{y}{p} \right\} \\
480 \quad g_i^*(y) &= \inf_{r_i} \left\{ \inf_{t_i} \left\{ \frac{p}{q} t_i^q \mid t_i \geq 0, t_i \geq \frac{r_i}{p} \right\} \mid D_i^T r_i = y \right\} \\
481 \quad &= \inf_{r_i, t_i} \left\{ \frac{p}{q} t_i^q \mid D_i^T r_i = y, t_i \geq 0, t_i \geq \frac{r_i}{p} \right\}.
\end{aligned}$$

482 By using the sum rule for conjugates, the multiply-with-a-constant rule and add-a-constant rule (Table 2,

483 lines 6, 4 and 1, respectively) we can now compute $f_2^*(y)$:

$$\begin{aligned}
484 \quad f_2^*(y) &= \frac{1}{\Delta^p} \inf_{y^i} \left\{ \sum_{i \in \mathcal{S}} v_i g_i^* \left(\frac{y^i}{v_i} \right) \mid \sum_{i \in \mathcal{S}} y^i = \Delta^p y \right\} + \gamma \\
485 &= \frac{1}{\Delta^p} \inf_{y^i} \left\{ \sum_{i \in \mathcal{S}} v_i \inf_{r_i, t_i} \left\{ \frac{p}{q} t_i^q \mid D_i^T r_i = \frac{y^i}{v_i}, t_i \geq 0, t_i \geq \frac{r_i}{p} \right\} \mid \sum_{i \in \mathcal{S}} y^i = \Delta^p y \right\} + \gamma \\
486 &= \frac{p}{q \Delta^p} \inf_{r_i, t_i} \left\{ \sum_{i \in \mathcal{S}} v_i t_i^q \mid \sum_{i \in \mathcal{S}} v_i D_i^T r_i = \Delta^p y, t_i \geq 0, p t_i \geq r_i \right\} + \gamma \\
487 &= \inf_{r, t} \left\{ \gamma + \frac{p}{q \Delta^p} v_S^T t^q \mid D_S^T V_S r = \Delta^p y, t \succeq 0, p t \succeq r \right\},
\end{aligned}$$

488 where t^q is the vector with entries t_i^q , $i \in \mathcal{S}$.

489 Non-negativity constraint

490 We finally have to deal with the non-negativity constraint $x \succeq 0$. Denoting the indicator function of \mathbf{R}_+^n
491 as $f_2(x)$, the corresponding support function is (cf. Table 3, line 2)

$$492 \quad f_1^*(y) = 0, \quad \text{dom } f_1^* = \{y \mid y \succeq 0\}.$$

493 As we know, Fenchel's dual problem is given by

$$494 \quad \sup_{\{y^i\}_{i=0}^2, u} \left\{ -f_0^*(y^0) - u f_1^* \left(\frac{y^1}{u} \right) - f_2^*(y^2) \mid \sum_{k=0}^2 y^k = 0, u \geq 0 \right\}.$$

495 We can already eliminate y^2 , yielding the equivalent problem

$$496 \quad \sup_{y^0, y^1, u} \left\{ -f_0^*(y^0) - u f_1^* \left(\frac{y^1}{u} \right) \mid y^0 + y^1 \succeq 0, u \geq 0 \right\}.$$

497 Thus we obtain the following dual for the treatment planning optimization problem:

$$\begin{aligned}
498 \quad & \sup_{y^0, y^1, u} \left\{ -N_0 \left\{ \inf_w \sum_{i \in \mathcal{T}} \left(w_i \log \frac{w_i}{v_i} - w_i \right) \mid -\alpha N_0 D_{\mathcal{T}}^T w = y^0, w \succeq 0 \right\} \right. \\
499 & \quad \left. - u \inf_{r, t} \left\{ \gamma + \frac{p}{q \Delta^p} v_S^T t^q \mid D_S^T V_S r = \Delta^p \frac{y^1}{u}, t \succeq 0, p t \succeq r \right\} \right. \\
500 & \quad \left. \mid y^0 + y^1 \succeq 0, u \geq 0 \right\}.
\end{aligned}$$

501 By omitting the inf operators we get

$$\begin{aligned}
502 \quad & \sup_{y^0, y^1, w, z, r, t, s, u} \left\{ -N_0 \left\{ \sum_{i \in \mathcal{T}} \left(w_i \log \frac{w_i}{v_i} - w_i \right) \mid -\alpha N_0 D_{\mathcal{T}}^T w = y^0, w \succeq 0 \right\} \right. \\
503 \quad & \left. -u \left\{ \gamma + \frac{p}{q \Delta^p} v_S^T t^q \mid D_S^T V_S r = \Delta^p \frac{y^1}{u}, t \succeq 0, pt \succeq r \right\} \right. \\
504 \quad & \left. \mid y^0 + y^1 \succeq 0, u \geq 0 \right\}.
\end{aligned}$$

505 This is equivalent to

$$\begin{aligned}
506 \quad & \sup_{w, r, t, u} \left\{ -N_0 \sum_{i \in \mathcal{T}} \left(w_i \log \frac{w_i}{v_i} - w_i \right) - u \gamma - \frac{u p v_S^T t^q}{q \Delta^p} \mid \right. \\
507 \quad & \left. w \succeq 0, t \succeq 0, pt \succeq r, u \geq 0, \right. \\
508 \quad & \left. -\alpha N_0 D_{\mathcal{T}}^T w + \frac{1}{\Delta^p} D_S^T V_S(ur) \succeq 0 \right\}.
\end{aligned}$$

509 By redefining s according to $r := ur, t := ut$ this becomes

$$\begin{aligned}
510 \quad & \sup_{w, r, t, u} \left\{ -N_0 \sum_{i \in \mathcal{T}} \left(w_i \log \frac{w_i}{v_i} - w_i \right) - u \gamma - \frac{u p}{q \Delta^p} v_S^T \left(\frac{t}{u} \right)^q \mid \right. \\
511 \quad & \left. w \succeq 0, t \succeq 0, pt \succeq r, u \geq 0, \right. \\
512 \quad & \left. -\alpha N_0 D_{\mathcal{T}}^T w + \frac{1}{\Delta^p} D_S^T V_S r + s \succeq 0 \right\}.
\end{aligned}$$

513 One easily verifies that this dual problem of the treatment plan optimization problem is indeed a convex
514 optimization problem.

515 7 Concluding remarks

516 Mathematical education usually includes the use of tables, e.g., tables of derivatives, primitive functions,
517 Laplace transforms, Fourier transforms, etc. It may have become clear from this paper that the availability
518 of tables as presented in in Appendix E simplifies the process of dualizing an optimization problem.
519 Combined with the use of Fenchel's duality theorem it yields a universal and structured way for deriving
520 dual problems for a wide spectrum of convex optimization problems.

521 Acknowledgement

522 The authors thank Lieven Vandenberghe (UCLA) for useful suggestions.

523 A Some examples of conjugate functions

524 By way of example we demonstrate how some of the results in Table 1 are obtained; this table also gives
525 appropriate references to the existing literature whenever available. The examples in this section are used
526 in Section 6 where we compute the dual problems of some more or less standard optimization problems.

527 A.1 Linear function

528 We start by proving the formula in line 1 of Table 1. With $f(x) = a^T x + b$, $a \in \mathbf{R}^n$ and $b \in \mathbf{R}$, we have

$$529 \quad f^*(y) = \sup_x \{y^T x - (a^T x + b)\} = \sup_x \{(y - a)^T x - b\}.$$

530 Obviously the sup equals $-b$ if $y = a$ and otherwise infinity (take $x = \lambda(y - a)$, $\lambda > 0$). Hence we obtain

$$531 \quad f^*(y) = -b, \quad \text{dom } f^* = \{a\}.$$

532 A.2 Log-sum-exp function

533 We next deal with line 5 in Table 1. Let $f(x) = \log(\sum_{i=1}^n e^{x_i})$, $x \in \mathbf{R}^n$. We then have

$$534 \quad f^*(y) = \sup_x \left\{ y^T x - \log \left(\sum_{i=1}^n e^{x_i} \right) \right\}.$$

535 Define $g(x) = y^T x - f(x)$. If $y_i < 0$ for some i , we take $x = -\lambda e_i$, where $\lambda > 0$ and e_i is the i -th unit
536 vector. Then $g(x) = -\lambda y_i - \log(n - 1 + e^{-\lambda})$, which goes to infinity if λ goes to infinity. If $y \succeq 0$ but
537 $\mathbf{1}^T y \neq 1$, take $x = \lambda \mathbf{1}$. Then $g(x) = \lambda \mathbf{1}^T y - \log(n e^\lambda) = \lambda(\mathbf{1}^T y - 1) - \log n$, which goes to infinity both
538 if $\mathbf{1}^T y > 1$ (when λ goes to infinity) and if $\mathbf{1}^T y < 1$ (when λ goes to minus infinity). Hence $y \in \text{dom } f^*$
539 holds only if $\mathbf{1}^T y = 1$ and $y \succeq 0$. We proceed by computing the maximal value of $g(x)$ under these
540 conditions. One has

$$541 \quad \frac{\partial g(x)}{\partial x_i} = y_i - \frac{e^{x_i}}{\sum_{i=1}^n e^{x_i}}.$$

542 By setting these partial derivatives equal to zero we obtain the condition

$$543 \quad y_i = \frac{e^{x_i}}{\sum_{i=1}^n e^{x_i}}, \quad i = 1, \dots, n.$$

544 Substitution of these values of y into $g(x)$ yields

$$545 \quad f^*(y) = \sum_{i=1}^n y_i \log y_i, \quad y \succ 0, \mathbf{1}^T y = 1.$$

546 By interpreting $0 \log 0 = 0$ this expression remains valid if some entries of y vanish. Thus we conclude
547 that

$$548 \quad f^*(y) = \sum_{i=1}^n y_i \log y_i, \quad \text{dom } f^* = \{y \mid \mathbf{1}^T y = 1, y \succeq 0\}.$$

549 A.3 Arbitrary norm

550 In this section we deal with line 6 in Table 1. Let $\|x\|$ be an arbitrary norm on \mathbf{R}^n and $f(x) := \|x\|$. We
551 then have

$$552 \quad f^*(y) = \sup_x \{y^T x - \|x\|\}.$$

553 Recall that the dual norm is defined by

$$554 \quad \|y\|_* = \sup_x \{y^T x \mid \|x\| \leq 1\}.$$

555 Since $\left\| \frac{x}{\|x\|} \right\| = 1$ if $x \neq 0$, we have for any nonzero $x \in \mathbf{R}^n$

$$556 \quad \|y\|_* \geq y^T \frac{x}{\|x\|} = \frac{y^T x}{\|x\|},$$

557 which gives

$$558 \quad y^T x \leq \|y\|_* \|x\|. \tag{11}$$

559 Since this inequality holds also for $x = 0$, it holds for each $x \in \mathbf{R}^n$. Define $g(x) := y^T x - \|x\|$. Then
560 $g(\lambda x) = \lambda g(x)$ for $\lambda \geq 0$.

561 If $\|y\|_* > 1$ then there exists an x with $\|x\| \leq 1$ and $y^T x > 1$. Then $g(x) > 0$. Hence $g(\lambda x)$ goes to
562 infinity if λ grows to infinity. We conclude from this that $y \in \text{dom } f^*$ holds only if $\|y\|_* \leq 1$.

563 If $\|y\|_* \leq 1$ then (11) implies $y^T x \leq \|y\|_* \|x\| \leq \|x\|$, whence $g(x) \leq 0$, for all $x \in \mathbf{R}^n$. Since $g(0) = 0$,
 564 we obtain

$$565 \quad f^*(y) = 0, \quad \text{dom } f^* = \{y \mid \|y\|_* \leq 1\}.$$

566 Let us recall that if $\|\cdot\|$ represents the p -norm then it is denoted as $\|\cdot\|_p$ and

$$567 \quad \|x\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{\frac{1}{p}}, \quad p \geq 1, \quad x \in \mathbf{R}^n.$$

568 In that case the dual norm is given by $\|\cdot\|_q$, with q such that $\frac{1}{p} + \frac{1}{q} = 1$.

569 **A.4 Quadratic function**

570 In this section we deal with line 13 in Table 1. Let $f(x) = \frac{1}{2}x^T P x, x \in \mathbf{R}^n$, where P is a (symmetric)
 571 positive semidefinite $n \times n$ matrix. We then have

$$572 \quad f^*(y) = \sup_x \{y^T x - \frac{1}{2}x^T P x\}.$$

573 Since $y^T x - \frac{1}{2}x^T P x$ is concave in x , its maximal value occurs if $y = P x$. If P is nonsingular then this
 574 implies $x = P^{-1}y$ and hence

$$575 \quad f^*(y) = y^T P^{-1}y - \frac{1}{2} (P^{-1}y)^T P P^{-1}y = \frac{1}{2}y^T P^{-1}y.$$

576 If P is singular, the equation $y = P x$ has a solution if and only if y belongs to the column space
 577 $L = \{P z \mid z \in \mathbf{R}^n\}$. If $y \in L$ then the solution is given by $x = P^\dagger y$, where P^\dagger is the generalized inverse
 578 of P , and we obtain

$$579 \quad f^*(y) = y^T P^\dagger y - \frac{1}{2} (P^\dagger y)^T P P^\dagger y = \frac{1}{2}y^T P^\dagger y,$$

580 where the last equality is due to the fact that $P^\dagger P P^\dagger = P^\dagger$ [2].

581 If $y \notin L$ we may write $y = y_1 + y_2$ with $y_1 \in L$ and $0 \neq y_2 \in L^\perp$. Take $\lambda > 0$ and $x = \lambda y_2$. Then
 582 $x^T y = x^T y_1 + x^T y_2 = x^T y_2 = \lambda y_2^T y_2 = \lambda \|y_2\|^2$. Since $x \in L^\perp$, we have $P x = 0$, whence $x^T P x = 0$.
 583 Hence $y^T x - \frac{1}{2}x^T P x = \lambda \|y_2\|^2$, which goes to infinity if λ grows to infinity. This proves that $y \notin \text{dom } f^*$

584 if $y \notin L$. Thus we have shown that

$$585 \quad f^*(y) = \frac{1}{2}y^T P^\dagger y, \quad \text{dom } f^* = \{y \mid y = Pz, z \in \mathbf{R}^n\}.$$

586 A.5 Conjugate of $\frac{1}{p}x^p$, $x \geq 0$, $p > 1$

587 We prove the formula in line 12 of Table 1. Let $f(x) = \frac{1}{p}x^p$, $x \geq 0$, $p > 1$. One has

$$588 \quad f'(x) = x^{p-1}$$

$$589 \quad f''(x) = (p-1)x^{p-2},$$

590 which shows that f is convex, because $x \geq 0$ and $p > 1$. Then

$$591 \quad f^*(y) = \sup_{x \geq 0} (yx - f(x)).$$

592 Since $yx - f(x) \leq 0$ if $y < 0$ the sup value the occurs for $x = 0$ and then it equals $0 - f(0) = 0$. Otherwise,

593 if $y \geq 0$, we have

$$594 \quad f^*(y) = \sup_{x \geq 0} \left(yx - \frac{1}{p}x^p \right).$$

595 The sup value is then attained when $y = f'(x) = x^{p-1}$, which gives $x = y^{\frac{1}{p-1}}$, and

$$596 \quad f^*(y) = y^{1+\frac{1}{p-1}} - \frac{1}{p}y^{\frac{p}{p-1}} = \left(1 - \frac{1}{p}\right)y^{\frac{p}{p-1}} = \frac{1}{q}y^q,$$

597 where q is such that $\frac{1}{p} + \frac{1}{q} = 1$. Summarizing,

$$598 \quad f^*(y) = \begin{cases} 0, & y < 0, \\ \frac{1}{q}y^q, & y \geq 0. \end{cases}$$

599 Since $f^*(y)$ is monotonically increasing, this can be written as

$$600 \quad f^*(y) = \inf_z \left\{ \frac{1}{q}z^q \mid z \geq y, z \geq 0 \right\}.$$

601 **A.6 Conjugate of $\frac{1}{px^p}$, $x > 0$, $p > 0$**

602 We prove the formula in line 11 of Table 1. Let $f(x) = \frac{1}{px^p}$, $x > 0$, $p > 0$. One has

$$603 \quad f'(x) = \frac{-1}{x^{p+1}}$$

$$604 \quad f''(x) = \frac{p+1}{x^{p+2}},$$

605 which shows that f is convex, because $x > 0$ and $p+1 > 1 > 0$. Also note that $f(x)$ is monotonically
606 decreasing to zero if x goes to ∞ . We have

$$607 \quad f^*(y) = \sup_{x>0} (yx - f(x)).$$

608 If $y > 0$ then $yx - f(x)$ goes to ∞ if x goes to ∞ . Hence the sup value is ∞ if $y > 0$. So, we may assume
609 that $y \leq 0$. Since then

$$610 \quad f^*(y) = \sup_{x>0} \left(yx - \frac{1}{px^p} \right),$$

611 the sup value is attained when $y = f'(x) = \frac{-1}{x^{p+1}}$, which gives $x = (-y)^{\frac{-1}{p+1}}$, and

$$612 \quad f^*(y) = y(-y)^{\frac{-1}{p+1}} - \frac{1}{p}(-y)^{\frac{p}{p+1}} = \left(-1 - \frac{1}{p}\right)(-y)^{\frac{p}{p+1}} = -\left(\frac{p+1}{p}\right)(-y)^{\frac{p}{p+1}}.$$

613 Summarizing,

$$614 \quad f^*(y) = -\left(\frac{p+1}{p}\right)(-y)^{\frac{p}{p+1}}, \quad \text{dom } f^* = \{y \mid y \leq 0\} = -\mathbf{R}_+.$$

615 **B Some identities for conjugate functions**

616 **B.1 Conjugate of $f(x) = \max_i f_i(x)$**

617 First note that we may write

$$618 \quad f(x) = \max_z \left\{ \sum_{i=1}^m z_i f_i(x) \mid \sum_{i=1}^m z_i = 1, z \succeq 0 \right\} = \max_{z \in S} \left\{ \sum_{i=1}^m z_i f_i(x) \right\},$$

619 where S denotes the simplex. By the definition of f^* we have

$$\begin{aligned}
620 \quad f^*(y) &= \sup_x \left\{ y^T x - \max_{z \in S} \left\{ \sum_{i=1}^m z_i f_i(x) \right\} \right\} \\
621 \quad &= \sup_x \min_{z \in S} \left\{ y^T x - \sum_{i=1}^m z_i f_i(x) \right\}.
\end{aligned}$$

622 Since the argument is convex in z and concave in x , and the domain of z is compact, we may interchange
623 the sup and the min operators, yielding

$$624 \quad f^*(y) = \min_{z \in S} \sup_x \left\{ y^T x - \sum_{i=1}^m z_i f_i(x) \right\} = \min_{z \in S} \left(\sum_{i=1}^m z_i f_i \right)^*(y).$$

625 Using the sum rule for conjugates (Table 2, line 6) we obtain

$$\begin{aligned}
626 \quad f^*(y) &= \min_{z \in S} \left\{ \sum_{i=1}^m (z_i f_i)^*(y^i) \mid \sum_{i=1}^m y^i = y \right\} \\
627 \quad &= \min_z \left\{ \sum_{i=1}^m (z_i f_i)^*(y^i) \mid \sum_{i=1}^m y^i = y, \sum_{i=1}^m z_i = 1, z \succeq 0 \right\}.
\end{aligned}$$

628 B.2 Deriving conjugates via the adjoint I

629 In this section we prove the formula in line 10 of Table 2. This formula enables us to write $f^*(y)$ as an
630 inf expression in cases where no closed formula is available for $f^*(y)$, whereas such a formula exists for
631 $(f^\diamond)^*(y)$. The proof below is an alternative for that in [14].

632 By the definition of $(f^\diamond)^*(y)$ we have

$$633 \quad \inf_z \{ z \mid (f^\diamond)^*(-z) \leq -y \} = \inf_z \left\{ z \mid \sup_x \{ -zx - f^\diamond(x) \} \leq -y \right\}.$$

634 Since $\text{dom } f^\diamond = \mathbf{R}_{++}$ we have $f^\diamond(x) = \infty$ if $x \leq 0$. Hence

$$635 \quad \inf_z \{ z \mid (f^\diamond)^*(-z) \leq -y \} = \inf_z \left\{ z \mid \sup_{x>0} \left\{ -zx - x f \left(\frac{1}{x} \right) \right\} \leq -y \right\}.$$

636 Due to Lagrange duality this implies

$$\begin{aligned}
637 \quad \inf_z \{ z \mid (f^\diamond)^*(-z) \leq -y \} &= \sup_{\lambda \geq 0} \inf_z \left\{ z + \lambda \left[y + \sup_{x>0} \left\{ -zx - x f \left(\frac{1}{x} \right) \right\} \right] \right\} \\
638 \quad &= \sup_{\lambda \geq 0} \inf_z \sup_{x>0} \left\{ \lambda y - \lambda x f \left(\frac{1}{x} \right) + z(1 - \lambda x) \right\}. \tag{12}
\end{aligned}$$

639 Observe that the argument in the last expression is the Lagrangian of the problem

$$640 \quad \sup_{x>0} \left\{ \lambda y - \lambda x f \left(\frac{1}{x} \right) \mid 1 - \lambda x = 0 \right\}.$$

641 Since this problem is Slater regular, its Lagrangian has a saddle point. As a consequence, in (12) we may

642 replace $\inf_z \sup_{x>0}$ by $\sup_{x>0} \inf_z$. One has

$$643 \quad \sup_{x>0} \inf_z \left\{ \lambda y - \lambda x f \left(\frac{1}{x} \right) + z(1 - \lambda x) \right\} = \sup_{x>0} \left\{ \lambda y - \lambda x f \left(\frac{1}{x} \right) \mid \lambda x = 1 \right\} = \lambda y - f(\lambda).$$

644 Substitution into (12) yields

$$645 \quad \inf_z \{ z \mid (f^\diamond)^*(-z) \leq -y \} = \sup_{\lambda \geq 0} \{ \lambda y - f(\lambda) \} = f^*(y).$$

646 The last equality holds because the domain of f is \mathbf{R}_+ .

647 **B.3 Deriving conjugates via the adjoint II**

648 In this section we prove formula 11 in Table 2. This formula can be applied if f has an inverse, which is

649 denoted as $h = f^{-1}$.

650 Since h is concave we have, by definition, $h_*(y) = \inf_x \{ yx - h(x) \}$. Thus we may write

$$651 \quad \begin{aligned} -(h_*)^\diamond(y) &= -yh_* \left(\frac{1}{y} \right) = -y \inf_x \left\{ \frac{1}{y}x - h(x) \right\} = y \sup_x \left\{ h(x) - \frac{x}{y} \right\} \\ 652 \quad &= y \sup_{x,t} \left\{ t - \frac{x}{y} \mid t \leq h(x) \right\}, \end{aligned}$$

653 where we used that h is (strictly) increasing. If $y < 0$ then $t - \frac{x}{y}$ goes to infinity if x goes to infinity.

654 Hence, we may assume that $y \geq 0$. Also using that $t \leq h(x)$ holds if and only if $f(t) \leq x$ we may proceed

655 as follows:

$$656 \quad \begin{aligned} -(h_*)^\diamond(y) &= \sup_{x,t} \{ ty - x \mid t \leq h(x) \} = \sup_{x,t} \{ ty - x \mid f(t) \leq x \} \\ 657 \quad &= \sup_t \{ ty - f(t) \} = f^*(y). \end{aligned}$$

658 **B.4 Linear substitution rule**

659 In this section we prove the formulas in lines 8a and 8b of Table 2. Let f be defined by $f(x) = h(Ax + b)$,
 660 where A is an $m \times n$ matrix, $b \in \mathbf{R}^m$ and $h : \mathbf{R}^m \rightarrow \mathbf{R}$ convex. We then have

$$\begin{aligned}
 661 \quad f^*(y) &= \sup_x \{y^T x - h(Ax + b)\} \\
 662 \quad &= \sup_{x,v} \{y^T x - h(v) \mid Ax + b = v\}.
 \end{aligned}$$

663 By writing the last problem as a minimization problem and then using Lagrange's duality theorem we
 664 get

$$\begin{aligned}
 665 \quad f^*(y) &= \min_z \sup_{x,v} \{y^T x - h(v) + z^T(v - b - Ax)\} \\
 666 \quad &= \min_z \sup_{x,v} \{z^T v - h(v) - z^T b + x^T(y - A^T z)\}.
 \end{aligned}$$

667 The contribution of the term containing x is finite if and only if $y - A^T z = 0$. Thus we obtain

$$\begin{aligned}
 668 \quad f^*(y) &= \min_z \sup_v \{z^T v - h(v) - z^T b \mid A^T z = y\} \\
 669 \quad &= \min_z \{h^*(z) - b^T z \mid A^T z = y\}.
 \end{aligned}$$

670 If A is square and nonsingular then $A^T z = y$ holds if and only if $z = A^{-T}y$. Hence, in that case we have

$$671 \quad f^*(y) = h^*(A^{-T}y) - b^T A^{-T}y.$$

672 **B.5 Composite function**

673 In this section we prove the formula in line 9 of Table 2. Let $f(x) = g(h(x))$, with g and h convex and g
 674 nondecreasing. We may write

$$675 \quad f(x) = \inf_z \{g(z) \mid h(x) \leq z\}.$$

676 Hence

$$\begin{aligned}
 677 \quad f^*(y) &= \sup_x \left\{ y^T x - \inf_z \{g(z) \mid h(x) \leq z\} \right\} \\
 678 \quad &= \sup_{x,z} \{y^T x - g(z) \mid h(x) \leq z\}.
 \end{aligned}$$

679 Using the Lagrange dual of the last problem, we get

$$\begin{aligned}
 680 \quad f^*(y) &= \min_{u \geq 0} \sup_{x, z} \{y^T x - g(z) + u(z - h(x))\} \\
 681 \quad &= \min_{u \geq 0} \left\{ \sup_z \{uz - g(z)\} + \sup_x \{y^T x - uh(x)\} \right\} \\
 682 \quad &= \min_{u \geq 0} \{g^*(u) + (uh)^*(y)\}.
 \end{aligned}$$

683 B.6 Adjoint function

684 In this section we prove the formula in line 12 of Table 2. With $g(x) = xf\left(\frac{a}{x}\right)$ we may write

$$685 \quad g^*(y) = \sup_{x \geq 0} \left\{ yx - xf\left(\frac{a}{x}\right) \right\} = \sup_{x \geq 0} \sup_w \{yx - xf(w) : xw = a\}.$$

686 Using Lagrange duality we proceed as follows:

$$\begin{aligned}
 687 \quad g^*(y) &= \sup_{x \geq 0} \sup_w \inf_z \{yx - xf(w) + z^T(a - xw)\} \\
 688 \quad &= \sup_{x \geq 0} \inf_z \left\{ yx + z^T a + x \sup_w \{-z^T w - f(w)\} \right\} \\
 689 \quad &= \inf_z \sup_{x \geq 0} \{z^T a + x(y + f^*(-z))\} \\
 690 \quad &= \inf_z \{z^T a \mid f^*(-z) \leq -y\}.
 \end{aligned}$$

691 It might be worth mentioning that if $a = 1$ then this formula implies the formula in line 9. This follows
 692 since the adjoint of the adjoint of a function yields the function itself.

693 B.7 Conjugates with ‘branches’

694 In the examples so far the conjugate of f was either a closed expression or the inf of a closed expression.
 695 This is not always the case, as we discuss in this section. Sometimes the natural domain of a function
 696 is restricted. Suppose $f : \mathbf{R} \rightarrow \mathbf{R}$ is a convex function, and a problem contains the function $f_D : \mathbf{R} \rightarrow$
 697 $[-\infty, \infty]$, defined as $f_D(x) = f(x)$ if $x \in D$ and $f_D(x) = \infty$ otherwise.

698 By way of example, let $f(x) = x^2$. The natural domain is \mathbf{R} , and $f^*(y) = \frac{1}{4}y^2$, with domain \mathbf{R} . Now let
 699 $D = [1, \infty)$. Then $f_D(x) = x^2$ if $x \geq 1$, otherwise $f_D(x) = \infty$. We then have

$$700 \quad f_D^*(y) = \sup_x \{yx - f_D(x)\}.$$

701 Define $g(x) = yx - f_D(x)$. Then the sup is finite only if $x \geq 1$ and $g(x) = yx - x^2$. Then the largest
 702 value of $g(x)$ occurs for $x = \frac{1}{2}y$, which satisfies $x \geq 1$ only if $y \geq 2$. If $y < 2$ then the largest value occurs
 703 at $x = 1$ and then it equals $y - f(1) = y - 1$. Thus we obtain

$$704 \quad f_D^*(y) = \begin{cases} y - 1 & \text{if } y \leq 2 \\ \frac{1}{4}y^2 & \text{if } y > 2. \end{cases}$$

705 Hence f_D^* is now given by two closed formulas, one for $y \leq 2$ and one for $y \geq 2$. Also note that the
 706 resulting function is continuous in $y = 2$, but not differentiable. More examples of this phenomenon can
 707 be found in [5, Table 2]. Many optimization algorithms cannot cope with nondifferentiable functions. A
 708 more tractable expression for $f_D^*(y)$ can be obtained by writing

$$709 \quad f_D(x) = f(x) + \delta_D(x).$$

710 By applying the sum rule for conjugates (i.e., line 6 in Table 2) we obtain

$$711 \quad f_D^*(y) = \inf_{y^1, y^2} \{f^*(y^1) + \delta_D^*(y^2) \mid y^1 + y^2 = y\}$$

$$712 \quad = \inf_{y^1} \{f^*(y^1) + \delta_D^*(y - y^1)\}. \quad (13)$$

713 This expression can be used immediately when forming Fenchel's dual problem, as it is the inf of a
 714 tractable function. In the above example one has

$$715 \quad \delta_D^*(y) = \sup_{x \geq 1} \{yx\} = \begin{cases} y & \text{if } y \leq 0 \\ \infty & \text{otherwise.} \end{cases}$$

716 Hence, when applying (13) to the example we obtain

$$717 \quad f_D^*(y) = \inf_z \{f^*(z) + \delta_D^*(y - z)\}$$

$$718 \quad = \inf_z \left\{ \frac{1}{4}z^2 + y - z \mid y - z \leq 0 \right\}$$

$$719 \quad = \inf_z \left\{ \left(\frac{1}{2}z - 1\right)^2 + y - 1 \mid y \leq z \right\}.$$

720 When forming a dual problem this expression is more convenient, since the 'branches' disappeared.

721 B.8 Support function of a cone

722 In this section we compute the support function of a convex cone \mathcal{K} in \mathbf{R}^n . One has

$$723 \quad \delta_{\mathcal{K}}^*(y) = \sup_{x \in \mathcal{K}} y^T x.$$

724 Recall that the dual cone of \mathcal{K} is defined by

$$725 \quad \mathcal{K}_* = \{y \mid y^T x \geq 0, \forall x \in \mathcal{K}\}.$$

726 If $y^T x > 0$ for some $x \in \mathcal{K}$ and $\lambda > 0$ then $\lambda x \in \mathcal{K}$ whereas $y^T(\lambda x) = \lambda y^T x$ goes to infinity if λ goes to
 727 infinity. Hence $\delta_{\mathcal{K}}^*(y)$ is finite only if $y^T x \leq 0$ for all $x \in \mathcal{K}$, and then the maximal value is attained at
 728 $x = 0$. As a consequence we have

$$729 \quad \delta_{\mathcal{K}}^*(y) = 0, \quad \text{dom } \delta_{\mathcal{K}}^* = -\mathcal{K}_*.$$

730 B.9 Discontinuity of the closure of a perspective function

731 In this section we deal with an example of a perspective function $(u_1 f_1)^*(y^1)$ that is finite on the closure of
 732 its domain, in spite of the fact that by definition (5) we have $(0 f_1)^*(y^1) = \infty$ if $y_1 \neq 0$. Notwithstanding
 733 this behaviour, Fenchel's dual problem works out well.

734 Let $f_0(x) = x$ and $f_1(x) := \delta_{[-1,1]}(x) - 1$. We have

$$735 \quad \min_x \{f_0(x) \mid f_1(x) \leq 0\} \equiv \min_x \{x \mid -1 \leq x \leq 1\} = -1.$$

736 Since $\delta_{[-1,1]}^*(y) = |y|$, one has $f_1^*(y) = |y| + 1$. Hence Fenchel's dual problem is

$$737 \quad \max_{y^0, y^1, u_1 \geq 0} \{-\{0 \mid y^0 = 1\} - (u_1 f_1)^*(y^1) \mid y^0 + y^1 = 0\} = \max_{u_1 \geq 0} \{-(u_1 f_1)^*(-1)\}.$$

738 One has

$$739 \quad (u_1 f_1)^*(-1) = \begin{cases} u_1 \left(\left| \frac{-1}{u_1} \right| + 1 \right) & \text{if } u_1 > 0 \\ \infty & \text{otherwise} \end{cases} = \begin{cases} 1 + u_1 & \text{if } u_1 > 0 \\ \infty & \text{otherwise.} \end{cases}$$

740 Hence the dual problem becomes

$$741 \quad \sup_{u_1 > 0} \{-1 - u_1\}.$$

742 The sup value equals -1 , as it should, but is not attained! Note that at optimality u_1 approaches zero,
 743 but $y^1 = -1$ stays away from zero.

744 C Fenchel's duality theorem

745 In this section we derive Fenchel's duality theorem from Lagrange's duality theorem. Let f and g be
 746 functions from \mathbf{R}^n to $\mathbf{R} \cup \{-\infty, \infty\}$, f proper convex and g proper concave. Then $f - g$ is convex. We
 747 consider the two problems

$$748 \quad (P) \quad \inf \{f(x) - g(x) \mid x \in \text{dom } f \cap \text{dom } g\},$$

$$749 \quad (D) \quad \sup \{g_*(y) - f^*(y) \mid y \in \text{dom } f^* \cap \text{dom } g_*\}.$$

750 **Theorem C.1 (Theorem 31.1 in [19])** *If $\text{ri.dom } f \cap \text{ri.dom } g$ is nonempty then the optimal values of*
 751 *(P) and (D) are equal and the maximal value of (D) is attained. If f is linear then $\text{ri.dom } f$ can be*
 752 *replaced by $\text{dom } f$. Similarly, if g is linear, then $\text{ri.dom } g$ can be replaced by $\text{dom } g$.*

753 **Proof:** We may reformulate (P) as follows:

$$754 \quad \inf_{u,v} \{f(u) - g(v) \mid u = v, u \in \text{dom } f, v \in \text{dom } g\}.$$

755 Defining

$$756 \quad \mathcal{C} = \{(u, v) \mid u \in \text{dom } f, v \in \text{dom } g\},$$

757 the problem can be reformulated as

$$758 \quad \inf_{u,v} \{f(u) - g(v) \mid u - v = 0, (u, v) \in \mathcal{C}\}.$$

759 The hypothesis in the theorem guarantees the existence of a point $x^* \in \text{ri.dom}(f) \cap \text{ri.dom}(g)$. As a
 760 consequence $(x^*, x^*) \in \text{ri}\mathcal{C}$. The constraints are linear and obviously (x^*, x^*) is feasible. Hence we may
 761 apply Lagrange's duality theorem. Denoting the optimal value of (P) as p^* we therefore have

$$762 \quad p^* = \max_y \inf_{u,v} \{f(u) - g(v) + (v - u)^T y \mid u \in \text{dom}(f), v \in \text{dom}(g)\}$$

$$763 \quad = \max_y \inf_{u,v} \{y^T v - g(v) - (u^T y - f(u)) \mid u \in \text{dom}(f), v \in \text{dom}(g)\}.$$

764 Note that y is fixed in the last minimization problem and u and v are independent variables. Hence we
 765 obtain

$$\begin{aligned}
 766 \quad p^* &= \max_y \left\{ \inf_{v \in \text{dom}(g)} \{y^T v - g(v)\} - \sup_{u \in \text{dom}(f)} \{u^T y - f(u)\} \right\} \\
 767 \quad &= \max_y \{g_*(y) - f^*(y)\},
 \end{aligned}$$

768 where we used that $g(v) = -\infty$ if $v \notin \text{dom}(g)$ and $f(v) = \infty$ if $v \notin \text{dom}(f)$. For the proof of the
 769 remaining statements in the theorem we refer to [19].

770 We just derived Fenchel's duality theorem from Lagrange's duality theorem. The converse, deriving
 771 Lagrange's duality theorem from Fenchel's duality theorem is also possible. So, in essence both theorems
 772 are equivalent. This has been worked out in [17].

773 D Derivation of Fenchel's dual problem

774 In this section we show how Fenchel's dual problem (D) of (P) in Section 3 can be obtained from Fenchel's
 775 duality theorem (Theorem C.1). Recall that (P) is the problem

$$776 \quad (P) \quad \inf_x \{f_0(x) \mid f_i(x) \leq 0, i = 1, \dots, m\},$$

777 where the functions $f_i : \mathbf{R}^n \rightarrow [-\infty, \infty]$ are proper convex functions for $i = 0, \dots, m$. Let $\mathcal{S}_i :=$
 778 $\{x \mid f_i(x) \leq 0\}$, for $i = 1, \dots, m$, and $\mathcal{S} = \bigcap_{i=1}^m \mathcal{S}_i$. Obviously, \mathcal{S} is the feasible region of (P). The
 779 assumption that (P) is Slater regular implies that the set \mathcal{S} is nonempty and $x \in \text{ri.dom } f$ for some
 780 $x \in \mathcal{S}$. Moreover, if f_i is nonlinear then $\text{ri } \mathcal{S}_i$ is nonempty.

781 Due to the definition of the indicator function $\delta_{\mathcal{S}}(x)$, we have

$$782 \quad \inf_{x \in \mathcal{S}} \{f_0(x)\} = \inf_x \{f_0(x) + \delta_{\mathcal{S}}(x)\}.$$

783 Since f_0 is proper convex, $-\delta_{\mathcal{S}}$ proper concave and

$$784 \quad \text{ri.dom } f_0 \cap \text{ri.dom } \delta_{\mathcal{S}} = \bigcap_{i=0}^m \text{ri.dom } f_i \neq \emptyset,$$

785 we may apply Fenchel's duality theorem. Hence we may write

$$\begin{aligned}
786 \quad \inf_{x \in \mathcal{S}} \{f_0(x)\} &= \inf_x \{f_0(x) - (-\delta_{\mathcal{S}}(x))\} \\
787 \quad &= \max_y \{(-\delta_{\mathcal{S}})^*(y) - f_0^*(y)\} \\
788 \quad &= \max_y \{-f_0^*(y) - \delta_{\mathcal{S}}^*(-y)\}, \tag{14}
\end{aligned}$$

789 where the last equality is due to (6).

790 At this stage we need to compute $\delta_{\mathcal{S}}^*(-y)$. Since the sets \mathcal{S}_i satisfy the Slater condition we may apply
791 [19, Corollary 16.4.1], which gives (see also Table 3, line 6)

$$792 \quad \delta_{\mathcal{S}}^*(-y) = \min_{\{y^i\}_{i=1}^m} \left\{ \sum_{i=1}^m \delta_{\mathcal{S}_i}^*(y^i) \mid \sum_{i=1}^m y^i = -y \right\}.$$

793 Substitution into (14) yields

$$\begin{aligned}
794 \quad \inf_{x \in \mathcal{S}} \{f_0(x)\} &= \max_y \left\{ -f_0^*(y) - \min_{\{y^i\}_{i=1}^m} \left\{ \sum_{i=1}^m \delta_{\mathcal{S}_i}^*(y^i) \mid \sum_{i=1}^m y^i = -y \right\} \right\} \\
795 \quad &= \max_{y, \{y^i\}_{i=1}^m} \left\{ -f_0^*(y) - \sum_{i=1}^m \delta_{\mathcal{S}_i}^*(y^i) \mid \sum_{i=1}^m y^i = -y \right\}, \tag{15}
\end{aligned}$$

796 where the last expression follows since the min operator is absorbed by the max operator. Thus we have
797 shown that the maximization problem in (15) has the same optimal value as (P), and hence it is a strong
798 dual problem for (P). Replacing y by y^0 this dual problem simplifies to

$$799 \quad \max_{\{y^i\}_{i=0}^m} \left\{ -f_0^*(y^0) - \sum_{i=1}^m \delta_{\mathcal{S}_i}^*(y^i) \mid \sum_{i=0}^m y^i = 0 \right\}. \tag{16}$$

800 A nice feature of the above dual problem is that its components are in one-to-one correspondence with
801 the basic elements of the primal problem: its objective function and each of the constraints functions.
802 By computing the conjugate of the objective function and the support function of each of the constraint
803 sets one obtains (16) in a structured way.

804 Usually an explicit formula for the support function is in general not available. As already mentioned
805 in Section 2, the support function of \mathcal{S}_i can be expressed in the perspective of the conjugate f_i^* of the
806 constraint function f_i . This is a consequence of the following lemma.

807 **Lemma D.1** *One has*

$$808 \quad \delta_{S_i}^*(y) = \min_{u \geq 0} \{(uf_i)^*(y)\}.$$

809 **Proof:** According to the definition of $\delta_{S_i}^*$ we have

$$810 \quad \delta_{S_i}^*(y) = \sup_{x \in S_i} \{y^T x\} = \sup_x \{y^T x \mid f_i(x) \leq 0\}.$$

811 By writing the last problem as a minimization problem and then applying Lagrange's duality theorem,
812 we obtain

$$813 \quad \delta_{S_i}^*(y) = \min_{u \geq 0} \sup_x \{y^T x - uf_i(x)\} = \min_{u \geq 0} \{(uf_i)^*(y)\},$$

814 which proves the lemma. □

815 By applying Lemma D.1 to (16) we obtain

$$816 \quad \inf_{x \in S} \{f_0(x)\} = \max_{\{y^i\}_{i=0}^m} \left\{ -f_0^*(y^0) - \sum_{i=1}^m \min_{u_i \geq 0} \{(u_i f_i)^*(y^i)\} \mid \sum_{i=0}^m y^i = 0 \right\}.$$

817 Once again the min operator is absorbed by the max operator, leaving us with the following dual problem
818 of (P):

$$819 \quad (D) \quad \max_{\{y^i\}_{i=0}^m, u} \left\{ -f_0^*(y^0) - \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=0}^m y^i = 0, u \succeq 0 \right\}.$$

820 E Tables of conjugates and support functions

821 Expressions for the convex (or concave) conjugates of many functions can be found in the existing
822 optimization literature. Strange enough, in text books on optimization tables of conjugates functions are
823 rare. One positive example is [7, p.50], but the table in this book is not exploited in a structured way to
824 derive dual problems.

825 An overview of conjugates is presented in Table 1, with references to existing literature. As we made
826 clear in Section 2, concave conjugates easily follow from these formulas by using (6).

827 In Table 2 we present some transformation rules that are useful for the computation of conjugate functions.

828 In line 9 of this table the function g is assumed to be real valued. For the more general case with g vector-
829 valued we refer to [15].

Table 1: Examples of convex conjugate functions available from the literature. In lines 6 and 7 $\|\cdot\|$ denotes any norm, and $\|\cdot\|_*$ the dual of this norm. In lines 8 - 10, the numbers p and q are related according to $\frac{1}{p} + \frac{1}{q} = 1$. In lines 13 and 14 P^\dagger denotes the generalized inverse of matrix P . In lines 15 and 16, \mathbf{S}_+^n and \mathbf{S}_{++}^n denote the set of positive semidefinite and positive definite $n \times n$ matrices, respectively.

no.	$f(x)$ (domain)	$f^*(y)$ (domain)	Reference
1	$a^T x + b$	$-b$ ($y = a$)	[8, p.91]
2	e^x	$y \log y - y$ ($y \geq 0, f^*(0) = 0$)	[19, p.105]
3	$-\log x$ ($x > 0$)	$-\log(-y) - 1$ ($y < 0$)	[19, p.106]
4	$x \log x$ ($x \geq 0, f(0) = 0$)	e^{y-1}	[8, p.92]
5	$\log(\sum_{i=1}^n e^{x_i})$	$\sum_{i=1}^n y_i \log y_i$ ($y \succeq 0, \mathbf{1}^T y = 1$)	[19, p.148]
6	$\ x\ $	0 ($\ y\ _* \leq 1$)	[8, p.93]
7	$\frac{1}{2}\ x\ ^2$	$\frac{1}{2}\ y\ _*^2$	[8, p.93]
8	$\frac{1}{p}x^p$ ($x \geq 0, p > 1$)	$\inf_z \left\{ \frac{1}{q}z^q \mid z \geq y, z \geq 0 \right\}$	Section A.5
9	$-\frac{1}{p}x^p$ ($x \geq 0, 0 < p < 1$)	$-\frac{1}{q} y ^q$ ($y < 0$)	[19, p.106]
10	$\frac{1}{p} x ^p$ ($p > 1$)	$\frac{1}{q} y ^q$	[19, p.106]
11	$\frac{1}{px^p}$ ($x > 0, p > 0$)	$-\frac{p+1}{p}(-y)^{\frac{p}{p+1}}$ ($y \leq 0$)	Section A.6
12	$-(a^2 - x^2)^{1/2}, (x \leq a)$	$a(1 + y^2)^{1/2}$	[19, p.106]
13	$\frac{1}{2}x^T P x, P \succeq 0$	$\frac{1}{2}y^T P^\dagger y$ ($y = Pz$)	[19, p.108]
14	$-\log \det X$ ($X \in \mathbf{S}_{++}^n$)	$-\log \det(-Y) - n$ ($-Y \in \mathbf{S}_{++}^n$)	[8, p.92]
15	$-\sqrt{c^T X c}, X \succeq 0$	$\inf_{z < 0} \left\{ -\frac{1}{4z} \mid -Y \succeq z c c^T \right\}$	[13, p.10]
16	$\cosh x$	$y \sinh^{-1} y - \sqrt{1 + y^2}$	[7, p.50]
17	$-\log(\cos x)$ ($ x < \frac{\pi}{2}$)	$y \tan^{-1} y - \frac{1}{2} \log(1 + y^2)$	[7, p.50]
18	$\log(\cosh x)$	$y \tanh^{-1} y + \frac{1}{2} \log(1 - y^2)$ ($ y < 1$)	[7, p.50]

Table 2: Rules for calculating the conjugate function. The functions $f(x)$, $f_i(x)$ and $h(x)$ in this table are real valued functions, $f(x)$ and $f_i(x)$ convex and $h(x)$ concave. In line 6 we assume $\cap_{i=1}^m \text{ri.dom } f_i \neq \emptyset$; if $f_i(x)$ is linear for some i then the corresponding $\text{ri.dom } f_i$ can be replaced by $\text{dom } f_i$ in this condition. We call this the *sum rule for conjugate functions*. In line 7, $x^1 : x^m$ is a partition of x . We will refer to line 8 as the *linear substitution rule*. In lines 10 and 11, f^\diamond denotes the *adjoint* of $f : \mathbf{R}_{++} \rightarrow \mathbf{R}$, which is defined as $f^\diamond(x) := xf(1/x)$. In line 13, S denotes the unit simplex.

no.	Function and assumptions	Conjugate function	Reference
1	$f(x) + a$	$f^*(y) - a$	[8, p.95]
2	$f(x + a)$	$f^*(y) - a^T y$	[19, p.107]
3	$f(x) + a^T x$	$f^*(y - a)$	[19, p.107]
4	$af(x)$, $a > 0$	$af^*\left(\frac{y}{a}\right)$	[19, p.140]
5	$f(ax)$, $a > 0$	$f^*\left(\frac{y}{a}\right)$	[19, p.107]
6	$\sum_{i=1}^m f_i(x)$	$\min_{y^i} \left\{ \sum_{i=1}^m f_i^*(y^i) \mid \sum_{i=1}^m y^i = y \right\}$	[19, p.145]
7	$\sum_{i=1}^m f_i(x^i)$	$\sum_{i=1}^m f_i^*(y^i)$	[8, p.95]
8	$f(Ax + b)$		
a	A nonsingular	$f^*(A^{-T}y) - b^T A^{-T}y$	[19, p.107]
b	otherwise	$\inf_z \{f^*(z) - b^T z \mid A^T z = y\}$	Section B.4
9	$f(g(x))$, f nondecreasing	$\inf_{z \geq 0} \{(zg)^*(y) + f^*(z)\}$	[15, p.39]
10	$\text{dom } f = \mathbf{R}_{++}$	$\inf_z \{z \mid (f^\diamond)^*(-z) \leq -y\}$	[14, p.453]
11	$h : \mathbf{R} \rightarrow \mathbf{R}$ strictly increasing	$(h^{-1})^*(y) = -(h_*)^\diamond(y)$, $y > 0$	[4, eq. (3.3)]
12	$xf\left(\frac{a}{x}\right)$, $\text{dom } f = \mathbf{R}^n$, $x \geq 0$	$\inf_z \{a^T z : f^*(-z) \leq -y\}$	Section B.6
13	$\max_i \{f_i(x) \mid i = 1, \dots, m\}$	$\min_{z \in S} \left\{ \sum_{i=1}^m (z_i f_i)^*(y^i) \mid \sum_{i=1}^m y^i = y \right\}$	[19, p.149]

Table 3: Support functions. In line 3, \mathcal{K} is a pointed cone and if \mathcal{K} is nonlinear, then it is assumed that S contains a strictly feasible point. The functions f_i and the sets S_i in lines 5-7 are closed convex. In line 5 we assume $\cap_{i=1}^m \text{ri.dom } f_i \neq \emptyset$; if $f_i(x)$ is linear for some i then the corresponding $\text{ri.dom } f_i$ can be replaced by $\text{dom } f_i$ in this condition. Similarly, in line 6 we assume $\cap_{i=1}^m \text{ri } S_i \neq \emptyset$, and if S_i is polyhedral then the corresponding $\text{ri } S_i$ can be replaced by S_i .

no.	$S, S \neq \emptyset$	$\delta_S^*(y)$	Reference
1	$\{x \mid Ax = b\}$	$\min_z \{b^T z \mid A^T z = y\}$	[8, p.380]
2	$\{x \mid Ax \preceq b\}$	$\min_z \{b^T z \mid A^T z = y, z \succeq 0\}$	[8, p.380]
3	$\{x \mid b - Ax \in \mathcal{K}\}$	$\min_z \{b^T z \mid A^T z = y, z \in \mathcal{K}_*\}$	[6, p.272]
4	$\{x \mid \ x\ _p \leq \rho\}$	$\rho \ y\ _q, \frac{1}{p} + \frac{1}{q} = 1$	[6, p.272]
5	$\{x \mid f_i(x) \leq 0, i = 1, \dots, m\}$	$\min_{u \succeq 0, y^i} \left\{ \sum_{i=1}^m (u_i f_i)^*(y^i) \mid \sum_{i=1}^m y^i = y \right\}$	[19, p.146]
6	$S = \bigcap_{i=1}^m S_i$	$\min_{y^i} \left\{ \sum_{i=1}^m \delta_{S_i}^*(y^i) \mid \sum_{i=1}^m y^i = y \right\}$	[19, p.146]
7	$S = S_1 \times \dots \times S_m$	$\sum_{i=1}^m \delta_{S_i}^*(y^i), y^1 : y^m \text{ is a partition of } y$	[6, p.294]

830 Table 3 gives a list of support functions, and Table 4 provides a list of cones and their dual cones.

Table 4: An overview of the most commonly used cones and their dual cones. The first three cones are self-dual. The set \mathbf{S}^n is the set of symmetric $n \times n$ matrices. The closure operator is denoted by cl .

Name	Cone \mathcal{K}	Dual cone \mathcal{K}_*
Nonnegative orthant	\mathbf{R}_+^n	\mathbf{R}_+^n
Second order cone	$\mathbf{L}^n = \{(x, t) \in \mathbf{R}^{n+1} \mid \ x\ \leq t\}$	\mathbf{L}^n
Positive semidefinite cone	$\mathbf{S}_n^+ = \{X \in \mathbf{S}^n \mid X \succeq 0\}$	\mathbf{S}_n^+
Power cone	$\{x \in \mathbf{R}^n \mid x_n \leq \prod_{i=1}^{n-1} x_i^{\alpha_i}\}$	$\{y \in \mathbf{R}^n \mid y_n \leq \prod_{i=1}^{n-1} (y_i / \alpha_i)^{\alpha_i}\}$
Exponential cone	$\text{cl} \{x \in \mathbf{R}^3 \mid x_2 e^{x_1/x_2} \leq x_3, x_2 > 0\}$	$\text{cl} \{y \in \mathbf{R}^3 \mid -y_1 e^{y_2/y_1-1} \leq y_3, y_1 < 0\}$
Copositive cone	$\{X \in \mathbf{S}^n \mid v^T X v \geq 0 \forall v \geq 0\}$	$\{Y \in \mathbf{S}^n \mid \exists A \in \mathbf{R}_+^{n \times m} Y = AA^T\}$

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