

Minimizing the difference of convex and weakly convex functions via bundle method

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Abstract: We consider optimization problems with objective and constraint being the difference of convex and weakly convex functions. This framework covers a vast family of nonsmooth and nonconvex optimization problems, particularly those involving certain classes of composite and nonconvex value functions. We investigate several stationary conditions and extend the proximal bundle algorithm of [van Ackooij et al., *Comput. Optim. Appl.*, 78 (2021), pp. 451–490] to compute critical points for problems in this class. Our modifications on that algorithm boil down to a different master program and an original rule to update the proximal parameter to ensure convergence in this more general setting. Thanks to this new rule, no pre-estimation of the underlying weakly-convex moduli is needed, opening the way to deal with optimization problems for which no practical and mathematically sound algorithms exist. Numerical experiments on some nonconvex stochastic problems illustrate the practical performance of the method.

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Dedicated to Professor Masao Fukushima.

1 Introduction.

This work is in the stream of methods for structured nonsmooth and nonconvex optimization problems, one of the optimization fields to which Professor Maso Fukushima significantly contributed. In particular, his seminal paper [23] with H. Mine proposes an algorithm for minimizing the sum of a convex function and a continuously differentiable function. Such an algorithm is nowadays known as the generalized conditional gradient method and was instrumental to several algorithmic developments in the field, particularly in DC programming. By considering a significantly broader setting, this work deals with nonsmooth and nonconvex optimization problems of the form

$$\min_{x \in X} f(x) \quad \text{s.t.} \quad c(x) \leq 0, \quad (1.1a)$$

where X is a nonempty bounded polyhedron contained in an open convex set $\mathcal{O} \subset \mathbb{R}^n$, and functions $f : \mathcal{O} \rightarrow \mathbb{R}$ and $c : \mathcal{O} \rightarrow \mathbb{R}$ are decomposable as the difference of *convex* and (locally) *weakly convex* functions. More specifically, we assume that the following convex-weakly convex (CwC) decompositions are available:

$$f(x) = f_1(x) - f_2(x) \quad \text{and} \quad c(x) = c_1(x) - c_2(x), \quad (1.1b)$$

with $f_1, c_1 : \mathcal{O} \rightarrow \mathbb{R}$ convex and $f_2, c_2 : \mathcal{O} \rightarrow \mathbb{R}$ *weakly convex functions on some neighbourhood of each* $x \in \mathcal{O}$. We adopt the more general definition of weakly convex functions (see Definition 2.2 below) given in [51, Def. 4.2] so that we can exploit the equivalence between the families of locally weakly convex and Lower- C^2 functions [33, Thm. 1.3, Cor. 1.3] to highlight the breadth of our approach. In particular, we have in mind the following settings for f_2 (as well as for c_2):

- i) $f_2(x) = \phi(x)$ is a (possibly nonsmooth) convex function;
- ii) $f_2(x) = -h(x)$ with h having Lipschitz continuous gradient;
- iii) $f_2(x) = \phi(x) - h(x)$, with ϕ and h as given above;

- iv) $f_2(x)$ is the optimal value of $\max_{t \in T} F(t, x)$, with T a (possibly nonconvex) compact set and F of class C^2 ;
- v) $f_2(x) = \phi(G(x))$, with $\phi : \mathbb{R}^m \rightarrow \mathbb{R}$ convex and Lipschitz and $G : \mathbb{R}^n \rightarrow \mathbb{R}^m$ a smooth mapping with Lipschitz Jacobian.

Analogous settings for c_2 , and their combinations with the ones for f_2 , are covered by our analysis (see Section 2 below for details).

Weakly convex functions enjoy favorable properties in so much as that they can be recast as *Difference-of-Convex* (DC) functions [10]. Hence, problem (1.1) can, in theory, be recast as a DC-constrained DC program, a setting that proves *practical if explicit DC decompositions are available*; see for instance [17, 16, 34, 24, 8, 50, 42, 35] and references therein. However, if no DC decomposition is known for f or c , the DC machinery is unsuitable, and the methods proposed in these references are not applicable. This is already the case for the more straightforward items ii) and iii) above if the underlying Lipschitz constant is unknown and no upper bound is readily available¹. The situation becomes even more complicated for items iv) and v): in general, there are no formulae, rules, or practical insights to obtain a DC decomposition for f_2 in these cases (see Example 1.1 below for a particular case of iv). A strategy to handle problem (1.1) via DC programming algorithms is to replace functions f_i and c_i ($i = 1, 2$) with $f_i(x) + \frac{\mu}{2}\|x\|^2$ and $c_i(x) + \frac{\mu}{2}\|x\|^2$ for a large parameter $\mu > 0$ estimating upper bounds on the unknown weakly-convex moduli μ_f and μ_c of f_2 and c_2 (see Proposition 2.4), hoping that $f_2(x) + \frac{\mu}{2}\|x\|^2$ and $c_2(x) + \frac{\mu}{2}\|x\|^2$ are convex on X . As, in general, there is no reliable way to assert the convexity of these latter functions, DC programming algorithms applied in this context must be understood as heuristics. Remarkably, the work [40] exploits such a strategy by combining a dynamic rule to update μ with a nonconvexity test so that convergence is achieved, but only in a probability sense. Differently, for a class of nonconvex two-stage stochastic problems, the authors of [20] exploit an implicitly convex-concave structure of the objective function and propose an algorithm based on the so-called partial Moreau envelope that disregards DC decompositions at the price of nonnegligible computational costs.

In contrast to the above references, this work investigates a bundle method approach for tackling (1.1), which neither requires explicit DC decompositions of the involved functions (in particular, bounds on the weakly-convex moduli μ_f and μ_c need not be known), nor relies on (often costly) Moreau envelopes. For the method to work, it suffices to dispose of a difference of *convex and weakly convex* (CwC) decomposition of the involved functions, as in (1.1b). Compared to DC, the latter structure appears more naturally in applications (see [15, § 7.5]) and has yet to be exploited to design optimality conditions and numerical algorithms. This work aims to fill this gap.

Our approach broadens and enhances the method proposed in [50] for dealing with DC-constrained DC-problems in two ways. First, the availability of DC decompositions is no longer needed, which makes our approach applicable to a larger scope of problems. Second, it is ensured to compute critical points for the original problem without any additional assumption on the second components: f_2 and c_2 need not be continuously differentiable as assumed in [50, Thm. 2]. In addition, it has a lower cost per iteration (the master subproblem has fewer constraints than the one of [50]). Similarly to [50], our approach builds upon a problem reformulation via improvement function, a well-known and successful strategy in the nonsmooth optimization literature [31, 1, 24]. However, due to the above modifications, the convergence analysis of our extension of the method proposed in [50] must be done anew. Furthermore, a new criticality definition for the reformulated problem links directly with (necessary) optimality conditions for the original problem (1.1), which makes it a major ingredient for these enhancements. Such a criticality concept is introduced and analyzed in Section 3 below, where we also extend the alternative characterization of Bouligand stationarity given in [27] to our CwC setting. Before that, we motivate this work with the following example that presents a class of problems (of great practical appeal) where the CwC decomposition arises upon applying a well-known interior-penalty strategy. An example of a real-life (chance-constrained optimal power flow) problem fitting our CwC structure without any approximation can be found in the Ph.D. thesis [15, § 7.5].

Example 1.1 (Nonconvex two-stage programming). Let $\Xi := \{\xi^1, \dots, \xi^S\}$ be a set of scenarios and $\pi_s > 0$ the probability of occurrence of event ξ^s , $s = 1, \dots, S$. Consider the following nonconvex two-stage program

$$\begin{cases} \min_{x \in X} & f_1(x) + \sum_{s=1}^S \pi_s Q(x; \xi^s) \\ \text{s.t.} & c_1(x) - c_2(x) \leq 0 \end{cases} \quad \text{with} \quad Q(x; \xi) := \begin{cases} \min_{y \in Y} & q(x, y; \xi) \\ \text{s.t.} & \psi_i(x, y; \xi) \leq 0, \quad i = 1, \dots, m. \end{cases} \quad (1.2)$$

Assume that:

- $f_1, c_1, c_2 : \mathbb{R}^n \rightarrow \mathbb{R}$ are convex (possibly nonsmooth) functions;

¹It is worth noting that in many practical problems, mainly those from data science, an upper bound on such a constant can be computed.

- $X \subset \mathbb{R}^n$, $Y \subset \mathbb{R}^{n_2}$ are two (non-empty) convex and bounded polytopes;
- $q, \psi_i : \mathbb{R}^n \times \mathbb{R}^{n_2} \times \Xi \rightarrow \mathbb{R}$, $i = 1, \dots, m$, possess the following characteristics: $q(\cdot, \cdot, \xi)$ and $\psi_i(\cdot, \cdot, \xi)$ are twice-continuously differentiable for every $\xi \in \Xi$ fixed and, moreover, $q(x, \cdot, \xi)$ and $\psi_i(x, \cdot, \xi)$ are convex for every x and ξ fixed;
- the constraints in the subproblem $Q(x; \xi)$ satisfy the Slater condition: for every $x \in X$ and $\xi \in \Xi$, there exists $y^o(x; \xi) \in Y$ such $\psi_i(x, y^o(x; \xi), \xi) < 0$, $i = 1, \dots, m$.

As presented in [50], the DC constraint $c_1(x) - c_2(x) \leq 0$ above is particularly useful in this stochastic programming setting to model chance constraints.

Under the above assumptions, evaluating the recourse function $Q(x; \xi)$ amounts to solving a well-defined convex optimization problem on variable y . Although this essential property is present, the recourse function itself fails to be convex on variable x (but $Q(\cdot; \xi)$ is continuous as a result of [2, Prop. 4.4]). Furthermore, without further assumptions, computing a (generalized) subgradient of $Q(\cdot; \xi)$ at x as well as asserting additional properties about this function are challenging tasks. This could for instance be done if the constraints satisfy a further Aubin or Lipschitz like property upon exploiting [25, Chapter 4]. Still though, most likely, at best we would be dealing with subdifferentials inclusions - and concrete algorithms to handle such general “marginal functions” would be unavailable.

A possible manner to curtail these difficulties is to approximate the recourse function with a more tractable one. As explained in [3], with the help of the log-barrier penalty function and a penalization parameter $\varepsilon > 0$, we may approximate $Q(x; \xi)$ with

$$Q^\varepsilon(x; \xi) := \min_{y \in Y} q(x, y; \xi) - \frac{1}{\varepsilon} \sum_{i=1}^m \log(-\psi_i(x, y; \xi)). \quad (1.3)$$

Given the above assumptions, it is well known that $Q^\varepsilon(x; \xi) \downarrow Q(x; \xi)$ as $\varepsilon \downarrow 0$ (e.g., [3, § 2.2] and [2, p. 266]), and thus the model

$$\begin{cases} \min_{x \in X} & f_1(x) - f_2(x) \\ \text{s.t.} & c_1(x) - c_2(x) \leq 0 \end{cases} \quad \text{with} \quad f_2(x) := \sum_{s=1}^S \pi_s [-Q^\varepsilon(x; \xi^s)]$$

is an accurate approximation of (1.2) when $\varepsilon > 0$ is small enough. Furthermore, as $-Q^\varepsilon(x; \xi) = \max_{y \in Y} \frac{1}{\varepsilon} \sum_{i=1}^m \log(-\psi_i(x, y; \xi)) - q(x, y; \xi)$ is a weakly convex function (c.f. item iv) above), this model fits the structure (1.1). We highlight that $Q^\varepsilon(x; \xi)$ is generally a nonsmooth (nonconvex) function; hence, the above problem is challenging. To our knowledge, no practical and mathematically sound optimization algorithm could tackle this class of problems before this work. Indeed, [23] requires f_2 to be smooth, [20] assumes $q(\cdot, \cdot, \xi)$ and $\psi_i(\cdot, \cdot, \xi)$ to be concave-convex functions, and [3] requires another degree of approximation by adding a Tikhonov regularization to (1.3) to force $Q^\varepsilon(x; \xi)$ be smooth. In all these references, function c_2 is absent. Being nonsmooth, we mention in passing that a (generalized) subgradient of f_2 at x can be computed and seen to be $\sum_{s=1}^S \pi_s g(y(x; \xi^s))$, where $g(\cdot) := \nabla_x [\frac{1}{\varepsilon} \sum_{i=1}^m \log(-\psi_i(x, \cdot; \xi^s)) - q(x, \cdot; \xi^s)]$ is the gradient w.r.t. x of the objective function of (1.3) multiplied by -1 , and $y(x; \xi)$ is an arbitrary optimal solution of (1.3) (see Proposition 2.1 and [30, Thm. 7.3]). \square

Our approach is still applicable in more general case, where the probability vector π in the above example is a function (of class C^2) of the first-stage variable x , i.e., $\pi_s(x)$, $i = 1, \dots, S$. Hence, this work’s class of optimization problems includes the challenging family of stochastic programming recourse models with decision-dependent uncertainty considered (e.g. [12] and [20]).

The remainder of this manuscript is organized as follows. Section 2 recalls essential definitions, key elements, and well-known concepts from variational analysis. Necessary optimality conditions for problem (1.1) are presented in Section 3 as well as the problem reformulation via an improvement function. Once the link between the reformulated and the original problem is established in the same section, Section 4 focuses on an improvement-function-based bundle method for problem (1.1). Section 5 presents the method’s convergence analysis to critical points, and finally, Section 6 illustrates the practical performance of our approach on some nonconvex stochastic optimization problems, including a real-life (chance-constrained optimal power flow) problem from [15, 36], and a compressed sensing problem.

Notation The following notation is employed throughout the text. For a real number a , we denote by $[a]_+$ the value $\max\{a, 0\}$. For any points $x, y \in \mathbb{R}^n$, $\langle x, y \rangle$ stands for the Euclidean inner product, and $\|\cdot\|$ for the associated norm, i.e., $\|x\| = \sqrt{\langle x, x \rangle}$. For a convex set X , $N_X(x)$ stands for its normal cone at the point x , i.e., the set $\{y : \langle y, z - x \rangle \leq 0 \text{ for all } z \in X\}$ if $x \in X$ and the empty set otherwise. The Bouligand tangent cone to a (possibly nonconvex) set $W \subset \mathbb{R}^n$ at a point $w \in W$ is the set $\mathcal{T}_W(w)$ of all tangent directions in the following sense: $d \in \mathcal{T}_W(w)$ if there exist a sequence of vectors $\{w^k\} \subset W$ converging to w and a sequence of positive scalars $t_k \rightarrow 0$ such that $d = \lim_{k \rightarrow \infty} (w^k - w)/t_k$. The indicator function of

$X \subset \mathbb{R}^n$ is defined as $i_X(x) = 0$ if $x \in X$ and $i_X(x) = +\infty$ otherwise. The convex hull of a set X is $\text{conv}X$ and the relative interior is denoted by $\text{ri} X$. The domain of a function $\varphi : \mathbb{R}^n \rightarrow (-\infty, +\infty]$ is represented by $\text{Dom}(\varphi) = \{x \in \mathbb{R}^n : \varphi(x) < +\infty\}$. Notation \mathcal{O} stands for an open convex set of the Euclidean space \mathbb{R}^n and, given the definitions of f and c , we have that $\mathcal{O} \subset \text{Dom}(f)$ and $\mathcal{O} \subset \text{Dom}(c)$. The component functions of f and c are f_1, f_2 , and c_1, c_2 respectively: f_1 and c_1 are convex, whereas f_2 and c_2 are weakly convex on some neighbourhood of every $x \in \mathcal{O}$. Finally, f^* stands for the Legendre-Fenchel transform of a function $f : \mathbb{R}^n \rightarrow (-\infty, +\infty]$.

2 Definition and prerequisites

This section starts by recalling the concept of (generalized) directional derivatives and subdifferentials. Basic subdifferential calculus is summarized in Proposition 2.1 below, followed by the definitions of weakly convex and lower- C^2 functions. The section closes with Proposition 2.4 asserting that the definition of (locally) weakly convex function can be globally extended to the whole convex and compact set X . Such a property is of crucial importance in this work.

A function $f : \mathcal{O} \rightarrow \mathbb{R}$ is said to be *locally Lipschitz continuous* if for each $x' \in \mathcal{O}$ there is a neighbourhood $V_{x'} \subset \mathcal{O}$ of x' such that, for some $L_{x'} \geq 0$,

$$|f(x) - f(y)| \leq L_{x'} \|x - y\| \quad \forall x, y \in V_{x'}.$$

The function f is said to be *Lipschitz continuous* on \mathcal{O} if $L_{x'} = L$ can be taken independent of $x' \in \mathcal{O}$, and $V_{x'}$ in the above inequality is replaced with \mathcal{O} .

Directional derivatives and subdifferentials. Let $\phi : \mathcal{O} \rightarrow \mathbb{R}$ be a convex function. Then ϕ is locally Lipschitz continuous and, for each $x \in \mathcal{O}$, the directional derivative

$$\phi'(x; d) := \lim_{\tau \downarrow 0} \frac{\phi(x + \tau d) - \phi(x)}{\tau}$$

exists (and is finite) in every direction $d \in \mathbb{R}^n$ [26, Prop. 2.81 and Cor. 2.82]. Such a derivative can be represented by $\phi'(x; d) = \max_{s \in \partial\phi(x)} \langle s, d \rangle$, where $\partial\phi(x)$ is the *subdifferential* of ϕ at point x :

$$\partial\phi(x) := \{s \in \mathbb{R}^n : \phi(y) \geq \phi(x) + \langle s, y - x \rangle \quad \forall y \in \mathbb{R}^n\}. \quad (2.1)$$

The elements of $\partial\phi(x)$ are referred to as the *subgradients* of ϕ at x . The *approximate subdifferential* is defined, for $\epsilon \geq 0$, by

$$\partial_\epsilon\phi(x) := \{s \in \mathbb{R}^n : \phi(y) \geq \phi(x) + \langle s, y - x \rangle - \epsilon \quad \forall y \in \mathbb{R}^n\}.$$

Let $f : \mathcal{O} \rightarrow \mathbb{R}$ be a locally Lipschitz continuous function. Then the generalized directional derivative defined by

$$f^\circ(x; d) := \limsup_{x' \rightarrow x, \tau \downarrow 0} \frac{f(x' + \tau d) - f(x')}{\tau}$$

is finite for all $x \in \mathcal{O}$ in every direction $d \in \mathbb{R}^n$ [4, Prop. 2.1.1(a)]. Such a mathematical concept permits us to define the *Clarke subdifferential* of f at $x \in \mathcal{O}$,

$$\partial^c f(x) := \{g \in \mathbb{R}^n : \langle g, d \rangle \leq f^\circ(x; d) \text{ for all } d \in \mathbb{R}^n\}, \quad (2.2)$$

which is a nonempty, convex, and compact subset of \mathbb{R}^n [4, Prop. 2.1.2(a)] satisfying $f^\circ(x; d) = \max_{g \in \partial^c f(x)} \langle g, d \rangle$.

The elements of $\partial^c f(x)$ are referred to as *generalized (or Clarke) subgradients*, as they are the usual subgradients, i.e., $\partial^c f = \partial f$, when f is convex [4, Prop. 2.2.7]. Furthermore, when f is continuously differentiable, $\partial^c f(x)$ reduces to the singleton $\{\nabla f(x)\}$. An alternative representation, in finite dimensions, of $\partial^c f(x)$ is (see [4, Thm. 2.5.1])

$$\partial^c f(x) := \text{conv} \left\{ \lim_{i \rightarrow \infty} \nabla f(x_i), x_i \rightarrow x, f \text{ differentiable at } x_i \right\}.$$

A fundamental result, often evoked in this work, is the following one [4, Prop. 2.1.2]: the mapping $\partial^c f$ is locally bounded in the interior of $\text{Dom}(f) := \{x \in \mathbb{R}^n : f(x) < \infty\}$. As a result, the image $\partial^c f(X)$ of every bounded set $X \subset \mathcal{O}$ ($\subset \text{Dom}(f)$) is bounded in \mathbb{R}^n . Useful calculus rules of subdifferentials are listed in Proposition 2.1 below and rely on the concept of regularity.

A locally Lipschitz continuous function $f : \mathcal{O} \rightarrow \mathbb{R}$ is *subdifferentially regular* (or simply *regular*) at $x \in \mathcal{O}$ if for every $d \in \mathbb{R}^n$ the ordinary directional derivative at x exists and coincides with the generalized one:

$$f'(x; d) = f^\circ(x; d) \quad \forall d \in \mathbb{R}^n.$$

It holds that smooth functions, as well as convex ones, are regular at every point in the interior of their domains. Moreover, a finite linear combination (by nonnegative scalars) of regular functions at x is regular [4, Prop. 2.3.6].

Proposition 2.1. Let $f_t : \mathcal{O} \rightarrow \mathbb{R}$, $t = 1, 2, \dots, m$, be locally Lipschitz functions and $x \in \mathcal{O}$ an arbitrary point. Then

- i) $\partial^c[\sum_{t=1}^m a_t f_t](x) \subset \sum_{t=1}^m a_t \partial^c f_t(x)$ for all $a \in \mathbb{R}^m$, and equality holds if
 - all but one of f_t are smooth [4, Prop. 2.3.3 and Cor. 2];
 - or if every f_t is regular at x and $a \in \mathbb{R}_+^m$ [4, Cor. 3];
- ii) $\partial^c f(x) \subset \text{conv}\{\partial^c f_t(x) : t \in I(x)\}$, for $f(x) = \max_{t=1, \dots, m} f_t(x)$ and $I(x) := \arg \max_{t=1, \dots, m} f_t(x)$, and equality holds and f is regular if every f_t is regular at x . [4, Prop. 2.3.12].

The last item can be strengthened when more structure is present, such as in the case of weakly convex functions (see Eq. (2.3) below).

Weakly convex functions: definition and main properties

Definition 2.2 (Def. 4.2 [51]). A function $f : \mathcal{O} \rightarrow \mathbb{R}$ is said to be (locally) *weakly convex* on \mathcal{O} if, on some neighbourhood $V_{x'} \subset \mathcal{O}$ of each $x' \in \mathcal{O}$, there exists $\mu_{x'} \geq 0$ such that, for all $\mu \geq \mu_{x'}$

$$\phi(x) := f(x) + \frac{\mu}{2} \|x\|^2 \quad \text{is finite and convex on } V_{x'}.$$

Furthermore, f is said to be *weakly convex in the global sense* on \mathcal{O} if the above property holds for $V_{x'} = \mathcal{O}$ and $\mu_{x'} = \bar{\mu} \geq 0$ regardless of $x' \in \mathcal{O}$. \square

Clearly, a convex function on \mathcal{O} is weakly convex in the global sense: it suffices to take $\mu_{x'} = 0$ and $V_{x'} = \mathcal{O}$ for all $x' \in \mathcal{O}$. When f is a smooth function with a Lipschitz continuous gradient, then f is weakly convex in the global sense with $\mu = L$ the Lipschitz constant of ∇f [6, Prop. 1]. Moreover, it follows from [7, Lemma 4.2] that the family of composite functions given in item v) of the Introduction is also weakly convex in the global sense.

Definition 2.2 implies that weakly convex functions are *locally DC*: the decomposition $f(x) = \phi(x) - \frac{\mu}{2} \|x\|^2$ holds on some neighbourhood of every $x' \in \mathcal{O}$. As a result, [29, Thm. 10.33] ensures that the class of weakly convex functions coincides with that of Lower- C^2 functions; see also [33, Thm. 1.3, Cor. 1.3].

Definition 2.3 (Def. 10.29 [29]). (LC² functions). A function $f : \mathcal{O} \rightarrow \mathbb{R}$ is said to be *Lower- C^2* or LC² on \mathcal{O} if, on some neighbourhood $V_{x'} \subset \mathcal{O}$ of each $x' \in \mathcal{O}$, there is a representation

$$f(x) = \max_{t \in T} f_t(x).$$

in which the functions f_t are of differentiability class C^2 on $V_{x'}$ and the index set T is a compact space such that $f_t(x)$, $\nabla f_t(x)$, and $\nabla^2 f_t(x)$ depend continuously not just on $x \in V_{x'}$ but jointly on $(t, x) \in T \times V_{x'}$. \square

In particular, if f is given by $f(x) = \max\{f_1(x), \dots, f_m(x)\}$ and all functions f_1, \dots, f_m are of class C^2 , then f is Lower- C^2 /weakly convex. Furthermore, the functions of item iv) are also weakly convex, since they are Lower- C^2 by definition.

An important property of LC²/weakly convex functions is regularity [28, Thm. 1]: for every $x \in \mathcal{O}$, the equality $f'(x; d) = f^\circ(x; d)$ holds in every direction $d \in \mathbb{R}^n$. Furthermore, Theorem 7.3 in [30] gives the following characterization of the Clarke subdifferential of f at $x \in \mathcal{O}$: for $I(x) = \arg \max_{t \in T} f_t(x)$,

$$\partial^c f(x) = \text{conv} \{ \nabla_x f_t(x) : t \in I(x) \} \quad \text{for all } x \in \mathcal{O}. \quad (2.3)$$

When constrained to a compact convex set $X \neq \emptyset$, we can say more about weakly convex functions. Indeed, the local property in Definition 2.2 globally extends to the whole X , and we have the following key result (whose proof can be found in the Appendix A).

Proposition 2.4. Let $f : \mathcal{O} \rightarrow \mathbb{R}^n$ be a weakly convex function, and $X \subset \mathcal{O}$ a compact and convex set. Then there exist a real number $\mu_f \geq 0$ and an open convex set \mathcal{O}' satisfying $X \subset \mathcal{O}' \subset \mathcal{O}$ such that, for all $\mu \geq \mu_f$:

- i) the function $\phi(x) := f(x) + \frac{\mu}{2} \|x\|^2$ is convex on \mathcal{O}' and $\partial \phi(x) = \partial^c f(x) + \mu x$ for all $x \in \mathcal{O}'$;
- ii) for all $s_f \in \partial^c f(x)$ with $x \in \mathcal{O}'$, the following inequality holds

$$f(y) \geq f(x) + \langle s_f, y - x \rangle - \frac{\mu}{2} \|y - x\|^2 \quad \forall y \in X. \quad (2.4)$$

Concerning the setting of this work where X is compact, the appealing DC decomposition $f(x) = \phi(x) - \frac{\mu_f}{2} \|x\|^2$ is, unfortunately, unavailable: the threshold μ_f in Proposition 2.4 is in general unknown. This fact precludes the application of DC techniques to optimization problems featuring general Lower- C^2 /weakly convex functions. Interested readers are referred to [41] for a strategy that uses approximated DC decompositions based on item i) of Proposition 2.4.

3 Necessary optimality conditions and problem reformulation

Let $f, c : \mathcal{O} \rightarrow \mathbb{R}$ be given by (1.1b). We highlight that some properties of their components can be transferred to these functions. (To ease the presentation, let us focus only on $f(x) = f_1(x) - f_2(x)$, as the same conclusions hold for c .) For instance, f is locally Lipschitz continuous because f_1 and f_2 are so. Furthermore, as f_1 is convex and f_2 is (locally) weakly convex, they are both directional differentiable and these properties extend to f as well: for every $x \in \mathcal{O}$, the directional derivative of f is finite in every direction $d \in \mathbb{R}^n$ as result of the following relation:

$$f'_1(x; d) - f'_2(x; d) = \lim_{\tau \downarrow 0} \left[\frac{f_1(x + \tau d) - f_1(x)}{\tau} - \frac{f_2(x + \tau d) - f_2(x)}{\tau} \right] = f'(x; d).$$

However, the important regularity condition of both f_1 and f_2 does not extend to f as a mere fact that the latter is not a linear combination with nonnegative coefficients of the two former functions (see Proposition 2.1.i). Hence, we cannot expect to have equality in the following inclusion

$$\partial^c f(x) \subset \partial f_1(x) - \partial^c f_2(x)$$

unless one of the component functions is smooth at x . Such an inclusion impacts stationary concepts as we will now discuss. Let us first consider the convexly-constrained problem

$$\min_{x \in X} f(x), \quad \text{with } f(x) = f_1(x) - f_2(x). \quad (3.1)$$

A point $\bar{x} \in X$ is said to be directional (d)-stationary for this problem if $f'(\bar{x}; d) \geq 0$ for all $d \in \mathcal{T}_X(\bar{x})$. The following result generalizes [27, Prop. 5], where a specific case of problem (3.1) with $f_2(x) = \max\{\psi_1(x), \dots, \psi_m(x)\}$ and convex ψ_1, \dots, ψ_m is considered.

Proposition 3.1. *A point $\bar{x} \in X$ is d -stationary of problem (3.1) if, and only if,*

$$\bar{x} \in \arg \min_{x \in X} f_1(x) - [f_2(\bar{x}) + \langle s_{f_2}, x - \bar{x} \rangle] \quad \forall s_{f_2} \in \partial^c f_2(\bar{x}).$$

Proof. Observe that $\mathcal{T}_X(\bar{x}) = \text{cl}\{d \in \mathbb{R}^n : d = t(x - \bar{x}), x \in X, t \in \mathbb{R}_+\}$ due to convexity of X . Therefore, the definition of d -stationarity can be equivalently written as $f'(\bar{x}; x - \bar{x}) \geq 0$ for all $x \in X$. Recall that f_2 is (locally) weakly convex and thus regular, which implies that $f'_2(\bar{x}; x - \bar{x}) = \max_{s_{f_2} \in \partial^c f_2(\bar{x})} \langle s_{f_2}, x - \bar{x} \rangle$. Hence,

$$\begin{aligned} f'(\bar{x}; x - \bar{x}) &\geq 0 && \forall x \in X \\ \Leftrightarrow f'_1(\bar{x}; x - \bar{x}) - f'_2(\bar{x}; x - \bar{x}) &\geq 0 && \forall x \in X \\ \Leftrightarrow f'_1(\bar{x}; x - \bar{x}) - \langle s_{f_2}, x - \bar{x} \rangle &\geq 0 && \forall s_{f_2} \in \partial^c f_2(\bar{x}), \forall x \in X \\ \Leftrightarrow \bar{x} \in \arg \min_{x \in X} f_1(x) - \langle s_{f_2}, x - \bar{x} \rangle &&& \forall s_{f_2} \in \partial^c f_2(\bar{x}). \end{aligned}$$

□

A point $\bar{x} \in X$ is said to be Clarke-stationary of problem (3.1) if

$$0 \in \partial^c f(\bar{x}) + N_X(\bar{x}). \quad (3.2)$$

Furthermore, by following the lead of DC programming (see for instance [6, §3.1]), $\bar{x} \in X$ is said to be a critical point if

$$0 \in \partial f_1(\bar{x}) - \partial^c f_2(\bar{x}) + N_X(\bar{x}). \quad (3.3)$$

It is not difficult to see that this inclusion means that

$$\bar{x} \in \arg \min_{x \in X} f_1(x) - [f_2(\bar{x}) + \langle s_{f_2}, x - \bar{x} \rangle] \quad \text{for some } s_{f_2} \in \partial^c f_2(\bar{x}).$$

Note that the concept of criticality is weaker than that of Clarke-stationarity, which in turn is weaker than d -stationarity (because $f'(\cdot; d) \leq f^\circ(\cdot; d)$ for all $d \in \mathbb{R}^n$). However, criticality and Clarke-stationarity coincide when at least one component function is smooth (in which case f is regular). Furthermore, we can see from Proposition 3.1 and the above alternative characterization of criticality that the three concepts coincide when f_2 is continuously differentiable at \bar{x} .

For the more general problem (1.1), $\bar{x} \in X$ is said to be a Bouligand (B)-stationary point of (1.1) if $f'(\bar{x}; d) \geq 0$ for all $d \in \mathcal{T}_{X^c}(\bar{x})$, with X^c the feasible set of (1.1). If the considered point strictly satisfies the nonconvex constraint, i.e. $c(\bar{x}) < 0$, then B -stationarity condition is equivalent to d -stationarity condition for problem (3.1), as $\mathcal{T}_{X^c}(\bar{x}) = \mathcal{T}_X(\bar{x})$. Analogously, B -stationarity boils down to d -stationarity if constraint $c(x) \leq 0$ is absent. Necessary and sufficient conditions for B -stationarity are given in [27, Prop. 4] for the case of DC-constrained DC problems. The next result deals with a more general case: we assume that only c_2 is convex, while f_2 remains a weakly-convex function.

Proposition 3.2. *In addition to our assumptions on problem (1.1), let $c_2 : \mathcal{O} \rightarrow \mathbb{R}$ be a convex function and $\bar{x} \in X^c := \{x \in X : c(x) \leq 0\}$ such that $c(\bar{x}) = 0$. Moreover, assume that the following constraint qualification (CQ) holds*

$$\text{c1} \{d \in \mathcal{T}_X(\bar{x}) : c'(\bar{x}; d) < 0\} = \{d \in \mathcal{T}_X(\bar{x}) : c'(\bar{x}; d) \leq 0\}. \quad (3.4)$$

Then, \bar{x} is a B -stationary point of problem (1.1) if and only if \bar{x} solves the convex problems

$$\left\{ \begin{array}{l} \min_{x \in X} f_1(x) - [f_2(\bar{x}) + \langle s_{f_2}, x - \bar{x} \rangle] \\ \text{s.t.} \quad c_1(x) - [c_2(\bar{x}) + \langle s_{c_2}, x - \bar{x} \rangle] \leq 0 \end{array} \right\} \quad \forall s_{f_2} \in \partial^c f_2(\bar{x}), \forall s_{c_2} \in \partial c_2(\bar{x}). \quad (3.5)$$

Proof. Denote $\bar{Y}(\bar{x}) = \{x \in X : c_1(x) \leq c_2(\bar{x}) + \langle s_{c_2}, x - \bar{x} \rangle\}$. As the CQ (3.4) holds, Proposition 2.1 of [42] ensures that $\mathcal{T}_{X^c}(\bar{x}) = \mathcal{T}_{\bar{Y}(\bar{x})}(\bar{x}) = \text{c1} \{d \in \mathbb{R}^n : d = t(x - \bar{x}), x \in \bar{Y}(\bar{x}), t \in \mathbb{R}_+\}$. Thus, the B -stationary definition is equivalent to

$$\begin{aligned} f'(\bar{x}; x - \bar{x}) &\geq 0 && \forall x \in \bar{Y}(\bar{x}) \\ \Leftrightarrow f'_1(\bar{x}; x - \bar{x}) &\geq \langle s_{f_2}, x - \bar{x} \rangle && \forall s_{f_2} \in \partial^c f_2(\bar{x}), \forall x \in \bar{Y}(\bar{x}). \end{aligned} \quad (3.6)$$

The stated result follows upon establishing the equivalence between (3.6) and (3.5).

[(3.6) \Rightarrow (3.5)]. Suppose (3.6) holds and let $s_{c_2} \in \partial c_2(\bar{x})$ be arbitrary. As

$$Y(s_{c_2}) := \{x \in X : c_1(x) \leq c_2(\bar{x}) + \langle s_{c_2}, x - \bar{x} \rangle\} \subset \bar{Y}(\bar{x})$$

due to convexity of c_2 , we conclude that $f'_1(\bar{x}; x - \bar{x}) \geq \langle s_{f_2}, x - \bar{x} \rangle$ for all $s_{f_2} \in \partial^c f_2(\bar{x})$ and all $x \in Y(s_{c_2})$. Convexity of the latter set implies that \bar{x} minimizes $f_1(x) - \langle s_{f_2}, x - \bar{x} \rangle$ over $Y(s_{c_2})$ for all $s_{f_2} \in \partial^c f_2(\bar{x})$. Thus, condition (3.5) holds because $s_{c_2} \in \partial c_2(\bar{x})$ was taken arbitrarily.

[(3.5) \Rightarrow (3.6)]. To show the reverse implication, we proceed with a proof by contrapositive. Suppose that there exist $s'_{f_2} \in \partial^c f_2(\bar{x})$ and $x' \in \bar{Y}(\bar{x})$ such that $f'_1(\bar{x}; x' - \bar{x}) < \langle s'_{f_2}, x' - \bar{x} \rangle$ (and hence $x' \neq \bar{x}$), i.e., (3.6) does not hold. Let $s'_{c_2} \in \partial c_2(\bar{x})$ be such that $c'_2(\bar{x}; x' - \bar{x}) = \langle s'_{c_2}, x' - \bar{x} \rangle$. Therefore, x' is feasible for the convex problem

$$\begin{array}{l} \min_{x \in X} f_1(x) - [f_2(\bar{x}) + \langle s'_{f_2}, x - \bar{x} \rangle] \\ \text{s.t.} \quad c_1(x) - [c_2(\bar{x}) + \langle s'_{c_2}, x - \bar{x} \rangle] \leq 0. \end{array}$$

Together with our assumption $f'_1(\bar{x}; x' - \bar{x}) < \langle s'_{f_2}, x' - \bar{x} \rangle$, we have that $d = x' - \bar{x}$ is a feasible descent direction for the above problem, and thus \bar{x} can not be one of its solution. Hence, \bar{x} does not satisfy (3.5). The proof is thus complete. \square

Note that convexity of c_2 plays an important role in the above proposition. Indeed, if c_2 is nonconvex, then the set $\{x \in X : c_1(x) \leq c_2(\bar{x}) + \langle s_{c_2}, x - \bar{x} \rangle\}$ is not necessarily a subset of X^c and when solving the linearized subproblem (3.5) we may get a point that is infeasible for the original problem (1.1).

Example 3.3. Let $f_1 = x$, $f_2 = 0$, $c_1 = 0$, $c_2 = \frac{x^3}{3}$ and $X = [-2, 2]$. We are not in the framework of Proposition 3.2, since c_2 is not convex on $[-2, 2]$, but weakly convex (with modulus $\mu = 4$). At $\bar{x} = 0$, which globally solves (1.1), the convex problem (3.5) becomes $\min_{x \in [-2, 2]} x$ (because we have dropped the trivial constraint $0 \leq 0$), and thus does not provide a feasible point for the original problem. \square

However, if the modulus μ is known for the weakly convex function c_2 , adding a quadratic term in the constraint of the convex problem (3.5) makes the corresponding set feasible for the original problem. Moreover, Proposition 3.2 is generalized in case of weakly convex c_2 .

Corollary 3.4. *Let $c_2 : \mathcal{O} \rightarrow \mathbb{R}$ be a weakly convex function and $\mu_c \geq 0$ be a real number from Proposition 2.4 corresponding to c_2 . Moreover, assume that the CQ (3.4) holds. Then, \bar{x} is a B -stationary point of problem (1.1) if and only if, for any given $\mu \geq \mu_c$, \bar{x} solves the convex problems*

$$\left\{ \begin{array}{l} \min_{x \in X} f_1(x) - [f_2(\bar{x}) + \langle s_{f_2}, x - \bar{x} \rangle] \\ \text{s.t.} \quad c_1(x) - [c_2(\bar{x}) + \langle s_{c_2}, x - \bar{x} \rangle] + \frac{\mu}{2} \|x - \bar{x}\|^2 \leq 0 \end{array} \right\} \quad \forall s_{f_2} \in \partial^c f_2(\bar{x}), \forall s_{c_2} \in \partial^c c_2(\bar{x}). \quad (3.7)$$

Proof. Consider the convex functions $\tilde{c}_1(x) = c_1(x) + \frac{\mu}{2} \|x\|^2$ and $\tilde{c}_2(x) = c_2(x) + \frac{\mu}{2} \|x\|^2$ with $\mu \geq \mu_c$. The result follows from Proposition 3.2 by using instead the DC decomposition $c(x) = \tilde{c}_1(x) - \tilde{c}_2(x)$ and by noting that, for an arbitrary $\tilde{s}_{c_2} \in \partial \tilde{c}_2(\bar{x})$, we obtain

$$\tilde{c}_1(x) - [\tilde{c}_2(\bar{x}) + \langle \tilde{s}_{c_2}, x - \bar{x} \rangle] = c_1(x) - [c_2(\bar{x}) + \langle s_{c_2}, x - \bar{x} \rangle] + \frac{\mu}{2} \|x - \bar{x}\|^2$$

with $s_{c_2} = \tilde{s}_{c_2} - \mu \bar{x}$, $s_{c_2} \in \partial^c c_2(\bar{x})$. \square

Except for some particular cases, checking B -stationarity numerically is out of reach. Therefore, weaker stationarity concepts need to come into play: $\bar{x} \in X$ is said to be Clarke-stationary for (1.1) if there exists a Lagrange multiplier $\bar{\lambda}$ such that

$$\begin{cases} 0 \in \partial^c f(\bar{x}) + \bar{\lambda} \partial^c c(\bar{x}) + N_X(\bar{x}) \\ c(\bar{x}) \leq 0, \bar{\lambda} c(\bar{x}) = 0, \bar{\lambda} \geq 0, \bar{x} \in X. \end{cases} \quad (3.8)$$

Analogously, \bar{x} is a critical point of (1.1) if there exists a Lagrange multiplier $\bar{\lambda}$ such that

$$\begin{cases} 0 \in \partial f_1(\bar{x}) - \partial^c f_2(\bar{x}) + \bar{\lambda} [\partial c_1(\bar{x}) - \partial^c c_2(\bar{x})] + N_X(\bar{x}) \\ c(\bar{x}) \leq 0, \bar{\lambda} c(\bar{x}) = 0, \bar{\lambda} \geq 0, \bar{x} \in X. \end{cases} \quad (3.9)$$

Observe that if f_1 or f_2 and c_1 or c_2 are smooth, then criticality boils down to Clarke stationarity. Next, we revisit the proximal bundle method of [50] and extend it to the more general setting of problem (1.1). To this end, the method must be modified, and its convergence analysis must be done anew.

3.1 Problem reformulation via improvement function

Nonsmooth and nonconvex constraints in optimization problems are in general numerically treated via exact penalization [18, 21, 34, 14], linearization of certain components [27, 40], and improvement functions [31, 1, 50]. The latter has a recognized good practical performance, does not require the additional assumptions normally assumed in exact penalization methods, and employs parameters that are simple to set. For these reasons, we handle problem (1.1) via the *improvement function* $H : \mathcal{O} \times \mathcal{O} \rightarrow \mathbb{R}$ given by

$$H(x; y) = \max \{ f(x) - \tau_f(y), c(x) - \tau_c(y) \}, \quad (3.10a)$$

$$\text{with } \tau_f(y) = f(y) + \rho[c(y)]_+ \text{ and } \tau_c(y) = \sigma[c(y)]_+, \text{ for } \rho \geq 0 \text{ and } \sigma \in [0, 1]. \quad (3.10b)$$

Observe that if \bar{x} is a global solution of (1.1), then $H(x; \bar{x}) \geq 0$ for all $x \in X$ and $H(\bar{x}; \bar{x}) = 0$.

Improvement functions (also known as progress functions) have been considered within bundle methods in [31, 48, 39] for convex problems, in [1] for a class of (nonconvex) optimal control problems, and in [24, 50] for DC-constrained DC programs. In what follows we exploit some relevant mathematical properties of (3.10) and its link to the original problem (1.1). To this end, we need to consider necessary conditions for a point \bar{x} to be a local solution of the reformulated problem

$$\min_{x \in X} H(x; \bar{x}). \quad (3.11)$$

As the second argument of H is fixed, it follows from (3.2) that $\bar{x} \in X$ is a Clarke-stationary point of (3.11) if

$$0 \in \partial_1^c H(\bar{x}; \bar{x}) + N_X(\bar{x}), \quad (3.12)$$

where $\partial_1^c H$ stands for the generalized subdifferential of H with respect to the first argument. Proposition 2.1 ii) yields

$$\partial_1^c H(\bar{x}; \bar{x}) \subset \begin{cases} \partial^c c(\bar{x}) & \text{if } f(\bar{x}) - \tau_f(\bar{x}) < c(\bar{x}) - \tau_c(\bar{x}) \\ \text{conv} \{ \partial^c f(\bar{x}), \partial^c c(\bar{x}) \} & \text{if } f(\bar{x}) - \tau_f(\bar{x}) = c(\bar{x}) - \tau_c(\bar{x}) \\ \partial^c f(\bar{x}) & \text{if } f(\bar{x}) - \tau_f(\bar{x}) > c(\bar{x}) - \tau_c(\bar{x}). \end{cases}$$

Since we do not work with generalized subgradients of either f or c , but only with subgradients of the functions yielding their CwC decompositions (1.1b), we must consider a weaker stationary definition: we say that $\bar{x} \in X$ is a critical point of the composite problem (3.11) with CwC decompositions (1.1b) if

$$0 \in N_X(\bar{x}) + \begin{cases} \partial c_1(\bar{x}) - \partial^c c_2(\bar{x}) & \text{if } f(\bar{x}) - \tau_f(\bar{x}) < c(\bar{x}) - \tau_c(\bar{x}) \\ \text{conv} \{ \partial f_1(\bar{x}) - \partial^c f_2(\bar{x}), \partial c_1(\bar{x}) - \partial^c c_2(\bar{x}) \} & \text{if } f(\bar{x}) - \tau_f(\bar{x}) = c(\bar{x}) - \tau_c(\bar{x}) \\ \partial f_1(\bar{x}) - \partial^c f_2(\bar{x}) & \text{if } f(\bar{x}) - \tau_f(\bar{x}) > c(\bar{x}) - \tau_c(\bar{x}). \end{cases} \quad (3.13)$$

Note that if both f and c are regular, then the above condition coincides with that of Clarke stationarity: recall (3.2), Proposition 2.1 i), and observe that the set defined by the expressions in the curly brackets above is nothing but $\partial_1^c H(\bar{x}; \bar{x})$. The following result, inspired by both [1, Lemma 5.1] that deals with the (stronger) Clarke stationarity and [50, Thm. 2] that works with the (weaker) criticality definition from DC programming, links condition (3.13) with criticality of the original problem.

Theorem 3.5. *Let $\bar{x} \in X$ be a point satisfying condition (3.13). Then, the following hold:*

i) *If $c(\bar{x}) > 0$, then \bar{x} is a critical point (in the sense of (3.3)) of the optimization problem*

$$\min_{x \in X} c_1(x) - c_2(x). \quad (3.14)$$

ii) *If $c(\bar{x}) = 0$ and \bar{x} is not a critical point of (3.14), then \bar{x} satisfies (3.9) for some $\bar{\lambda} > 0$.*

iii) If $c(\bar{x}) < 0$, then \bar{x} satisfies (3.9) with $\bar{\lambda} = 0$.

Proof. With the τ function defined in (3.10b), note that

$$f(\bar{x}) - \tau_f(\bar{x}) - [c(\bar{x}) - \tau_c(\bar{x})] = \begin{cases} -[\rho + (1 - \sigma)]c(\bar{x}) < 0 & \text{if } c(\bar{x}) > 0, \\ 0 & \text{if } c(\bar{x}) = 0, \\ -c(\bar{x}) > 0 & \text{if } c(\bar{x}) < 0. \end{cases}$$

Hence,

$$\begin{aligned} c(\bar{x}) > 0 &\Leftrightarrow f(\bar{x}) - \tau_f(\bar{x}) < c(\bar{x}) - \tau_c(\bar{x}), \\ c(\bar{x}) = 0 &\Leftrightarrow f(\bar{x}) - \tau_f(\bar{x}) = c(\bar{x}) - \tau_c(\bar{x}), \\ c(\bar{x}) < 0 &\Leftrightarrow f(\bar{x}) - \tau_f(\bar{x}) > c(\bar{x}) - \tau_c(\bar{x}), \end{aligned}$$

and items i) and iii) follow directly from (3.13). To show item ii), recall that $c(\bar{x}) = 0$ and condition (3.13) ensures the existence of $\lambda \in [0, 1]$ such that

$$0 \in \lambda[\partial f_1(\bar{x}) - \partial^c f_2(\bar{x})] + (1 - \lambda)[\partial c_1(\bar{x}) - \partial^c c_2(\bar{x})] + N_X(\bar{x}).$$

By assumption, \bar{x} is not a critical point of (3.14). Then $0 \notin \partial c_1(\bar{x}) - \partial^c c_2(\bar{x}) + N_X(\bar{x})$, implying that λ above must be strictly positive. Dividing the displayed inclusion by $\lambda > 0$ we obtain the criticality condition (3.9) with $\bar{\lambda} = (1 - \lambda)/\lambda > 0$. \square

At item ii) above, the assumption that \bar{x} is not a critical point of (3.14) can be seen as a constraint qualification, which turns out to be more restrictive than (3.4). Indeed, the latter excludes d -stationary points of (3.14), but not necessarily critical ones. The following example gives a critical point \bar{x} of (3.14) that satisfies (3.4) but not the criticality condition (3.9) for the nonlinearly-constrained problem (1.1).

Example 3.6. Take $c_1(x) = \max\{x, 2x\}$, $c_2(x) = \max\{2x, 4x\}$, $X = [-2, 2]$ and $\bar{x} = 0$. Then $\mathcal{T}_X(\bar{x}) = \mathbb{R}$, $N_X(\bar{x}) = \{0\}$, and \bar{x} is a critical point of (3.14) because $0 \in \partial c_1(\bar{x}) - \partial c_2(\bar{x}) = [1, 2] - [2, 4] = [-3, 0]$. Furthermore, note that

$$c'_1(\bar{x}; d) = \max_{s \in [1, 2]} sd = \begin{cases} 2d & \text{if } d \geq 0 \\ d & \text{if } d \leq 0 \end{cases} \quad \text{and} \quad c'_2(\bar{x}; d) = \max_{s \in [2, 4]} sd = \begin{cases} 4d & \text{if } d \geq 0 \\ 2d & \text{if } d \leq 0, \end{cases}$$

thus $c'(\bar{x}; d) = \min\{-d, -2d\}$. We conclude that $\{d \in \mathcal{T}_X(\bar{x}) : c'(\bar{x}; d) < 0\} = \mathbb{R}_+$, whereas $\{d \in \mathcal{T}_X(\bar{x}) : c'(\bar{x}; d) \leq 0\} = \mathbb{R}_+ \cup \{0\}$, showing that $\bar{x} = 0$ satisfies the CQ (3.4). However, if we take $f_1(x) = 0$ and $f_2(x) = -\frac{1}{2}x^2 + x$, the following system does not have a solution:

$$\begin{cases} 0 & \in \partial f_1(0) - \partial^c f_2(0) + \bar{\lambda}[\partial c_1(0) - \partial c_2(0)] \\ \bar{\lambda} & \geq 0 \end{cases} \equiv \begin{cases} 0 & \in -1 + \bar{\lambda}[-3, 0] \\ \bar{\lambda} & \geq 0 \end{cases} \equiv \begin{cases} 0 & \in [-3\bar{\lambda} - 1, -1] \\ \bar{\lambda} & \geq 0, \end{cases}$$

i.e., \bar{x} does not satisfy (3.9). Figure 1 illustrates the objective and constraint function in this example: it is clear that \bar{x} is indeed a global maximizer of $f(x)$ under the constraints $x \in X$ and $c(x) \leq 0$. \square

This example shows that, at item ii) of Theorem 3.5, we cannot replace the assumption that \bar{x} is not a critical point of (3.14) with the CQ (3.4).

3.2 The DC setting.

In the DC setting, functions f_2 and c_2 are convex and the improvement function (3.10) is DC. Indeed, for \bar{x} fixed, we can write

$$H(x; \bar{x}) = F(x; \bar{x}) - G(x), \quad \text{with} \quad \begin{cases} F(x; \bar{x}) & = \max\{f_1(x) + c_2(x) - \tau_f(\bar{x}), f_2(x) + c_1(x) - \tau_c(\bar{x})\}, \\ G(x) & = f_2(x) + c_2(x). \end{cases} \quad (3.15)$$

Since F and G are convex functions, the criticality condition (3.3) for (3.11) (under this DC decomposition) reads as

$$0 \in \partial_1 F(\bar{x}; \bar{x}) - \partial G(\bar{x}) + N_X(\bar{x}), \quad (3.16)$$

where $\partial_1 F$ stands for the subdifferential of F with respect to the first argument. It turns out that our new condition (3.13) is stronger than (3.16), used in [50].

Lemma 3.7. *In addition to our basic assumptions on problem (1.1), suppose that f_2 and c_2 are convex. Then the necessary optimality condition (3.13) implies (3.16).*

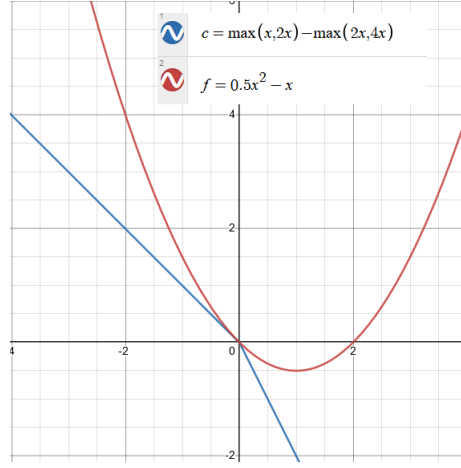


Figure 1: Function $f(x) = \frac{1}{2}x^2 - x$ in red and $c(x) = \max\{x, 2x\} - \max\{2x, 4x\}$ in blue.

Proof. Let $\bar{x} \in X$ be a point satisfying (3.13). Let us first observe that since f_2, c_2 are convex and thus regular, we have $\partial G(\bar{x}) = \partial f_2(\bar{x}) + \partial c_2(\bar{x})$. A similar observation can be made concerning the computation for F . Our analysis splits into three possible cases.

- a) $f(\bar{x}) - \tau_f(\bar{x}) < c(\bar{x}) - \tau_c(\bar{x})$. It follows from (3.13) that

$$0 \in N_X(\bar{x}) + \partial c_1(\bar{x}) - \partial c_2(\bar{x}) \subset N_X(\bar{x}) + \partial f_2(\bar{x}) + \partial c_1(\bar{x}) - [\partial f_2(\bar{x}) + \partial c_2(\bar{x})].$$

We claim that this inclusion implies (3.16). To see that, observe that the above inequality implies $f_1(\bar{x}) + c_2(\bar{x}) - \tau_f(\bar{x}) < f_2(\bar{x}) + c_1(\bar{x}) - \tau_c(\bar{x})$, which in turn gives $\partial_1 F(\bar{x}; \bar{x}) = \partial f_2(\bar{x}) + \partial c_1(\bar{x})$. Therefore, the right-hand-side of the above inclusion is (3.16).

- b) $f(\bar{x}) - \tau_f(\bar{x}) = c(\bar{x}) - \tau_c(\bar{x})$. It follows from (3.13) that there exists $\lambda \in [0, 1]$ such that

$$\begin{aligned} 0 &\in N_X(\bar{x}) + \lambda[\partial f_1(\bar{x}) - \partial f_2(\bar{x})] + (1 - \lambda)[\partial c_1(\bar{x}) - \partial c_2(\bar{x})] \\ &= N_X(\bar{x}) + \lambda[\partial f_1(\bar{x}) + \partial c_2(\bar{x})] - \lambda\partial f_2(\bar{x}) + (1 - \lambda)\partial c_1(\bar{x}) - \partial c_2(\bar{x}) \\ &\subset N_X(\bar{x}) + \lambda[\partial f_1(\bar{x}) + \partial c_2(\bar{x})] + (1 - \lambda)[\partial f_2(\bar{x}) + \partial c_1(\bar{x})] - [\partial f_2(\bar{x}) + \partial c_2(\bar{x})] \\ &\subset N_X(\bar{x}) + \partial_1 F(\bar{x}; \bar{x}) - \partial G(\bar{x}). \end{aligned}$$

- c) $f(\bar{x}) - \tau_f(\bar{x}) > c(\bar{x}) - \tau_c(\bar{x})$. Again, (3.13) gives

$$0 \in N_X(\bar{x}) + \partial f_1(\bar{x}) - \partial f_2(\bar{x}) \subset N_X(\bar{x}) + \partial f_1(\bar{x}) + \partial c_2(\bar{x}) - [\partial f_2(\bar{x}) + \partial c_2(\bar{x})].$$

The proof is complete because in this case $\partial_1 F(\bar{x}; \bar{x}) = \partial f_1(\bar{x}) + \partial c_2(\bar{x})$ due to the fact that $f_1(\bar{x}) + c_2(\bar{x}) - \tau_f(\bar{x}) > f_2(\bar{x}) + c_1(\bar{x}) - \tau_c(\bar{x})$. □

Remark 3.8. In the DC setting, the three concepts of criticality (3.12), (3.13), and (3.16) are equivalent when f_2 and c_2 are continuously differentiable at \bar{x} . Indeed, in this case f and c are regular at \bar{x} and (3.12) coincides with (3.13) (regardless of convexity of f_2 and c_2). Theorem 2 in [50] ensures that, under these assumptions, (3.16) is equivalent to (3.12). □

The following example shows that (3.16) does not necessarily imply (3.13) in the nondifferentiable DC case.

Example 3.9. Let $X = [-1, 1]$, $f_1 = 2x$, $f_2 = |x|$, $c_1 = 4x$ and $c_2 = 2|x|$. At $\bar{x} = 0$, $f(\bar{x}) = c(\bar{x}) = 0$ and thus, $\tau_f(\bar{x}) = \tau_c(\bar{x}) = 0$ due to (3.10b). Furthermore, we have that

$$\partial f_1(0) = \{2\}, \quad \partial f_2(0) = [-1, 1], \quad \partial c_1(0) = \{4\}, \quad \text{and} \quad \partial c_2(0) = [-2, 2].$$

As a result, $\partial f_1(0) - \partial f_2(0) = [1, 3]$, $\partial c_1(0) - \partial c_2(0) = [2, 6]$, $\partial f_1(0) + \partial c_2(0) = [0, 4]$, and $\partial f_2(0) + \partial c_1(0) = [3, 5]$. As in $N_X(0) = \{0\}$, we conclude that

$$0 \notin [1, 6] = \text{conv}\{\partial f_1(0) - \partial f_2(0), \partial c_1(0) - \partial c_2(0)\} + N_X(0),$$

whereas

$$\text{conv}\{\partial f_1(0) + \partial c_2(0), \partial c_1(0) + \partial f_2(0)\} + N_X(0) = [0, 5]$$

and $\partial f_2(0) + \partial c_2(0) = [-3, 3]$, showing that (3.16) is satisfied but not (3.13). □

The paper [50] proposes a bundle method for DC-constrained DC programs employing the DC decomposition $H = F - G$ above. Once a critical point satisfying (3.16) is computed, the link with criticality of the original problem is adequate if f_2 and c_2 are continuously differentiable at \bar{x} . In the next section we modify that method to compute a point satisfying the stronger criticality condition (3.13). As a result, the link with criticality of the original problem is nicely established by Theorem 3.5 without any additional assumption. In fact, f_2 and c_2 need not even be convex, but weakly convex on some neighbourhood of each $x \in \mathcal{O}$. We, therefore, strengthen the analysis provided in [50] even though significantly fewer assumptions are required: [50] works in the DC configuration, whereas here, we deal with the more general CwC structure. These improvements, together with the optimality conditions presented above, feature the main contributions of this work.

4 Proximal bundle method with improvement function

This section extends the proximal bundle method of [50] for computing a critical point of problem (1.1). The main tool in our analysis is the improvement function H given in (3.10). In the DC setting, the algorithm of [50] works with the explicit DC decomposition (3.15) of H and computes a point $\bar{x} \in X$ satisfying the classic criticality condition in DC programming (3.16). In this section we do not decompose H and consider the milder assumption that f_2 and c_2 are weakly convex and target the stronger criticality condition (3.13).

4.1 The method's main ingredients: model, subproblem, and descent test

The algorithm requires four oracles (black-boxes) providing, for every given $x \in X$, $i \in \{1, 2\}$, the function values $f_i(x)$, $c_i(x)$, arbitrary subgradients $s_{f_1} \in \partial f_1(x)$, $s_{c_1} \in \partial c_1(x)$ (c.f., (2.1)) and arbitrary generalized subgradients $s_{f_2} \in \partial^c f_2(x)$, $s_{c_2} \in \partial^c c_2(x)$ (c.f. (2.2)). We do not impose any assumption on these (generalized) subgradients, as they are assumed to be computed by (external) oracles that do not accept any intervention from the algorithm. (This is particularly useful in industrial applications where companies do not want or cannot share information on the underlying functions with optimizers.)

At iteration $k \in \mathbb{N}$, given a trial point $x^k \in X$, we construct a linearization of every component (here $s_{f_i}^k$, $s_{c_i}^k$, $i \in \{1, 2\}$), denote the respective - generalized - subgradients at x^k :

$$\bar{f}_i^k(x) := f_i(x^k) + \langle s_{f_i}^k, x - x^k \rangle \quad (i = 1, 2) \quad (4.1a)$$

$$\bar{c}_i^k(x) := c_i(x^k) + \langle s_{c_i}^k, x - x^k \rangle \quad (i = 1, 2). \quad (4.1b)$$

Due to convexity of f_1 and c_1 , we have the following inequalities

$$\bar{f}_1^k(x) \leq f_1(x) \quad \text{and} \quad \bar{c}_1^k(x) \leq c_1(x) \quad \text{for all } x \in \mathbb{R}^n. \quad (4.2)$$

Since X is compact and components f_2 and c_2 are assumed to be only weakly convex on some neighbourhood of each $x \in \mathcal{O}$, we have weaker inequalities for these functions. Let $\mathcal{O}' \subset \mathcal{O}$ be an open convex set and μ_f, μ_c real numbers ensured by Proposition 2.4. As $x^k \in X \subset \mathcal{O}'$, the following inequalities are due to Proposition 2.4, item ii)

$$\bar{f}_2^k(x) \leq f_2(x) + \frac{\bar{\mu}}{2} \|x - x^k\|^2 \quad \text{and} \quad \bar{c}_2^k(x) \leq c_2(x) + \frac{\bar{\mu}}{2} \|x - x^k\|^2 \quad \text{for all } x \in X, \quad (4.3)$$

where $\bar{\mu} := \max\{\mu_f, \mu_c\} > 0$. Observe that the threshold $\bar{\mu}$ is in general unknown, and the inequalities in (4.3) are only supposed to hold for x in X , in contrast with the (subgradient) inequalities in (4.2).

Let \mathcal{B}_f^k and \mathcal{B}_c^k be two index sets gathering the *bundle* of information (function values and subgradients) given by the oracles. In general, $\mathcal{B}_f^k, \mathcal{B}_c^k \subset \{0, \dots, k\}$ but other possibilities exist making it possible to design a limited-memory method (see Remark 5.5 below). These index sets are useful to define the following individual cutting-plane models for the convex functions f_1 and c_1 :

$$\check{f}_1^k(x) := \max_{j \in \mathcal{B}_f^k} \bar{f}_1^j(x) \leq f_1(x) \quad \text{for all } x \in \mathbb{R}^n$$

$$\check{c}_1^k(x) := \max_{j \in \mathcal{B}_c^k} \bar{c}_1^j(x) \leq c_1(x) \quad \text{for all } x \in \mathbb{R}^n.$$

Furthermore, let $\ell_k \in \{0, \dots, k\}$ be the iteration index of the best candidate solution (*stability center*, in the parlance of bundle methods) among the trial points $\{x^0, \dots, x^k\}$: whenever a better candidate solution x^{k+1} is computed by the algorithm, at a so-called *serious step*, such a point becomes the new stability center

and the counter ℓ is increased by one: for $\kappa \in (0, \frac{1}{2})$, we declare a *serious step* and let $\ell_{k+1} := k + 1$ if $x^{k+1} \neq x^{\ell_k}$ and the inequality

$$H(x^{k+1}; x^{\ell_k}) \leq H(x^{\ell_k}; x^{\ell_k}) - \frac{\kappa}{2} \|x^{k+1} - x^{\ell_k}\|^2 \quad (4.5)$$

holds, and declare a *null step* and let $\ell_{k+1} := \ell_k$ otherwise. Since the descent test is independent of the model, the following result from [50] also holds in our framework.

Lemma 4.1 (Lemma 1 in [50]). *Let $x^{\ell_k} \in X$ be the stability center at iteration k . Then $H(x^{\ell_k}; x^{\ell_k}) \geq 0$ and if inequality (4.5) holds, we have that either*

- i) $f(x^{k+1}) \leq f(x^{\ell_k}) - \frac{\kappa}{2} \|x^{k+1} - x^{\ell_k}\|^2$ and $c(x^{k+1}) \leq 0$ when $c(x^{\ell_k}) \leq 0$; or
- ii) $c(x^{k+1}) \leq c(x^{\ell_k}) - \frac{\kappa}{2} \|x^{k+1} - x^{\ell_k}\|^2$ when $c(x^{\ell_k}) > 0$.

The rationale of serious iterates is to ensure sufficient decrease on one component function of $H(\cdot; x^{\ell_k})$ while maintaining feasibility for (1.1) once reached. Having all these ingredients at our disposal, we can now define our convex model for the improvement function (3.10) at iteration k :

$$\check{H}^k(x; x^{\ell_k}) = \max \left\{ \check{f}_1^k(x) - \check{f}_2^k(x) - \tau_f(x^{\ell_k}), \check{c}_1^k(x) - \check{c}_2^k(x) - \tau_c(x^{\ell_k}) \right\}. \quad (4.6)$$

(Even in the particular setting where f_2 and c_2 are convex functions, this model differs from the one employed in [50] and is crucial to obtain convergence results stronger than the ones in that paper.) Given a prox-parameter $\mu^k > 0$ estimating the threshold $\bar{\mu}$ in (4.3), the next iterate is the solution of the following strict convex subproblem

$$x^{k+1} = \arg \min_{x \in X} \check{H}^k(x; x^{\ell_k}) + \frac{\mu^k}{2} \|x - x^{\ell_k}\|^2, \quad (4.7)$$

which can be transformed into a QP (provided X is a polyhedron) by adding an extra variable $r \in \mathbb{R}$

$$\begin{cases} \min_{x, r} & r + \frac{\mu^k}{2} \|x - x^{\ell_k}\|^2 \\ \text{s.t.} & \check{f}_1^j(x) - \check{f}_2^k(x) - r \leq \tau_f(x^{\ell_k}) \quad \forall j \in \mathcal{B}_f^k \\ & \check{c}_1^j(x) - \check{c}_2^k(x) - r \leq \tau_c(x^{\ell_k}) \quad \forall j \in \mathcal{B}_c^k \\ & x \in X, \quad r \in \mathbb{R}. \end{cases} \quad (4.8)$$

The optimality condition for (4.7) gives

$$x^{k+1} = x^{\ell_k} - \frac{1}{\mu^k} [p^{k+1} + s_X^{k+1}], \quad \text{with} \quad \begin{cases} p^{k+1} \in \partial_1 \check{H}^k(x^{k+1}; x^{\ell_k}) \\ s_X^{k+1} \in N_X(x^{k+1}). \end{cases} \quad (4.9)$$

As usual in bundle methods, we may remove from the model the inactive linearizations to keep (4.8) small. To this end, we denote by $\check{\mathcal{B}}_f^k \subset \mathcal{B}_f^k$ and $\check{\mathcal{B}}_c^k \subset \mathcal{B}_c^k$ the index set of active linearizations in the QP subproblem (4.8), i.e.,

$$\check{\mathcal{B}}_f^k := \{j \in \mathcal{B}_f^k : \alpha_f^j > 0\} \quad \text{and} \quad \check{\mathcal{B}}_c^k := \{j \in \mathcal{B}_c^k : \alpha_c^j > 0\} \quad (4.10)$$

where $\alpha_f^j \geq 0$, $j \in \mathcal{B}_f^k$, denote the Lagrange multipliers associated with the first set of constraints and $\alpha_c^j \geq 0$, $j \in \mathcal{B}_c^k$, the ones associated with the second family of constraints. We mention in passing that the index sets \mathcal{B}_f^k and \mathcal{B}_c^k can be kept bounded at the price of including artificial (aggregate) linearizations. We postpone this discussion to Remark 5.5, right after the analysis of null steps (the only place in the convergence analysis where bundle management plays a role.)

We can now present the following proximal bundle method algorithm for CwC-constrained CwC programs (1.1), which modifies [50, Alg. 1] in two ways. First, the convex model (4.6) for the improvement function is distinct. On the one hand, it is a key element to obtain the stronger criticality condition (3.13), and on the other hand, it leads to a simpler/smaller strongly convex QP (4.8) (more details can be found in Subsection 4.2 below). Second, Algorithm 1 employs an ad-hoc rule to update the proximal parameter μ^k so that no pre-estimation of the underlying weakly-convex moduli is needed. The proposed rule employs the following value

$$\nu^k := 2 \max \left\{ \frac{\check{f}_2^k(x^{k+1}) - f_2(x^{k+1})}{\|x^{k+1} - x^{\ell_k}\|^2}, \frac{\check{c}_2^k(x^{k+1}) - c_2(x^{k+1})}{\|x^{k+1} - x^{\ell_k}\|^2}, 0 \right\}. \quad (4.11)$$

(For numerical performance it can be preferable to replace the last term (equal to zero) in (4.11) by a small machine epsilon.)

Algorithm 1 Proximal Bundle Method for CwC-constrained CwC programs - CwC-PBM

- Step 0 (Initialization)** Let $x^0 \in X$, $\kappa \in (0, \frac{1}{2})$, $\kappa \leq \mu^0$, $\rho \geq 0$, $\sigma \in [0, 1)$, and $\text{To1} \geq 0$ be given.
 Call the oracles to compute $f_i(x^0)$, $c_i(x^0)$, and (generalized) subgradients $s_{f_i}^0$, $s_{c_i}^0$, $i = 1, 2$.
 Define $k := \ell_k = 0$ and $\mathcal{B}_f^0 = \mathcal{B}_c^0 := \{0\}$.
- Step 1 (Trial point)** Compute x^{k+1} by solving the QP (4.8).
- Step 2 (Stopping test)** If $\|x^{k+1} - x^{\ell_k}\| \leq \text{To1}$, then stop and return x^{ℓ_k} .
- Step 3 (Oracles call)** Compute $f_i(x^{k+1})$, $c_i(x^{k+1})$, and subgradients $s_{f_i}^{k+1}$, $s_{c_i}^{k+1}$, $i = 1, 2$.
- Step 4 (Descent test)**
- (a) If (4.5) holds, then declare a *serious step*: define $\ell_{k+1} := k + 1$, choose $\mathcal{B}_f^{k+1}, \mathcal{B}_c^{k+1} \subset \{0, \dots, k + 1\}$ with $\{k + 1\} \in \mathcal{B}_f^{k+1} \cap \mathcal{B}_c^{k+1}$ and arbitrarily select $\mu^{k+1} \in (0, \mu^k]$.
 - (b) Else, declare a *null step*: define $\ell_{k+1} := \ell_k$ and choose $\mathcal{B}_f^{k+1}, \mathcal{B}_c^{k+1} \subset \{0, \dots, k + 1\}$ with $\bar{\mathcal{B}}_f^k \cup \{k + 1, \ell_k\} \subset \mathcal{B}_f^{k+1}$ and $\bar{\mathcal{B}}_c^k \cup \{k + 1, \ell_k\} \subset \mathcal{B}_c^{k+1}$ ($\bar{\mathcal{B}}_f^k$ and $\bar{\mathcal{B}}_c^k$ as in (4.10)).
 Compute ν^k by (4.11). If $\nu^k \geq \mu^k - 2\kappa$, set $\mu^{k+1} = \nu^k + 1$; otherwise $\mu^{k+1} = \mu^k$.
- Step 5 (Loop)** Set $k := k + 1$ and go back to Step 1.
-

A drawback of the rule for updating the prox-parameter is that μ^k only increases after a null step when the inequality $\nu^k \geq \mu^k - 2\kappa$ is verified. As a result, μ^k may never increase: this is, for instance, the case when f_2 and c_2 are convex (thus $\nu^k = 0$ for all k). The motivation for this rule is to eventually keep the prox-parameter fixed if the algorithm performs an infinite sequence of null steps after a last serious step (see Lemma 5.1). This is a condition necessary to prove Proposition 5.4 below. We care to mention that increasing μ^k after a null step is a simple strategy that pays off in practice: it helps the algorithm to either stop or produce a new serious step, and thus accelerate the numerical performance.

4.2 The DC setting: a comparison with the earlier bundle method for DC programs

In the DC setting, both functions f_2 and c_2 are convex and the improvement function (3.10) is DC. The DC decomposition given in (3.15), with \bar{x} replaced with x^{ℓ_k} , was exploited in the bundle method of [50] through the following model for the improvement function $H(\cdot; x^{\ell_k})$ (see Eq. (18) therein):

$$\max \left\{ \check{f}_1^k(x) + \check{c}_2^k(x) - \tau_f(x^{\ell_k}), \check{f}_2^k(x) + \check{c}_1^k(x) - \tau_c(x^{\ell_k}) \right\} - [\check{f}_2^{\ell_k}(x) + \check{c}_2^{\ell_k}(x)]. \quad (4.12)$$

Differently from our model (4.6), the above gathers also cutting-planes for f_2 and c_2 and, although gathering more information, only the weaker criticality condition (3.16) is ensured by the method of [50]. Hence, the proposed model (4.6) is more advantageous than (4.12) from both practical and theoretical point of view:

- the quadratic program (QP) issued by our model has only half of the linearizations, and is thus simpler to solve;
- convexity of f_2 and c_2 are required in (4.12), but not in (4.6);
- both models (4.6) and (4.12) are iteratively updated to ensure that every cluster point $\bar{x} \in X$ of the sequence of stability centers satisfies a criticality condition. To show that such a point is also critical for (the DC counterpart of) (1.1), [50, Thm. 2] requires both f_2 and c_2 to be continuously differentiable at \bar{x} . As we will show in Theorem 5.8 below, neither convexity nor differentiability of f_2 and c_2 are required to establish that \bar{x} issued by Algorithm 1 is also critical for (1.1) in the sense of (3.9). Thus, Algorithm 1 strengthens the results of [50] even though significantly fewer assumptions are required.

Although the apparently small changes concerning [50, Alg. 1], the convergence analysis in that paper cannot be reused here. The reason is that the analysis in [50] strongly depends on the DC decomposition

of the employed model for the improvement function. That reasoning is no longer valid for our new model, even if f_2 and c_2 were convex. Furthermore, our more general setting requires extra steps to cope with the weakly convex functions.

5 Convergence Analysis

The goal of this section is to show that every cluster point \bar{x} of the sequence $\{x^{\ell_k}\}_k \subset X$ generated by Algorithm 1 satisfies the necessary optimality condition (3.13). To this end, we first observe that the sequence of prox-parameters issued by Algorithm 1 is bounded.

Lemma 5.1. *The value $\mu_{\max} := \sup_{k \in \mathbb{N}} \mu^k$ is finite. Furthermore, if the algorithm produces an infinite sequence of null steps after a last serious step, then the prox-parameter becomes eventually constant.*

Proof. Let $\bar{\mu} := \max\{\mu_{f_2}, \mu_{c_2}, \mu^0\} > 0$ be given, where μ_{f_2} and μ_{c_2} are as in Proposition 2.4 for the weakly convex functions f_2 and c_2 , and μ^0 is the parameter given to the algorithm at initialization. Then, by taking $y := x^{k+1}$ and $x := x^{\ell_k}$ in (2.4) it follows that

$$2 \frac{\bar{f}_2^{\ell_k}(x^{k+1}) - f_2(x^{k+1})}{\|x^{k+1} - x^{\ell_k}\|^2} \leq \bar{\mu}, \quad \text{and} \quad 2 \frac{\bar{c}_2^{\ell_k}(x^{k+1}) - c_2(x^{k+1})}{\|x^{k+1} - x^{\ell_k}\|^2} \leq \bar{\mu} \quad \text{for all } k \text{ with } x^{k+1} \neq x^{\ell_k}.$$

As a result, $\nu^k \leq \bar{\mu}$ for all k . Note that the prox-parameter is only increased after a null step such that $\nu^k \geq \mu^k - 2\kappa$. In this case, the rule employed in Step 4 of the algorithm sets $\mu^{k+1} = \nu^k + 1$, which gives $\mu^{k+1} = \nu^k + 1 \leq \bar{\mu} + 1$. Since the algorithm does not increase the prox-parameter after a serious step or null step such that $\nu^k < \mu^k - 2\kappa$, we conclude that $\mu_{\max} := \sup_{k \in \mathbb{N}} \mu^k \leq \bar{\mu} + 1$ is finite. Finally, note that the prox-parameter is sharply increased after a null step such that $\nu^k \geq \mu^k - 2\kappa$: $\mu^{k+1} = \nu^k + 1 \geq \mu^k - 2\kappa + 1 > \mu^k + \delta$ because $\kappa \in (0, \frac{1}{2})$, with $\delta = \frac{1}{2} - \kappa > 0$. As a result, if the algorithm produces an infinite sequence of null steps after a last serious step, then the inequality $\nu^k < \mu^k - 2\kappa$ will be satisfied for all k large enough and the prox-parameter will become constant (otherwise μ^k would increase indefinitely, which contradicts that μ_{\max} is finite). \square

We now define the following function $\bar{H} : \mathcal{O} \times \mathcal{O} \rightarrow \mathbb{R}$, which is of key importance in our analysis:

$$\bar{H}(x; y) := \max \{ f_1(x) - [f_2(y) + \langle s_{f_2}, x - y \rangle] - \tau_f(y), c_1(x) - [c_2(y) + \langle s_{c_2}, x - y \rangle] - \tau_c(y) \}, \quad (5.1)$$

with $s_{f_2} \in \partial^c f_2(y)$ and $s_{c_2} \in \partial^c c_2(y)$. As these subgradients are not specified, the above definition is ambiguous. However, when y is a point previously computed by the algorithm, say $y = x^j$ for $j \leq k$, then $s_{f_2}^j \in \partial^c f_2(x^j)$ and $s_{c_2}^j \in \partial^c c_2(x^j)$ are the subgradients provided by the oracles and ambiguity disappears:

$$\bar{H}(x; x^j) := \max \{ f_1(x) - \bar{f}_2^j(x) - \tau_f(x^j), c_1(x) - \bar{c}_2^j(x) - \tau_c(x^j) \}.$$

It follows from convexity of f_1 and c_1 that, for every $y \in \mathcal{O}$ fixed, the function $\bar{H}(\cdot; y)$ is convex and satisfies $\bar{H}(\cdot; x^{\ell_k}) \geq \check{H}^k(\cdot; x^{\ell_k})$ for all k . Furthermore, as $\ell_k \in \mathcal{B}_f^k \cap \mathcal{B}_c^k$ for all k , we have that $\bar{f}_i^{\ell_k}(x^{\ell_k}) = f_i(x^{\ell_k})$, $\bar{c}_i^{\ell_k}(x^{\ell_k}) = c_i(x^{\ell_k})$, $i = 1, 2$, and thus

$$\bar{H}(x^{\ell_k}; x^{\ell_k}) = \check{H}^k(x^{\ell_k}; x^{\ell_k}) = H(x^{\ell_k}; x^{\ell_k}). \quad (5.2)$$

The following lemma is of particular interest in the remainder of this work.

Lemma 5.2. *Suppose that \bar{x} minimizes $\bar{H}(\cdot; \bar{x})$ over X . Then, \bar{x} satisfies the necessary optimality condition (3.13).*

Proof. Convexity of $\bar{H}(\cdot; \bar{x})$ in the first argument and assumption on $\bar{x} \in X$ imply that $0 \in \partial_1 \bar{H}(\bar{x}; \bar{x}) + N_X(\bar{x})$. The result follows by noting that, for some pair of generalized subgradients $\bar{s}_{f_2} \in \partial^c f_2(\bar{x})$ and $\bar{s}_{c_2} \in \partial^c c_2(\bar{x})$, the following set

$$\partial_1 \bar{H}(\bar{x}; \bar{x}) = \begin{cases} \partial c_1(\bar{x}) - \bar{s}_{c_2} & \text{if } f(\bar{x}) - \tau_f(\bar{x}) < c(\bar{x}) - \tau_c(\bar{x}) \\ \text{conv} \{ \partial f_1(\bar{x}) - \bar{s}_{f_2}, \partial c_1(\bar{x}) - \bar{s}_{c_2} \} & \text{if } f(\bar{x}) - \tau_f(\bar{x}) = c(\bar{x}) - \tau_c(\bar{x}) \\ \partial f_1(\bar{x}) - \bar{s}_{f_2} & \text{if } f(\bar{x}) - \tau_f(\bar{x}) > c(\bar{x}) - \tau_c(\bar{x}) \end{cases}$$

is contained in the one defined by the curly brackets in (3.13). \square

We begin the convergence analysis for the case $\text{To1} = 0$ with the remark that the sequence of stability centers $\{x^{\ell_k}\}_k$ has at least one cluster point, since it is contained in the compact set X . We split the analysis into three cases: the algorithm performs only finitely many steps; the algorithm performs infinitely many steps and the sequence $\{x^{\ell_k}\}_k$ is either finite or infinite.

Proposition 5.3 (Finitely many iterations). *Assume that Algorithm 1 stops at iteration k with $\text{ToI} = 0$. Then, the last stability center $\bar{x} := x^{\ell_k} = x^{k+1}$ satisfies condition (3.13).*

Proof. It follows from the model's definition (4.6) and (5.1) that $\bar{H}(x; x^{\ell_k}) \geq \check{H}^k(x; x^{\ell_k})$ for all $x \in \mathcal{O}$. Hence, as $x^{\ell_k} \in X$ we have that

$$\begin{aligned} \bar{H}(x^{\ell_k}; x^{\ell_k}) &\geq \min_{x \in X} \bar{H}(x; x^{\ell_k}) + \frac{\mu_k}{2} \|x - x^{\ell_k}\|^2 \\ &\geq \min_{x \in X} \check{H}^k(x; x^{\ell_k}) + \frac{\mu_k}{2} \|x - x^{\ell_k}\|^2 \\ &= \check{H}^k(x^{k+1}; x^{\ell_k}) + \frac{\mu_k}{2} \|x^{k+1} - x^{\ell_k}\|^2 \\ &= \check{H}^k(x^{\ell_k}; x^{\ell_k}) = \bar{H}(x^{\ell_k}; x^{\ell_k}), \end{aligned}$$

where the first equality is due to (4.7), the second one follows by the fact that $x^{k+1} = x^{\ell_k}$ since the algorithm stops at iteration k with $\text{ToI} = 0$, and the last one is due to (5.2). Hence, x^{ℓ_k} minimizes $\bar{H}(\cdot; x^{\ell_k}) + \frac{\mu_k}{2} \|\cdot - x^{\ell_k}\|^2$ over X and the quadratic term vanishes in the corresponding optimality condition: $0 \in \partial_1 \bar{H}(\bar{x}; x^{\ell_k}) + N_X(\bar{x})$ and the stated result follows from Lemma 5.2. \square

If the algorithm performs finitely many serious steps and infinite number of null steps, the following result shows that the last stability center satisfies (3.13).

Proposition 5.4 (Finitely many serious steps). *Suppose that Algorithm 1 with $\text{ToI} = 0$ does not stop but produces only finitely many serious steps. Then the last stability center \bar{x} satisfies the condition (3.13), and $\lim_{k \rightarrow \infty} x^k = \bar{x}$.*

Proof. Let $\ell \in \mathbb{N}$ denote the last serious iteration, then $\bar{x} = x^\ell$ and note that, for all subsequent (null) iterations $k > \ell$, $\ell_k = \ell$ and the linearizations \bar{f}_2^ℓ and \bar{c}_2^ℓ are fixed in the model $\check{H}^k(\cdot; \bar{x})$, which is in this case a cutting-plane model for the convex function $\check{H}(\cdot; \bar{x})$. Here we take $\tau_f^\ell = \tau_f(x^\ell)$, $\tau_c^\ell = \tau_c(x^\ell)$, and function $\bar{H}(\cdot; \bar{x})$ defined with the fixed linearizations \bar{f}_2^ℓ and \bar{c}_2^ℓ , i.e.,

$$\bar{H}(\cdot; x^\ell) := \max \left\{ f_1(\cdot) - \bar{f}_2^\ell(\cdot) - \tau_f^\ell, c_1(\cdot) - \bar{c}_2^\ell(\cdot) - \tau_c^\ell \right\}.$$

We highlight that the updating rule for μ^k in Algorithm 1 ensures that the sequence $\{\mu^k\}_{k > \ell}$ is non-decreasing and becomes constant at a certain value $\mu' \in (0, \mu_{\max}]$ after finitely many steps $k' > \ell$, as a consequence of Lemma 5.1. More precisely, the updating rule at Step 4(b) of Algorithm 1 ensures that

$$\mu^k = \mu' \quad \text{and} \quad \nu^k + 2\kappa < \mu' \quad \text{for all } k > k'. \quad (5.3)$$

Hence, from iteration k' on, Algorithm 1 becomes a cutting-plane procedure to compute the unique solution \tilde{x} of

$$\min_{x \in X} \bar{H}(x; \bar{x}) + \frac{\mu'}{2} \|x - \bar{x}\|^2. \quad (5.4)$$

As the algorithm keeps all the active linearizations in the bundles (Step 4(b)), standard arguments from the convex bundle methods' theory (see [5, Prop. 4.3]) ensure that

$$\lim_{k \rightarrow \infty} x^k = \tilde{x} \quad \text{and} \quad \lim_{k \rightarrow \infty} [\check{H}^k(x^{k+1}; \bar{x}) - \bar{H}(x^{k+1}; \bar{x})] = 0.$$

(The last inequality implies that the convex model asymptotically coincides with the function at the limit point.) We claim that $\tilde{x} = \bar{x}$. To show that, let us assume the opposite, i.e., $\tilde{x} \neq \bar{x}$, and arrive to a contradiction. In this case, for some $\delta > 0$, we may find an index k_1 such that $\|x^{k+1} - \bar{x}\|^2 > \delta$ for all $k \geq k_1$. We may furthermore find an index k_2 such that $\check{H}^k(x^{k+1}; \bar{x}) - \bar{H}(x^{k+1}; \bar{x}) \geq -\frac{\kappa}{2}\delta$ for all $k \geq k_2$ as the left-hand side vanishes. Therefore, for $k'' \geq \max\{k_1, k_2, k'\}$, we have

$$\check{H}^k(x^{k+1}; \bar{x}) - \bar{H}(x^{k+1}; \bar{x}) \geq -\frac{\kappa}{2} \|x^{k+1} - \bar{x}\|^2 \neq 0 \quad \text{for all } k > k''.$$

The following chain of inequalities holds at every iteration $k > k''$:

$$\begin{aligned}
\check{H}^k(\bar{x}; \bar{x}) &\geq \check{H}^k(x^{k+1}; \bar{x}) + \frac{\mu'}{2} \|x^{k+1} - \bar{x}\|^2 && \text{(by (4.7) and (5.3))} \\
&= [\check{H}^k(x^{k+1}; \bar{x}) - \bar{H}(x^{k+1}; \bar{x})] + \bar{H}(x^{k+1}; \bar{x}) + \frac{\mu'}{2} \|x^{k+1} - \bar{x}\|^2 \\
&\geq -\frac{\kappa}{2} \|x^{k+1} - \bar{x}\|^2 + \bar{H}(x^{k+1}; \bar{x}) + \frac{\mu'}{2} \|x^{k+1} - \bar{x}\|^2 \\
&\geq -\frac{\kappa}{2} \|x^{k+1} - \bar{x}\|^2 + \max \left\{ \begin{array}{l} f_1(x^{k+1}) - \bar{f}_2^\ell(x^{k+1}) - \tau_f^\ell + \frac{\mu'}{2} \|x^{k+1} - \bar{x}\|^2 \\ c_1(x^{k+1}) - \bar{c}_2^\ell(x^{k+1}) - \tau_c^\ell + \frac{\mu'}{2} \|x^{k+1} - \bar{x}\|^2 \end{array} \right\} && \text{(by (5.1))} \\
&> -\frac{\kappa}{2} \|x^{k+1} - \bar{x}\|^2 + \max \left\{ \begin{array}{l} f_1(x^{k+1}) - \bar{f}_2^\ell(x^{k+1}) - \tau_f^\ell + \frac{\nu^k + 2\kappa}{2} \|x^{k+1} - \bar{x}\|^2 \\ c_1(x^{k+1}) - \bar{c}_2^\ell(x^{k+1}) - \tau_c^\ell + \frac{\nu^k + 2\kappa}{2} \|x^{k+1} - \bar{x}\|^2 \end{array} \right\} && \text{(by (5.3))} \\
&\geq -\frac{\kappa}{2} \|x^{k+1} - \bar{x}\|^2 + \max \left\{ \begin{array}{l} f_1(x^{k+1}) - f_2(x^{k+1}) - \tau_f^\ell + \frac{2\kappa}{2} \|x^{k+1} - \bar{x}\|^2 \\ c_1(x^{k+1}) - c_2(x^{k+1}) - \tau_c^\ell + \frac{2\kappa}{2} \|x^{k+1} - \bar{x}\|^2 \end{array} \right\} && \text{(by (4.11))} \\
&= H(x^{k+1}; \bar{x}) + \frac{\kappa}{2} \|x^{k+1} - \bar{x}\|^2. && \text{(by (3.10))}
\end{aligned}$$

As $\bar{x} = x^\ell$ and $\check{H}^k(x^\ell; x^\ell) = H(x^\ell; x^\ell)$ due to (5.2), we have shown that the descent test (4.5) is satisfied at $x^{k+1} \neq x^\ell$:

$$H(x^{k+1}; x^\ell) \leq H(x^\ell; x^\ell) - \frac{\kappa}{2} \|x^{k+1} - x^\ell\|^2,$$

contradicting thus the assumption that only null steps are performed for $k > \ell$. Hence, $\tilde{x} = \bar{x}$ and the last stability center solves (5.4). This allows us to conclude (thanks to convexity of $\bar{H}(\cdot; \bar{x})$) that $\bar{x} = x^\ell$ solves $\min_{x \in X} \bar{H}(x; x^\ell)$. Lemma 5.2 then concludes the proof. \square

Remark 5.5 (Bundle compression). It is worth mentioning that the index sets \mathcal{B}_f^k and \mathcal{B}_c^k gathering the information bundle can be kept bounded; each one having at most M_{\max} indices, for a chosen integer $M_{\max} \geq 3$. Indeed, it suffices to keep in the bundles the linearizations issued by the stability center x^{ℓ_k} , the new trial point x^{k+1} and the so-called aggregate linearization as in [5, Eq. 4.5]. When transcribed to our setting, the aggregate linearizations for f_1 and c_1 read as

$$\begin{aligned}
\bar{f}_1^{\alpha_f^k}(x) &:= \bar{f}_1^k(x^{k+1}) + \langle p_f^k, x - x^{k+1} \rangle \leq f_1(x) \quad \forall x \in \mathbb{R}^n \\
\bar{c}_1^{\alpha_c^k}(x) &:= \bar{c}_1^k(x^{k+1}) + \langle p_c^k, x - x^{k+1} \rangle \leq c_1(x) \quad \forall x \in \mathbb{R}^n,
\end{aligned}$$

with $p_f^k := \sum_{j \in \mathcal{B}_f^k} \alpha_f^j s_{f_1}^j$, $p_c^k := \sum_{j \in \mathcal{B}_c^k} \alpha_c^j s_{c_1}^j$ and multipliers α_f, α_c as in (4.10). We claim that the following economical rule for managing \mathcal{B}_f^k and \mathcal{B}_c^k (in Step 4 of Algorithm 1) is enough to ensure convergence:

Serious step: set $\mathcal{B}_f^{k+1} = \{k+1\}$ and $\mathcal{B}_c^{k+1} = \{k+1\}$;

Null step: set $\mathcal{B}_f^{k+1} = \{k+1, \ell_k, a_f^k\}$ and $\mathcal{B}_c^{k+1} = \{k+1, \ell_k, a_c^k\}$.

Indeed, Proposition 5.4 is still valid if the algorithm employs the above economical rule for updating the bundles: the key Proposition 4.3 from [5] still applies and thus the displayed equations right after (5.4) hold. As it can be noted in the sequel, no bundle management restriction (besides the requirement that $k+1 \in \mathcal{B}_f^{k+1} \cap \mathcal{B}_c^{k+1}$) is required after a serious steps. \square

We consider now the case of infinitely many serious steps. To this end, we need the following auxiliary result.

Lemma 5.6. *There exist constants $L, M > 0$ such that, for all $k \in \mathbb{N}$, the three following conditions hold for $p^{k+1} \in \partial_1 \check{H}^k(x^{k+1}; x^{\ell_k})$, $s_X^{k+1} \in N_X(x^{k+1})$, and $e^{k+1} = L \|x^{k+1} - x^{\ell_k}\|$:*

$$\|p^{k+1} + s_X^{k+1}\| \leq \mu_{\max} \|x^{k+1} - x^{\ell_k}\| \leq M, \quad (5.5a)$$

$$p^{k+1} + s_X^{k+1} \in \partial_{e^{k+1}} [\bar{H}(x^{\ell_k}; x^{\ell_k}) + i_X(x^{\ell_k})], \quad (5.5b)$$

$$p^{k+1} \in \partial_{e^{k+1}} \bar{H}(x^{\ell_k}; x^{\ell_k}). \quad (5.5c)$$

Proof. As $\mu^k \in (0, \mu_{\max}]$ (c.f. Lemma 5.1), expression (4.9) yields the first inequality in (5.5a). Recall that the iterates x^{k+1} and x^{ℓ_k} are contained in the bounded set X for all k . The second inequality in (5.5a) then

follows. Convexity of the function $\check{H}^k + i_X$ and (4.9) gives that, for all $x \in \mathbb{R}^n$,

$$\begin{aligned} \check{H}^k(x; x^{\ell_k}) + i_X(x) &\geq \check{H}^k(x^{k+1}; x^{\ell_k}) + \langle p^{k+1} + s_X^{k+1}, x - x^{k+1} \rangle \\ &\geq \check{H}^k(x^{k+1}; x^{\ell_k}) + \langle p^{k+1} + s_X^{k+1}, x - x^{\ell_k} \rangle + \langle p^{k+1} + s_X^{k+1}, x^{\ell_k} - x^{k+1} \rangle \\ &\geq \check{H}^k(x^{k+1}; x^{\ell_k}) + \langle p^{k+1} + s_X^{k+1}, x - x^{\ell_k} \rangle - M \|x^{\ell_k} - x^{k+1}\|, \end{aligned} \quad (5.6)$$

where the last inequality is due to (5.5a) and Cauchy-Schwarz inequality. Definition (4.6) of $\check{H}^k(\cdot; x^{\ell_k})$ as well as the fact that $\ell_k \in \mathcal{B}_f^k \cap \mathcal{B}_c^k$ give the following chain of inequalities:

$$\begin{aligned} \check{H}^k(x^{k+1}; x^{\ell_k}) &\geq \max \left\{ \bar{f}_1^{\ell_k}(x^{k+1}) - \bar{f}_2^{\ell_k}(x^{k+1}) - \tau_f(x^{\ell_k}), \bar{c}_1^{\ell_k}(x^{k+1}) - \bar{c}_2^{\ell_k}(x^{k+1}) - \tau_c(x^{\ell_k}) \right\} \\ &= \max \left\{ f_1(x^{\ell_k}) - f_2(x^{\ell_k}) + \langle s_{f_1}^{\ell_k} - s_{f_2}^{\ell_k}, x^{k+1} - x^{\ell_k} \rangle - \tau_f(x^{\ell_k}), \right. \\ &\quad \left. c_1(x^{\ell_k}) - c_2(x^{\ell_k}) + \langle s_{c_1}^{\ell_k} - s_{c_2}^{\ell_k}, x^{k+1} - x^{\ell_k} \rangle - \tau_c(x^{\ell_k}) \right\} \\ &\geq \max \left\{ f_1(x^{\ell_k}) - f_2(x^{\ell_k}) - \tau_f(x^{\ell_k}), c_1(x^{\ell_k}) - c_2(x^{\ell_k}) - \tau_c(x^{\ell_k}) \right\} \\ &\quad + \min \left\{ \langle s_{f_1}^{\ell_k} - s_{f_2}^{\ell_k}, x^{k+1} - x^{\ell_k} \rangle, \langle s_{c_1}^{\ell_k} - s_{c_2}^{\ell_k}, x^{k+1} - x^{\ell_k} \rangle \right\}. \end{aligned}$$

Since $X \subset \mathcal{O}$ is compact, we have that $\partial f_1(X)$, $\partial c_1(X)$, $\partial^c f_2(X)$, and $\partial^c c_2(X)$ are bounded sets (see Section 2). Hence, there exist $K_f > 0$ and $K_c > 0$ such that, $\|s_{f_1}^{\ell_k} - s_{f_2}^{\ell_k}\| \leq K_f$ and $\|s_{c_1}^{\ell_k} - s_{c_2}^{\ell_k}\| \leq K_c$ for all k . Applying the Cauchy-Schwarz inequality to the inequalities above and recalling that

$$\max \left\{ f_1(x^{\ell_k}) - f_2(x^{\ell_k}) - \tau_f(x^{\ell_k}), c_1(x^{\ell_k}) - c_2(x^{\ell_k}) - \tau_c(x^{\ell_k}) \right\} = H(x^{\ell_k}; x^{\ell_k}) = \bar{H}(x^{\ell_k}; x^{\ell_k})$$

by definition, we get

$$\check{H}^k(x^{k+1}; x^{\ell_k}) \geq \bar{H}(x^{\ell_k}; x^{\ell_k}) - L_0 \|x^{k+1} - x^{\ell_k}\|, \quad \text{with } L_0 = \max\{K_f, K_c\}. \quad (5.7)$$

Recall that $\bar{H}(x; x^{\ell_k}) \geq \check{H}^k(x; x^{\ell_k})$ for all $x \in \mathbb{R}^n$ and combine (5.6) with (5.7) to obtain

$$\begin{aligned} \bar{H}(x; x^{\ell_k}) + i_X(x) &\geq \check{H}^k(x; x^{\ell_k}) + i_X(x) \\ &\geq \check{H}^k(x^{k+1}; x^{\ell_k}) + \langle p^{k+1} + s_X^{k+1}, x - x^{\ell_k} \rangle - M \|x^{\ell_k} - x^{k+1}\| \\ &\geq \bar{H}(x^{\ell_k}; x^{\ell_k}) - (L_0 + M) \|x^{k+1} - x^{\ell_k}\| + \langle p^{k+1} + s_X^{k+1}, x - x^{\ell_k} \rangle. \end{aligned}$$

We have thus shown (5.5b) with $L = M + L_0$. To prove the last inclusion (5.5c), observe that this chain of inequalities remains true if the term $i_X(x)$ is excluded together with corresponding subdifferential s_X^{k+1} : for all $x \in \mathbb{R}^n$,

$$\begin{aligned} \bar{H}(x; x^{\ell_k}) &\geq \check{H}^k(x; x^{\ell_k}) \geq \check{H}^k(x^{k+1}; x^{\ell_k}) + \langle p^{k+1}, x - x^{\ell_k} \rangle - M \|x^{\ell_k} - x^{k+1}\| \\ &\geq \bar{H}(x^{\ell_k}; x^{\ell_k}) + \langle p^{k+1}, x - x^{\ell_k} \rangle - (L_0 + M) \|x^{k+1} - x^{\ell_k}\|. \end{aligned}$$

□

Lemma 5.6 can be applied to any iteration $k + 1$ between the serious step ℓ_k and ℓ_{k+1} (ℓ_{k+1} included). Indeed, equation (4.9) used to show (5.5a) holds true. Other arguments used in the proof remain valid for any iteration between ℓ_k and ℓ_{k+1} , as from the point of view of the algorithm, the only change is the bundle information, which is not explicitly used in the proof.

Observe that $\check{H}^k(\cdot; x^{\ell_k})$ given in (4.6) is the pointwise maximum of finitely many affine functions. Hence, its subdifferential is the convex hull of the ‘‘active’’ linearization slopes, i.e., Proposition 2.1 ii) asserts that

$$\partial_1 \check{H}^k(x^{k+1}; x^{\ell_k}) := \text{conv} \left\{ \left\{ s_{f_1}^j - s_{f_2}^{\ell_k} \right\}_{j \in \bar{\mathcal{B}}_f^k}, \left\{ s_{c_1}^j - s_{c_2}^{\ell_k} \right\}_{j \in \bar{\mathcal{B}}_c^k} \right\}, \quad (5.8)$$

with $\bar{\mathcal{B}}_f^k$ and $\bar{\mathcal{B}}_c^k$ given in (4.10). Since $X \subset \mathcal{O}$ is compact, we have that $\partial f_1(X)$, $\partial c_1(X)$, $\partial^c f_2(X)$, and $\partial^c c_2(X)$ are bounded sets (see Section 2). Thus, (4.9) and (5.8) certificate that the sequence of model’s subgradients $\{p^k\}$ is bounded. This property is used in the proof of the following proposition.

Proposition 5.7 (Infinitely many serious steps). *Assume that the algorithm performs infinitely many serious steps. Then, any cluster point $\bar{x} \in X$ of the sequence $\{x^{\ell_k}\}_k$ satisfies the necessary optimality condition (3.13).*

Proof. We first show that

$$\lim_{k \rightarrow \infty} \|x^{\ell_{k+1}} - x^{\ell_k}\| = 0. \quad (5.9)$$

To this end, we must analyze the two cases of Lemma 4.1. In case i), Algorithm 1 produces a feasible point for (1.1) after finitely many serious steps and all subsequent points are feasible. Let $x^{\ell_{k_1}}$ be the first feasible serious iterate. Then, Lemma 4.1 i) yields

$$f(x^{\ell_{k+1}}) \leq f(x^{\ell_k}) - \frac{\kappa}{2} \|x^{\ell_{k+1}} - x^{\ell_k}\|^2 \text{ and } c(x^{\ell_{k+1}}) \leq 0 \text{ for all } k \geq \ell_{k_1}.$$

The telescopic sum of the first inequality above yields

$$\sum_{k=k_1}^{\infty} \|x^{\ell_{k+1}} - x^{\ell_k}\|^2 \leq \frac{2}{\kappa} \sum_{k=k_1}^{\infty} \left(f(x^{\ell_k}) - f(x^{\ell_{k+1}}) \right) \leq \frac{2}{\kappa} \left(f(x^{\ell_{k_1}}) - \lim_{k \rightarrow \infty} f(x^{\ell_{k+1}}) \right).$$

Since f is finite-valued and continuous over the bounded set X , the right-hand side of the above inequality is finite. Hence, (5.9) holds. Assume now that the sequence $\{x^{\ell_k}\}$ is infeasible for (1.1). Lemma 4.1 ii) yields

$$0 < c(x^{\ell_{k+1}}) \leq c(x^{\ell_k}) - \frac{\kappa}{2} \|x^{\ell_{k+1}} - x^{\ell_k}\|^2 \text{ for all } \ell.$$

Once again, by using the telescopic sum we get (5.9).

As $X \subset \mathcal{O}$ is compact, with \mathcal{O} an open set contained in the domains of component functions, and the generalized subdifferential is locally compact, we conclude that for $s_{f_2}^{\ell_k} \in \partial^c f_2(x^{\ell_k})$ and $s_{c_2}^{\ell_k} \in \partial^c c_2(x^{\ell_k})$

$$\{x^{\ell_k}\}, \quad \{s_{f_2}^{\ell_k}\} \quad \text{and} \quad \{s_{c_2}^{\ell_k}\} \quad \text{are bounded sequences.}$$

By taking subsequences, we can define an index set $\mathcal{K} \subset \{0, 1, 2, \dots\}$ such that

$$\lim_{\mathcal{K} \ni k \rightarrow \infty} x^{\ell_k} = \bar{x} \in X, \quad \lim_{\mathcal{K} \ni k \rightarrow \infty} s_{f_2}^{\ell_k} = \bar{s}_{f_2} \in \partial^c f_2(\bar{x}) \quad \text{and} \quad \lim_{\mathcal{K} \ni k \rightarrow \infty} s_{c_2}^{\ell_k} = \bar{s}_{c_2} \in \partial^c c_2(\bar{x}),$$

where the two last limits are due to the fact that the generalized subdifferential is outer-semicontinuous [4, Prop. 2.1.5(b)]. Let us now define $\varphi_k(\cdot) = \bar{\mathbf{H}}(\cdot; x^{\ell_k})$ and $\varphi(\cdot) = \bar{\mathbf{H}}(\cdot; \bar{x})$, with the latter defined in (5.1) with $y = \bar{x}$ and the pair of generalized subgradients $(\bar{s}_{f_2}, \bar{s}_{c_2})$ above. For every $x \in \mathbb{R}^n$ fixed, the above limits imply

$$\lim_{\mathcal{K} \ni k \rightarrow \infty} \bar{\mathbf{H}}(x; x^{\ell_k}) = \bar{\mathbf{H}}(x; \bar{x}),$$

i.e., $\{\varphi_k\}_{k \in \mathcal{K}}$ converges pointwise to φ . Applying Lemma 5.6 for $k+1 = \ell_{k+1}$, we obtain $p^{\ell_{k+1}} \in \partial_{e^{\ell_{k+1}}} \varphi_{\ell}(x^{\ell_k})$ (Eq. (5.5c)). Furthermore, as the sequence $\{p^{\ell_{k+1}}\}$ is bounded (see the paragraph right before this proposition), we may take another subsequence indexed by $\mathcal{K}' \subset \mathcal{K}$ so that $\lim_{\mathcal{K}' \ni \ell \rightarrow \infty} p^{\ell_{k+1}} = \bar{p} \in \mathbb{R}^n$ and $\lim_{\mathcal{K}' \ni \ell \rightarrow \infty} e^{\ell_{k+1}} = 0$ in view of (5.9) and definition of $e^{\ell_{k+1}}$ given in Lemma 5.6. With these conditions at hand, Lemma A.1 (in Appendix A) ensures that $\bar{p} \in \partial \varphi(\bar{x})$, i.e., $\bar{p} \in \partial_1 \bar{\mathbf{H}}(\bar{x}; \bar{x})$. Next, observe that $\{s_X^{\ell_{k+1}}\}$ is a bounded sequence as the inequality

$$\|p^{\ell_{k+1}} + s_X^{\ell_{k+1}}\| \leq \mu_{\max} \|x^{\ell_{k+1}} - x^{\ell_k}\| \quad (5.10)$$

holds due to Lemma 5.6 (with μ_{\max} finite due to Lemma 5.1). By definition of the convex normal cone, it follows that there exists a suitable subsequence of $\{s_X^{\ell_{k+1}}\}_{k \in \mathcal{K}'}$, with $\mathcal{K}'' \subset \mathcal{K}'$ converging to a cluster point $\bar{s} \in N_X(\bar{x}) = \partial i_X(\bar{x})$. Hence, since X is polyhedral and $\text{ri}(\text{Dom}(\bar{\mathbf{H}}(\cdot; \bar{x}))) = \mathcal{O} \neq \emptyset$,

$$\bar{p} + \bar{s} \in \partial_1 \bar{\mathbf{H}}(\bar{x}; \bar{x}) + \partial i_X(\bar{x}) = \partial_1 [\bar{\mathbf{H}}(\bar{x}; \bar{x}) + i_X(\bar{x})]. \quad (5.11)$$

Finally, inequality (5.10) combined with (5.9) yield $\bar{p} + \bar{s} = 0$. Hence, $0 \in \partial_1 [\bar{\mathbf{H}}(\bar{x}; \bar{x}) + i_X(\bar{x})]$, showing that \bar{x} minimizes $\bar{\mathbf{H}}(\cdot; \bar{x})$ over X . Lemma 5.2 thus concludes the proof. \square

The following theorem sums up the algorithm's convergence analysis.

Theorem 5.8 (Convergence analysis). *Let $X \neq \emptyset$ be a bounded polyhedron contained in the open set \mathcal{O} , $f_1, c_1 : \mathcal{O} \rightarrow \mathbb{R}$ convex, and $f_2, c_2 : \mathcal{O} \rightarrow \mathbb{R}$ weakly convex functions on some neighbourhood of each $x \in \mathcal{O}$. If in Algorithm 1 the stopping test tolerance ToI is set to zero, then any cluster point \bar{x} of the sequence of stability centers $\{x^{\ell_k}\}$ satisfies the necessary optimality condition (3.13).*

Furthermore, concerning the original problem (1.1), the following holds:

- i) If $c(\bar{x}) > 0$ (which cannot happen if x^0 is feasible), then \bar{x} is a critical point of $\min_{x \in X} c_1(x) - c_2(x)$.
- ii) If $c(\bar{x}) = 0$ and \bar{x} is not a critical point of $\min_{x \in X} c_1(x) - c_2(x)$, then \bar{x} satisfies the criticality condition (3.9) with $\bar{\lambda} > 0$.

iii) If $c(\bar{x}) < 0$, then \bar{x} satisfies the criticality condition (3.9) with $\bar{\lambda} = 0$.

If $\text{To1} > 0$, then the algorithm stops after finitely many steps $k \in \mathbb{N}$ with an approximate critical point $x^{\ell k}$ of (3.13).

Proof. For the case $\text{To1} = 0$, condition (3.13) follows directly from Proposition 5.3 if $\{x^{\ell k}\}$ is finite, from Proposition 5.4 if the algorithm produces only finitely many serious steps followed by an infinite sequence of null steps, and from Proposition 5.7 if infinitely many serious steps are produced.

Furthermore, the connection with the necessary optimality condition for the original problem (1.1) is established by Theorem 3.5.

Proposition 5.4 ensures that $\lim_{k \rightarrow \infty} \|x^{k+1} - x^\ell\| = 0$ if x^ℓ is the last stability center. Otherwise, $\lim_{k \rightarrow \infty} \|x^{\ell k+1} - x^{\ell k}\| = 0$, as shown in the proof of Proposition 5.7. Thus, Algorithm 1 stops after finitely many steps provided $\text{To1} > 0$. \square

6 Illustrative numerical examples

A deep analysis of the numerical performance of Algorithm 1 (CwC-PBM) is out of the scope of this paper. Instead, this section aims to illustrate our approach to solving some challenging test problems. Here we have two goals: show that it provides good-quality critical points (examples of Subsection 6.2 - Subsection 6.5) and is able to solve problems that, to the best of our knowledge, could not be resolved with other solvers (Subsection 6.1). We consider four nonconvex stochastic optimization problems and one coming from signal processing. Notice that the stochastic problems (example of Subsection 6.1 - Subsection 6.3) do not have explicit DC decompositions, and thus DC programming algorithms are not directly applicable.

Numerical experiments were performed using MATLAB R2020a and Gurobi 9.5.1 (for solving the master QP subproblem (4.8) in Algorithm 1) on a personal computer with the following characteristics: Windows 10 Professional, 32 Go, Intel i7-10850H (6 cores). Our implementation allows for simpler problems without non-linear constraints (as in the case of the problem in Subsection 6.2). We invite the interested reader to check Appendix B for a brief presentation of how the approach can be simplified in this case.

Unless otherwise specified, the choice of the parameters in Algorithm 1 is as follows: $\rho = \sigma = \frac{1}{2}$, $\kappa = 0.3$, $\mu^0 = 10^2$ and $\text{To1} = 10^{-4}$.

6.1 Highly nonconvex chance-constrained problem

In this section we investigate the following optimization problem (having weakly-convex constraint):

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & c^\top x \\ \text{s.t.} \quad & \mathbb{P} \left[\frac{1}{2} \xi^\top Q_j(x) \xi + q_j(x)^\top \xi + d_j(x) \leq 0, j = 1, \dots, k \right] \geq p \\ & \underline{x} \leq x \leq \bar{x}. \end{aligned} \quad (6.1)$$

We first note that as a result of [47] and the upfollowing concrete data, that the probability function is continuously differentiable. Furthermore the underlying feasible set is compact and so we are in case ii) of the introduction: $c_1(x) = p$ and $-c_2$ indicating the probability function. The underlying data is not convex in the parameter replaced by the random vector. As a result, the underlying feasible set is not expected to be convex. Concretely we will consider the following data, for $k = 2$:

$$Q_1(x) = \begin{bmatrix} 3(x_1 - 1) & -x_2 \\ -x_2 & 3(x_1 - 1) \end{bmatrix}, \quad Q_2(x) = \begin{bmatrix} -2x_2 & x_1 - 1 \\ x_1 - 1 & -2x_2 \end{bmatrix}$$

as well as

$$q_1(x) = \begin{bmatrix} 3 \\ 1 \end{bmatrix} (x_1 - 1), \quad q_2(x) = \begin{bmatrix} 1 \\ 4 \end{bmatrix} x_2$$

$d_1(x) = -2, d_2(x) = -2$. We have also picked $p = 0.7$ together with $c = (-1, -1)$, $\underline{x} = (-2, -2)$, $\bar{x} = (2, 2)$. The random vector is taken to be multivariate Gaussian with mean vector 0 and covariance matrix

$$\Sigma = \begin{bmatrix} 2 & -1 \\ -1 & 2 \end{bmatrix}.$$

This optimization problem is quite challenging. First it can be observed that an alternative sample average approximation along the lines of [22] would be a MILP. It is thus tempting to first try to solve the resulting optimization problem with such a formulation. We have done so for the following sample sizes $\{100, 1000, 5000, 10000, 50000\}$, using CPLEX 12.10. The resulting computation times are 0.5, 4, 22, 100, >

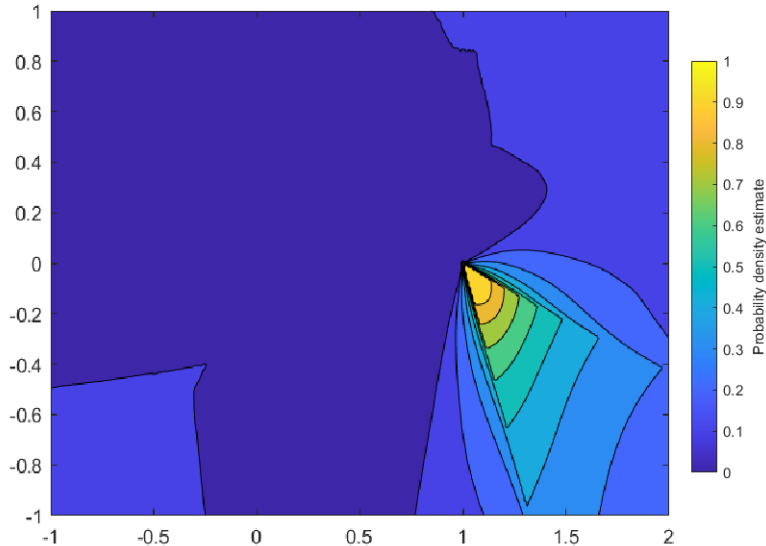


Figure 2: Probability distribution associated with the chance constraint in (6.1)

8000 seconds respectively. The last computation was aborted still showcasing a 25.9% gap after more than 2 hours of computation. Unfortunately, none of the obtained solutions turn out to be feasible, quite to the contrary: the typically obtained final probability value is roughly 0.04 being far from the required 0.7. We have also performed a run of a sampled problem with 10000 random realizations, but with a significantly higher probability level of 0.9. In this case, after roughly one hour of computation, the resulting solution being at a MIPgap of 3.3 %, is still highly infeasible having only a probability value of 0.02.

Furthermore, we have tested **IPOPT solver** for the problem resolution: tests have been performed for the six initial points listed in Table 1 and default tolerance 10^{-4} . After at most 26 seconds of computation, IPOPT halted with highly infeasible points with a probability constraint value equal to 0. The difficulty of generating feasible points might come from the form of probability distribution as the probability level sharply raises from zero (blue, Figure 2) to nearly 1 (yellow, Figure 2) when approaching the feasible area from most directions, which causes the loss of gradient information in a large zone of zero probability. But of course this gradient information is not exploited by the MILP formulation at all.

In contrast, as Table 1 shows, the **CwC-PBM algorithm** manages to improve the probability level if the starting point is infeasible, and to improve the objective function value while satisfying probability constraint for a feasible initial point. Moreover, for one of the tested starting points we have managed to generate the near (globally) optimal solution (1.2400; -0.1126). Since the problem is indeed very difficult, a precise internal sampling scheme for the probability function is required. This amounts to the number of samples used to compute a formula of the type (6.4), which subsequently leads to design of the oracle for c_2 component. We can prematurely end further sampling, very much like the implementation of Genz' code [9], by checking if sampling variance - in fact the confidence interval bounds - are sufficiently small. Unlike Genz' code a crude antithetic Monte-Carlo scheme has been used for sampling, thus leaving much room for significant improvement in terms of speed and precision, by using for instance QMC as, for instance, in [11].

Initial point	Time (s)	Iterations	Initial probability level	Final probability level	Objective value
$(1.2, -0.1)^\top$	47.6	17	0.80	0.70	-1.1097
$(1.5, -0.4)^\top$	78.4	28	0.43	0.70	-1.1346
$(1.6, -0.3)^\top$	36.9	14	0.42	0.70	-1.0815
$(1.7, -0.4)^\top$	59.8	20	0.36	0.70	-1.1443
$(1.1, -0.4)^\top$	21.3	10	0.43	0.43	-0.7001
$(1, -0.8)^\top$	10.8	9	0.18	0.18	-0.2001

Table 1: Results obtained with the CwC-PBM algorithm for problem (6.1) depending on initialization.

6.2 Investment like problems

We will see how the structure of Example 1.1 can appear in practice. Here we will follow the general discussion in [46]. We are interested in the situation wherein we dispose of a set of different technologies $i = 1, \dots, m$ capable of generating electricity. Each technology comes with a specific and detailed set of constraints P_i , cost function c_i attributing to $p_i \in \mathbb{R}^T$ the cost of generation. Altogether, the various technologies are meant to ensure the satisfaction of a given customer load $d \in \mathbb{R}^T$. We are interested in finding the optimal mix. Thus for $i = 1, \dots, m$, we are given $\theta_i \in \mathbb{N}$, the number of “units” of a given type we would like to invest in. The vector θ comes with an investment cost $F(\theta)$. In a deterministic setting this would amount to solving

$$\begin{aligned} \min_{\theta \in \Theta, p_i^j \in P_i} \quad & F(\theta) + \sum_{i=1}^m \sum_{j=1}^{\theta_i} c_i(p_i^j) \\ \text{s.t.} \quad & \sum_{i=1}^m \sum_{j=1}^{\theta_i} p_i^j \geq d. \end{aligned}$$

Now should for each i , the mappings c_i as well as the feasible sets P_i be convex, then it must be so that the averaged solution: $p_i^* = \frac{1}{\theta_i} \sum_{j=1}^{\theta_i} (p_i^j)^*$, in which each power plant of technology i produces exactly this amount is also optimal. This follows from using convexity of P_i showing feasibility of p_i^* and through using Jensen’s inequality for c_i . This is also exactly what would happen if we would solve the subproblems of the Lagrangian dual (w.r.t. the load constraint) for a given fixed investment vector. The convexifying effect of the Lagrangian is well known, e.g., [19, 44] and thus for this dual setting convexity of c_i or P_i would not be essential. Either way, as a result we may thus assume that each power plant of the same technology produces the same amount. This would thus lead to the simpler problem (less variables):

$$\begin{aligned} \min_{\theta \in \Theta, p_i \in P_i} \quad & F(\theta) + \sum_{i=1}^m \theta_i c_i(p_i) \\ \text{s.t.} \quad & \sum_{i=1}^m \theta_i p_i \geq d. \end{aligned}$$

We will investigate a two-stage stochastic version of the last problem, wherein d is uncertain. We thus define:

$$Q(\theta, \xi) = \min_{p_i \in P_i} \sum_{i=1}^m \theta_i c_i(p_i) \text{ s.t. } \sum_{i=1}^m \theta_i p_i \geq d. \quad (6.2)$$

and consider the optimization problem

$$\min_{\theta \in \Theta} F(\theta) + \mathbb{E}[Q(\theta, \xi)], \quad (6.3)$$

where for the sake of simplicity we will assume θ to be allowed to take continuous values (Θ is a polytope). We will also assume that the feasible set P_i is convex, although one could investigate problem (6.3) without this assumption - for instance by arguing through Lagrangian duality.

Let us now look at a concrete instance. We will consider a time horizon of $t = 1, \dots, T$ time steps where each time step is considered to be Δt hours long. The problem disposes of m types of technology, having the following characteristics. Each technology type has a maximum power output level p_i^{mx} , proportional cost c_i and gradient condition g_i . Additionally, each unit is assumed to dispose of a carbon emission rate e_i , and the system subject to a carbon cost f . Concretely this means that the proportional cost gets updated through the formula $c_i \leftarrow c_i + fe_i$.

The system is moreover endowed with a given customer load that we will assume to be multivariate Gaussian with a given mean and positive definite Covariance matrix. We refer to [37, § 5] for the description of P_i (polyhedral). Furthermore for the various technologies we will assume that $F(\theta) = F^\top \theta$. The purpose of our experiment is to showcase how concretely the new algorithm can process specifically structured problems such as these.

Following the description of Example 1.1, we need to compute $Q^\varepsilon(\theta, \xi)$ at each given θ . The latter involves the solution of a convex optimization problem, wherein ψ_t is given by the t -th component of $d - \sum_{i=1}^m \theta_i p_i$. As a result of the logarithm, the objective function defining Q^ε is convex in y . We will therefore use a cutting-plane approach to internally compute Q^ε , as well as it's gradient. The inner optimization is initialized from the optimal solution y_0 of the inner optimization problem of $Q(x, \xi)$. The latter can be computed by solving a linear program. This will give us the oracle for f_2 (in the notation of Subsection 4.1).

Table 2 provides the concrete data.

	1	2	3	4	5	6	7	8	9	10
p^{mx} (MW)	900	900	900	300	300	200	200	200	100	10000
g (MW/h)	100	100	100	30	30	20	20	70	70	5000
c (€/MWh)	30	35	37	45	47	60	100	110	150	10000
F^{inv} (€)	493151	493151	493151	41096	41096	32877	32877	32877	21918	0
e (t/MW)	0	0	0	1	1	0.5	0.5	0.5	1.1	0

Table 2: Data for the stylized investment problem

We can observe that the last unit described in the previous table is an imbalance unit - a computational trick to ensure that one can always meet the load, in this case even despite a potentially completely unbalanced set of invested assets. In terms of constraints on investment, we do not allow investment in this last unit, the capacity will remain at 1. The cost of investment was set up using typical values of investment cost per kW, upon rescaling to match T and while accounting for life span of the various technologies. The data of the case is stylized and the general purpose of the study is more a demonstration of the capabilities of the algorithm rather than an attempt to provide practical insights into investment regarding the electrical system.

We have performed tests with the **CwC-PBM algorithm** and **IPOPT solver**, both applied to approximation with Q^ε , $\varepsilon = 10^{-2}$, for $f = 0$ (zero carbon cost) and $f = 100$. The latter case enables to see the potential impact of such a penalized setting for emitting technologies. The methods provide comparable results in terms of objective value with the average relative difference of 0.98% (both for default tolerance 10^{-4}), see Table 3. (Both algorithms exploit the problem's decomposable structure and employ the same oracles.) The average execution time (among 6 considered cases) is 2579 seconds for the CwC-PBM algorithm, and constitutes to 2241 seconds for IPOPT (including function evaluation). For initial states $v_2 = (1111101011)^\top$, the resulting total amount of installed capacity is less for the case with $f = 100$ (around 3500 MW for both methods) compared to the case with zero carbon cost (more than 3900 MW for both methods). While the latter remains close to the nominal investment vector, the methods update the set of installed capacities quite significantly, by shifting essentially all generation from carbon emitting technologies to technologies 1 – 3 not emitting CO2 at all for the case with $f = 100$. This can be explained since the original setting was slightly overcapacitated - and as a result of the introduction of the "fictive" imbalance unit - infeasibility is no longer an issue. The maximum load over the considered scenarios being roughly 3500 as well. A similar situation occurs for the initial state $v_3 = (1111111111)^\top$ (roughly, 4000 MW for $f = 100$ against 3700 MW for zero carbon cost). For $v_1 = (1111100001)^\top$, the solution found by IPOPT corresponds to 3700 MW of installed capacity for $f = 100$, compared to approximately 3500 MW suggested by the CwC-PBM

algorithm. This numerical experiment thus clearly shows that the CwC-PBM algorithm provides meaningful good quality solutions for this type of problems.

Initial state	CO2 cost	CwC-PBM algorithm		IPOPT solver	
		Iterations	Objective value	Iterations	Objective value
$v_1 = (11111100001)^\top$	0	49	5 781 983	16	5 864 083
	100	33	5 886 797	10	6 001 847
$v_2 = (11111101011)^\top$	0	29	5 965 697	9	5 865 690
	100	31	5 938 787	15	5 920 723
$v_3 = (11111111111)^\top$	0	16	5 985 021	6	5 970 187
	100	32	5 945 049	11	5 930 397

Table 3: Results obtained with CwC-PBM algorithm and IPOPT solver for (6.3) for carbon cost $f = 0$ and 100. Average execution time is 2 579 seconds for the CwC-PBM algorithm and 2 241 seconds for IPOPT.

Setting the prox-parameter μ^0 equal to its value at the first serious step μ^{l^1} predictably speeds up the CwC-PBM algorithm’s performance. For initialization at v_2 and v_3 , the prox-parameter μ^{l^1} is of the order of 10^6 , while it is equal to its default value of 10^2 for initialization at v_1 . The number of iterations decreases significantly for the former cases—down to 15 and 24 in the case of v_2 , and 9 and 28 in the case of v_3 —with an average improvement of 27% in execution time. When altering the stopping test tolerance To1 within the set $\{10^{-4}, 10^{-6}, 10^{-8}, 10^{-12}\}$ (with fixed $\mu^0 = \mu^{l^1}$), we observe that the objective value does not change for all considered initial states and carbon costs. To progress to the next tolerance level within the specified set, the algorithm undergoes between 1 and 4 iterations, resulting in duration ranging from 44 to 456 seconds depending on the case. Moreover, the objective value shows minimal variation with alterations in the parameter κ . Varying κ within the set $\{0.001, 0.01, 0.03, 0.1, 0.3\}$, the objective value remains within the limits of $\pm 0.01\%$ (tests have been performed with the initial state v_2). The algorithm is thus robust with respect to the choice of To1 and κ and can be accelerated with the selection of an appropriate prox-parameter μ .

6.3 Decision dependent probability constraints in two stage problems

In this section we consider the following stochastic problem having different random vectors:

$$\begin{cases} \min_{x \in X} & f_1(x) + \sum_{s=1}^S \pi_s Q(x; \xi^s) \\ \text{s.t.} & \mathbb{P}[A_1 x + b_1 \geq \omega_1] \geq p_1 \end{cases} \quad \text{with} \quad Q(x; \xi) := \begin{cases} \min_{y \in Y} & q(x, y; \xi) \\ \text{s.t.} & \mathbb{P}_x[A_2(\xi)y + b_2(\xi) \geq \omega_2(x)] \geq p_2. \end{cases}$$

In this problem, $\xi \in \Xi := \{\xi^1, \dots, \xi^S\}$, $\omega_1 \sim \mathcal{N}(\mu_1, \Sigma_1)$, and $\omega_2(x) \sim \mathcal{N}(\mu_2(x), \Sigma_2(x))$. The latter random vector depends on the first-stage decisions. We assume that the covariance matrices Σ_1 and $\Sigma_2(x)$ are positive definite for all $x \in X$. As a result, the probability functions are twice-differentiable [13, 45]. Furthermore, as the multivariate Gaussian distribution is log-concave, we get that $c_1(x) = \log(p_1) - \log(\mathbb{P}[A_1 x + b_1 \geq \omega_1])$ is a convex function and so is the objective of the penalized subproblem

$$Q^\varepsilon(x; \xi) = \min_{y \in Y} q(x, y; \xi) - \frac{1}{\varepsilon} \log \left(\mathbb{P}_x[A_2(\xi)y + b_2(\xi) \geq \omega_2(x)] - p_2 \right).$$

We are thus in the setting of Example 1.1 with $f_2(x) = \sum_{s=1}^S \pi_s [-Q^\varepsilon(x; \xi^s)]$. We can observe that the just given optimization problem is a version of two-stage stochastic program having unhedgeable, or post-decision random realizations.

Now in order to compute the gradient of both of the involved probability functions, we can rely on two different formulæ for the gradients. The mapping c_1 is continuously differentiable and its gradient can be evaluated by employing the results shown in [49]. The second stage probability function is also differentiable

under a mild regularity condition, its gradient can be evaluated using the formulæ from [43, Thm. 5.1]. Indeed with $L_2(x)$ the matrix resulting from the Cholesky decomposition $\Sigma_2(x) = L_2(x)L_2(x)^\top$, we may write

$$c_2^i(x, y) = \mathbb{P}[-A_2(\xi^i)y - b_2(\xi^i) + \mu_2(x) + L_2(x)\omega_2 \leq 0],$$

where $\omega_2 \in \mathbb{R}^m$, $\omega_2 \sim \mathcal{N}(0, I)$. Hence we can observe that:

$$\nabla c_2^i(x, y) = \int_{v \in \mathbb{S}^{m-1}: J^*(v) \neq \emptyset, |J^{**}(v)|=1} -\frac{\chi(\hat{\rho}(v))}{(L_2(x))_{j(v)}v} \left(\hat{\rho}(v) \nabla (L_2(x))_{j(v), \cdot} v + \nabla \mu_2(x), -(A_2)_{j(v), \cdot}^\top \right) d\mu_\zeta(v) \quad (6.4)$$

with

$$\begin{aligned} J^*(v) &= \{j = 1, \dots, r : (L_2(x))_j v > 0\} \\ \hat{\rho}(v) &= \min_{j \in J^*(v)} \frac{A_2 y + b_2 - \mu_2(x)}{(L_2(x))_j v} \\ J^{**}(v) &= \left\{ j \in J^*(v) : \hat{\rho}(v) = \frac{A_2 y + b_2 - \mu_2(x)}{(L_2(x))_j v} \right\}. \end{aligned}$$

and $j(v)$ being the unique element of $J^{**}(v)$. In this case since L_2 has linearly independent rows - which is the case since Σ_2 is positive definite - the aforementioned regularity condition (R2CQ) holds true. In fact (R2CQ) holds true locally and as a consequence it is indeed so that both c_1 and c_2 are twice continuously differentiable. This was already clear for c_1 upon using well known classic arguments.

Let us now consider the following concrete example of a problem of this kind. We are interested in a situation considering a manufacturer capable of producing two different products. The first-stage decision variables of the problem consist of setting prices for the products and an advertisement levels. The price will be assumed to be in relation to the average second-stage demand for the given product. We will use the following rule $\mu_2(x) = (\bar{\mu}_1/x_1, \bar{\mu}_2/x_2)$, with x_1, x_2 being the price levels for product 1 and 2 respectively. Advertisement is assumed to have a beneficial effect on the variance of the demand, but simultaneous advertisement for both products will be counterproductive. In other words:

$$\Sigma_2(x) := \begin{bmatrix} (0.1\bar{\mu}_1/x_1 x_3)^2 & -0.4(0.01\bar{\mu}_1/x_1 x_3 \bar{\mu}_2/x_2 x_4) \\ -0.4(0.01\bar{\mu}_1/x_1 x_3 \bar{\mu}_2/x_2 x_4) & (0.1\bar{\mu}_2/x_2 x_4)^2 \end{bmatrix}.$$

The second stage decision y involves the production of the goods. The production process of the goods is subject to some possible unreliability as the amount of actually produced goods are concerned. The matrix A_2 is thus a diagonal matrix, where the first entry is a uniform random variable over the interval $[0.9, 1]$ - on average only 95% of the commissioned products actually get manufactured. The second diagonal entry is uniform over the interval $[0.8, 1]$ - the process of production here is more unreliable. However producing with the more unreliable process is slightly cheaper. Any products that are manufactured but not sold, will incur a penalty. The second stage cost function is thus given by

$$q(x, y, \xi) = (2 - x_1)y_1 + (1 - x_2)y_2 + 12\mathbb{E}[\max\{y_1 - (\omega_2)_1, 0\}] + 12\mathbb{E}[\max\{y_2 - (\omega_2)_2, 0\}].$$

The last two terms correspond to the penalization of produced, but not sold goods. It turns out that the latter expectations can be computed ‘‘analytically’’ as they are related to the computation of an expectation of a truncated Gaussian random variable. Therefore, we can observe that the following identity holds true:

$$\begin{aligned} \mathbb{E}[\max\{y_1 - (\omega_2)_1, 0\}] &= \Phi((y_1 - (\mu_2(x))_1)x_1 / (0.1\bar{\mu}_1 x_3))(y_1 - (\mu_2(x))_1) \\ &\quad - \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}((y_1 - (\mu_2(x))_1)x_1 / (0.1\bar{\mu}_1 x_3))^2} 0.1\bar{\mu}_1 / x_1 x_3. \end{aligned}$$

The second formula is of course immediately deduced as it is analogous. Both products require a different setup of the factory, so $y_1 + y_2 \leq 10$. Furthermore, the first-stage cost function is related to the cost of advertisement $f_1(x) = q_1 x_3^2 + q_2 x_4^2$. Furthermore all first-stage variables are bounded.

The implementation of this example requires first the implementation of the formulæ for the gradient of the probability function. Here we can exploit the earlier given formula immediately. It can be observed (see the more extensive discussion in [38]) that the probability value itself can be computed with exactly the same cost. Subsequently the algorithm scheme is very similar to the one of the investment problem. In particular, combining the computations for probability function value and subgradient (6.4) with the reasoning of the previous example, we will obtain the oracle for f_2 component (in the notation of Subsection 4.1), while the oracle for f_1 is straightforward from the formula of the advertisement cost.

Therefore, we have also run this case with the **CwC-PBM algorithm** and found an approximate critical solution after a total of 12 iterations (1600 seconds on personal laptop). The found solution is $x = (3.37, 3.21, 0.096, 0.784)$, showing that there is interest in balancing the prices, i.e., not taking maximal

prices, while also investing in advertisement. We have done the same test with **IPOPT solver**. The computation was aborted after 50 000 seconds with the resulting infeasible point $x = (9.997, 9.983, 0.0004, 0.167)$ slowly approaching the bound $(10, 10, 0, 0)$.

This example thus shows that the new algorithm allows us to consider settings beyond classic convexity, even when dealing with probability functions - in this case with decision dependent random vectors.

6.4 Compressed sensing problem

In this section we focus on the problem of compressed sensing considered in [52]:

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & \|x\|_1 - \|x\| \\ \text{s.t.} \quad & \|Ax - b\|_{LL_2, \gamma}^2 \leq \delta, \end{aligned} \tag{6.6}$$

where $A \in \mathbb{R}^{q \times n}$ is a full row rank matrix and $b \in \mathbb{R}^q$. For given $\gamma > 0$, Lorentzian norm $\|\cdot\|_{LL_2, \gamma}$ of a vector $y \in \mathbb{R}^q$ is defined as

$$\|y\|_{LL_2, \gamma}^2 = \sum_{i=1}^q \log \left(1 + \frac{y_i^2}{\gamma^2} \right).$$

As discussed in [52], the problem (6.6) is DC with twice continuously differentiable constraint, whose modulus of Lipschitz gradient is known. This allows us to construct the oracle for the constraint component. To compute subgradient of the component $f_1(x) = \|x\|_1$, we have chosen the sign function.

As in [52], we have generated $A \in \mathbb{R}^{q \times n}$ with normally distributed random entries normalizing it so that each column has a unit norm. To set the original point, we have chosen a subset of size $s_0 = \lfloor \frac{q}{9} \rfloor$ among basis vectors and generated a s_0 -sparse vector x_{orig} with i.i.d. normally distributed random entries. We have taken $b = Ax_{orig} + 0.01\eta$, each η_i having a standard Cauchy distribution, and $\delta = 1.1\|0.01\eta\|_{LL_2, \gamma}$ with $\gamma = 0.02$.

We have performed tests with the **CwC-PBM algorithm** and **SCP_{ls} algorithm** coded based on Algorithm 2.1 of the paper [52] for $q = 720i$, $n = 2560i$, with $i = 1$ (Figure 3) and $i = 5$ (Figure 4). Both algorithms were initialized at $x_0 = A^+b$ with matrix A^+ denoting the Moore-Penrose pseudoinverse of A . The gap between solution provided by the SCP_{ls} algorithm (with tolerance 10^{-4}) and CwC-PBM algorithm is 4.45% and 2.71% for 2400 and 3000 iterations, respectively, for the case $i = 1$. It constitutes 2.2% and 1.16% for 2400 and 3000 iterations, respectively, for the case $i = 5$, Table 4. However, the execution time of the CwC-PBM algorithm is higher compared to SCP_{ls}, which was designed for a more specific framework (constraint functions have Lipschitz continuous gradient, and the objective function is decomposed as a sum of a smooth function and DC function), Table 4.

We have also run the CwC-PBM algorithm with initialization at zero vector. It manages to obtain a feasible solution after 50 iterations (for both $i = 1$ and 5), as well as to recover an optimal solution within tolerance 3.2×10^{-4} for $i = 1$ (10 000 iterations) and 1.2×10^{-3} for $i = 5$ (3 000 iterations) with execution time of 1347 seconds and 2717 seconds, respectively. The SCP_{ls} algorithm can not be applied in this case, as a feasible initial point is required.

	CwC-PBM algorithm			SCP _{ls} algorithm	
	Iterations	Time (s)	Objective value	Time (s)	Objective value
i=1	2 400	324.27	54.16	16.02	51.85
	3 000	592.92	53.26		
i=5	2 400	2 396.65	289.52	133.37	283.27
	3 000	3 180.52	286.56		

Table 4: Results obtained with CwC-PBM and SCP_{ls} algorithms for (6.6) with the initial point $x_0 = A^+b$

6.5 Chance-Constrained Optimal Power Flow

This section is dedicated to the operational planning problem of distribution energy grid under uncertainties related to the probabilistic nature of nodal generation and consumption considered in [15, 36]. It is formulated as a chance-constrained Optimal Power Flow (OPF) where the objective function constitutes to the operational planning cost, deterministic constraints are convex and reflect contractual engagements related

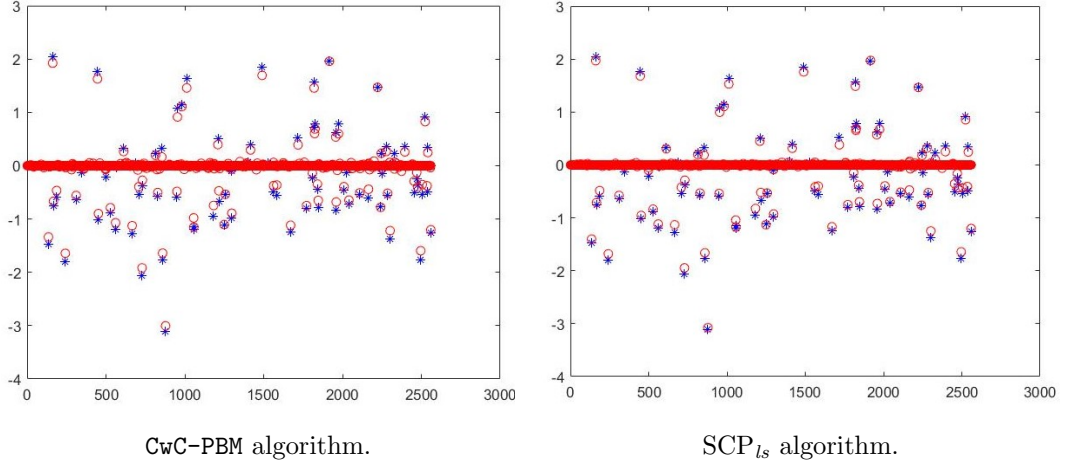


Figure 3: Computed solutions (marked by circle) of (6.6) and x_{orig} (marked by asterisk) for $i = 1$

to the grid operation, and the joint probability constraint incorporates the stochastic nature of the model. More precisely, it ensures that for an operational decision, grid operating conditions remain within technical limits with a given probability $1 - \alpha$ (security level). An operational decision bears on power modulation of a grid user (limitation of power injected to the grid for producers and power curtailment for consumers). Power modulation conditions can be individually contractualized with producers (generators), which may reduce the operational cost due to lower price of power modulation within the specified limits.

Following notations in [36], we denote the set of producers with such individual contracts by \mathcal{N}_{SCP} , and the corresponding decisions on power modulation within contractualized limits by $(\mathbf{p}_i^\gamma, \mathbf{q}_i^\gamma)$, $i \in \mathcal{N}_{SCP}$. Decisions on power modulation above those limits as well as decisions concerning other users are denoted by $(\mathbf{p}_i^\beta, \mathbf{q}_i^\beta)$ with corresponding set of indexes $i \in \mathcal{N}_{FIT}$. To simplify the notation, if there is no need to distinguish user types, we will denote decision variables by

$$\begin{aligned} \mathbf{p} &:= (\{\mathbf{p}_i^\gamma\}_{i \in \mathcal{N}_{SCP}}, \{\mathbf{p}_i^\beta\}_{i \in \mathcal{N}_{FIT}}) \in \mathbb{R}^n \\ \mathbf{q} &:= (\{\mathbf{q}_i^\gamma\}_{i \in \mathcal{N}_{SCP}}, \{\mathbf{q}_i^\beta\}_{i \in \mathcal{N}_{FIT}}) \in \mathbb{R}^n \end{aligned}$$

with $n = |\mathcal{N}_{SCP}| + |\mathcal{N}_{FIT}|$.

In [36] the chance-constrained optimal power flow is reformulated as a DC-constrained DC-problem (with convex objective function, i.e. $f_2 = 0$)

$$\begin{aligned} \min_{(\mathbf{p}, \mathbf{q}) \in X} \quad & \sum_{i \in \mathcal{N}_{SCP}} C^i |\mathbf{p}_i^\gamma| + \sum_{i \in \mathcal{N}_{FIT}} \tilde{C}^i |\mathbf{p}_i^\beta| + \sum_{i \in \mathcal{N}_{SCP}} c^i (\mathbf{p}_i^\gamma)^2 + \sum_{i \in \mathcal{N}_{FIT}} \tilde{c}^i (\mathbf{p}_i^\beta)^2 \\ \text{s.t.} \quad & c_1(\mathbf{p}, \mathbf{q}) - c_2(\mathbf{p}, \mathbf{q}) \leq 0. \end{aligned} \quad (6.7)$$

Here all the coefficients $C^i, c^i, i \in \mathcal{N}_{SCP}$, as well as $\tilde{C}^i, \tilde{c}^i, i \in \mathcal{N}_{FIT}$ are positive, c_1, c_2 are convex functions. This reformulation is made based on an oracle enabling to find a DC decomposition of the constraint under probability sign ([36], Section 3.1), which imposes the OPF solution to be in the required bounds. Applying the sample average approximation, it leads to a DC approximation of the probability constraint ([36], Section 3.2). This procedure thus provides oracles for the DC constraint in problem (6.7). The proximal bundle algorithm from [50] (PBMDC²) is applied to solved it.

As in [36], we consider the case of 33-bus network accommodating 31 loads with 3 generation units among them. One producer has individual contract with corresponding costs $C^1 = 4.2 \cdot 10^{-5}$, and $c^1 = 0$ for power modulation up to 50% of his generation. Above this volume, the costs are $\tilde{C}^1 = 4.2 \cdot 10^{-3}$ and $\tilde{c}^1 = 0.01$. For two other producers, the costs are $\tilde{C}^2 = \tilde{C}^1$, $\tilde{C}^3 = 0.02$ and $\tilde{c}^2 = \tilde{c}^1$, $\tilde{c}^3 = 0.1$, while for the remaining 28 consumers $\tilde{C}^i = \tilde{c}^i = 1$. We use data from Enedis Open Data² on July 27, 2020, to construct load and generation profiles for $N = 1000$ scenarios, only 545 among which satisfy the technical limits (i.e. the security level is 0.545 without power modulation).

²<https://data.enedis.fr/explore/dataset/coefficients-des-profils/table>

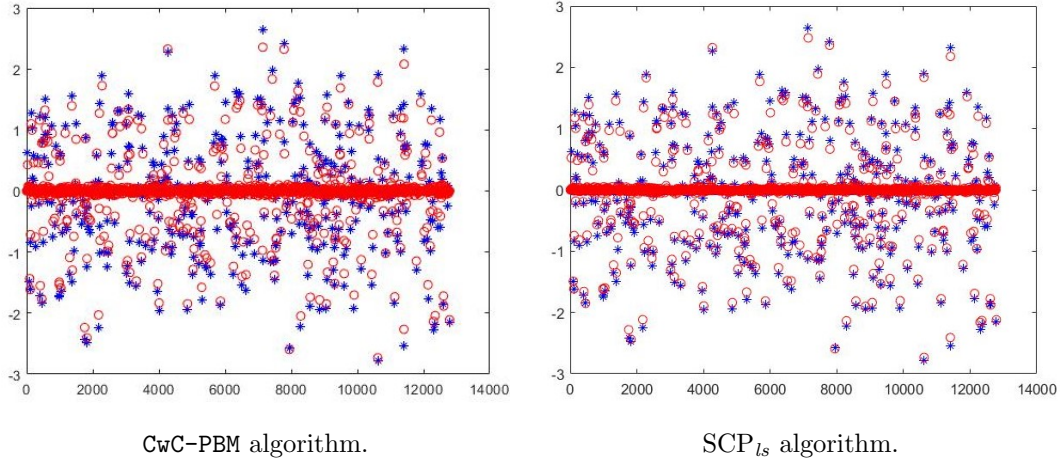


Figure 4: Computed solutions (marked by circle) of (6.6) and x_{orig} (marked by asterisk) for $i = 5$

We have run this problem with the **CwC-PBM algorithm** and **PBMDC² algorithm**. The parameters ρ , σ , κ and **To1** for both methods are set to: $\rho = 10^7$, $\sigma = 0.5$, $\kappa = 0.3$ and **To1** = 10^{-5} . The parameter μ^0 is chosen to be 80 for the CwC-PBM algorithm, while $\mu^0 = 10^2$, $\mu_{\min} = 10^{-6}$ and $\mu_{\max} = 10^6$ for PBMDC². We set 11 values of the security level $1 - \alpha$ ranging from 0.75 to 1 with a step size of 0.025. The initial vector $(\mathbf{p}^0, \mathbf{q}^0)$ is equal to zero, which corresponds to the case with no power modulation.

For all values of security level $1 - \alpha$, the CwC-PBM algorithm yields an approximate critical point of (6.7). It is not necessary the case for PBMDC², as c_2 needs to be differentiable [50][Thm. 2], but this condition does not hold for (6.7) (see Section 3.1 of [36]).

The CwC-PBM algorithm provides better solutions in terms of the objective value for all values of security level $1 - \alpha$, Figure 5 (a). The maximal improvement is of 13.33% is obtained for $1 - \alpha = 0.775$, and constitutes 3.56% in average. As explained in [36], the obtained security level is less than $1 - \alpha$ (Figure 5 (b)) due to DC approximation of probability constraint. The relative maximal difference between the targeted and obtained security level is 5.07% for the CwC-PBM algorithm compared to 4.93% for the PBMDC² algorithm (both at $1 - \alpha = 0.75$). Nevertheless, solutions supplied by both algorithms are feasible for (6.7). The average execution time is less for the PBMDC² algorithm: 1340 seconds compared to 4232 for CwC-PBM. However, the objective value corresponding to PBMDC²'s solution is attained after 2387 seconds in average. The latter difference can be explained by the fact that we have not exploited the DC structure in CwC-PBM.

This use case thus shows that the new algorithm is applicable to real-life industrial problems and capable to generate approximate critical points in DC-constrained framework, without assumptions on differentiability of DC components. However, achieving criticality comes at the cost of an increase in execution time.

7 Conclusion

In this manuscript, we have considered nonsmooth and nonconvex optimization problems where the objective function and nonlinear constraint are represented as the difference of convex and weakly convex functions (CwC). Our work studies various stationary conditions and a bundle method approach enabling to compute critical (generalized KKT) points. The latter broadens and enhances the algorithm developed in [50] for the case of Difference-of-Convex (DC)-constrained DC-problems, and likewise relies on problem reformulation via an improvement function. To the best of our knowledge, proposed method is the first one that directly exploits the CwC-structure of the involved functions and does not require additional assumptions or transformations as, for instance, explicit Difference-of-Convex decompositions or Moreau envelopes. We have illustrated the method performance with a few stochastic problems, including two-stage and chance-constrained problems and a compressed sensing problem with nonlinear constraint. Preliminary results are meaningful and show that the algorithm enables tackling settings beyond the classic Difference-of-Convex setting.

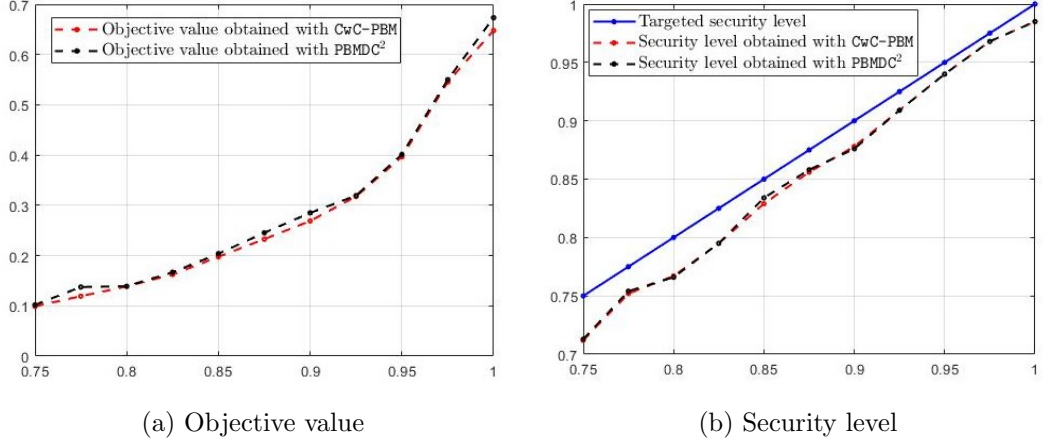


Figure 5: Objective value and security level obtained with CwC-PBM and PBMDC² algorithms.

Appendices

A Some mathematical results

Proof of Proposition 2.4. Since $f : \mathcal{O} \rightarrow \mathbb{R}^n$ is weakly convex, it follows by definition that, relative to some neighbourhood $V_{x'}$ of each point $x' \in \mathcal{O}$, there exist $\mu_{x'} > 0$ such that for all $\mu \geq \mu_{x'}$ the function $\phi(x) = f(x) + \frac{\mu}{2}\|x\|^2$ is finite and convex on $V_{x'}$. In such a representation, there is no loss of generality in assuming that $V_{x'} \subset \mathcal{O}$ (if necessary we can define a new/smaller neighbourhood as $V_{x'} \cap \mathcal{O}$ for which the above conclusion obviously stands). By considering all the points in X , let $V := \{V_{x'} : x' \in X\}$ be the collection of all such neighbourhoods. Then, by construction, V is an open cover of the compact set X and, by definition of compactness, it has a finite open subcover, i.e., there exists finitely many points $\{x'_1, \dots, x'_m\} \subset X$ such that $\mathcal{O}' := \cup_{i=1}^m V_{x'_i} \supset X$, and by construction \mathcal{O}' is an open subset of \mathcal{O} . The first part of item i) thus follows by taking $\mu_f := \max_{i=1, \dots, m} \mu_{x'_i} < \infty$. By writing $f(x) = \phi(x) - \frac{\mu}{2}\|x\|^2$ and recalling Proposition 2.1 i) we get $\partial^c f(x) = \partial \phi(x) - \mu x$ for all $x \in \mathcal{O}'$. This concludes item i).

To show item ii), let us now define $\tilde{\phi}(x) = f(x) + \frac{\mu}{2}\|x\|^2 + i_{\mathcal{O}'}(x)$, an extended real-valued convex function: $\tilde{\phi} : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$. Note that for each $x \in \mathcal{O}'$, there exists a neighbourhood $V_x \subset \mathcal{O}'$ such that $\tilde{\phi}(x') = \phi(x')$ for all $x' \in V_x$. This fact permits us to conclude that $\partial \tilde{\phi}(x) = \partial \phi(x)$ for all $x \in \mathcal{O}'$. It thus follows from item i) that, for every $x \in \mathcal{O}'$ and every $s \in \partial \tilde{\phi}(x)$, there exists $s_f \in \partial^c f(x)$ such that $s = s_f + \mu x$ and the subgradient inequality reads as

$$\tilde{\phi}(y) \geq \tilde{\phi}(x) + \langle s_f + \mu x, y - x \rangle \quad \forall y \in \mathbb{R}^n,$$

i.e., $f(y) + \frac{\mu}{2}\|y\|^2 + i_{\mathcal{O}'}(y) \geq f(x) + \frac{\mu}{2}\|x\|^2 + i_{\mathcal{O}'}(x) + \langle s_f + \mu x, y - x \rangle$ for all $y \in \mathbb{R}^n$. The latter simplifies to

$$f(y) + i_{\mathcal{O}'}(y) \geq f(x) + \langle s_f, y - x \rangle - \frac{\mu}{2}\|y - x\|^2 \quad \forall y \in \mathbb{R}^n.$$

By restricting y to the set X and recalling that $s_f = s - \mu x \in \partial^c f(x)$ is an arbitrary subgradient (because no restriction was imposed to $s \in \partial \tilde{\phi}(x)$), the above inequality becomes (2.4). \square

Lemma A.1. Let $\varphi : \mathbb{R}^n \rightarrow \mathbb{R}$ be a convex function, and $\{\varphi_\ell\}_{\ell \in \mathbb{N}}$ a sequence of convex functions $\varphi_\ell : \mathbb{R}^n \rightarrow \mathbb{R}$ converging pointwise to φ , i.e., $\lim_{\ell \rightarrow \infty} \varphi_\ell(x) = \varphi(x)$ for every given point x . Furthermore, let $\{x^\ell\} \subset \mathbb{R}^n$ be such that $\lim_{\ell \rightarrow \infty} x^\ell = \bar{x}$ and $\{\epsilon^\ell\} \subset \mathbb{R}_+$ satisfy $\lim_{\ell \rightarrow \infty} \epsilon^\ell = 0$. If $g^\ell \in \partial_{\epsilon^\ell} \varphi_\ell(x^\ell)$ for all ℓ and $\lim_{\ell \rightarrow \infty} g^\ell = \bar{g}$, then $\bar{g} \in \partial \varphi(\bar{x})$.

Proof. First, let us prove that $\liminf_{\ell} \varphi_\ell(x^\ell) \geq \varphi(\bar{x})$. Since $\text{dom}(\varphi_\ell) = \text{dom}(\varphi) = \mathbb{R}^n$, it follows from [32, Cor. 2C] that the pointwise convergence of $\{\varphi_\ell\}_{\ell \in \mathbb{N}}$ is equivalent to epi-convergence, which in turn is

equivalent (see [32, Eq. (4.2)]) to epi-convergence of $\{\varphi_\ell^*\}_{\ell \in \mathbb{N}}$, the sequence of conjugate functions to φ_ℓ . Hence, it follows that $\lim_\ell \varphi_\ell^*(x) = \varphi^*(x)$ for every given $x \in \mathbb{R}^n$. Now consider the following development:

$$\varphi_\ell(x^\ell) = (\varphi_\ell^*)^*(x^\ell) = \sup_{y \in \mathbb{R}^n} [\langle y, x^\ell \rangle - \varphi_\ell^*(y)] \geq \langle y, x^\ell \rangle - \varphi_\ell^*(y) \quad \forall y \in \mathbb{R}^n.$$

Accordingly, $\liminf_\ell \varphi_\ell(x^\ell) \geq \liminf_\ell [\langle y, x^\ell \rangle - \varphi_\ell^*(y)] = \langle y, \bar{x} \rangle - \varphi^*(y)$ for all $y \in \mathbb{R}^n$, showing that

$$\liminf_\ell \varphi_\ell(x^\ell) \geq \sup_y [\langle y, \bar{x} \rangle - \varphi^*(y)] = (\varphi^*)^*(\bar{x}) = \varphi(\bar{x}).$$

Recall that $g^\ell \in \partial_{\epsilon^\ell} \varphi_\ell(x^\ell)$. Then, $\varphi_\ell(x) \geq \varphi_\ell(x^\ell) + \langle g^\ell, x - x^\ell \rangle - \epsilon^\ell$ for all $x \in \mathbb{R}^n$. By taking the limit when ℓ goes to infinity we get

$$\begin{aligned} \varphi(x) &= \liminf_\ell \varphi_\ell(x) = \liminf_\ell \varphi_\ell(x) \geq \liminf_\ell [\varphi_\ell(x^\ell) + \langle g^\ell, x - x^\ell \rangle - \epsilon^\ell] \\ &\geq \liminf_\ell \varphi_\ell(x^\ell) + \liminf_\ell \langle g^\ell, x - x^\ell \rangle - \limsup_\ell \epsilon^\ell \\ &\geq \varphi(\bar{x}) + \langle \bar{g}, x - \bar{x} \rangle, \end{aligned}$$

showing that $\bar{g} \in \partial\varphi(\bar{x})$. □

B Simplified algorithm for the case without nonlinear constraints

This section describes how Algorithm 1 can be simplified to deal with the simpler convexly-constrained problem

$$\min_{x \in X} f(x), \quad \text{with } f(x) = f_1(x) - f_2(x). \quad (\text{B.1})$$

In this case, the problem's model (4.6) reduces to $\check{H}^k(x; x^{\ell k}) = \check{f}_1^k(x) - \check{f}_2^{\ell k}(x)$, and the descent test (4.5) becomes $f(x^{k+1}) \leq f(x^{\ell k}) - \frac{\kappa}{2} \|x - x^{\ell k}\|^2$. Hence, Algorithm 1 boils down to the following plainer scheme.

Convergence analysis for Algorithm 2 follows from that of Algorithm 1 upon several simplifications. Instead of doing this exercise, we simply state the following result.

Theorem B.1. *Consider problem (B.1) with $X \neq \emptyset$ a bounded polyhedron contained in the open set \mathcal{O} , $f_1 : \mathcal{O} \rightarrow \mathbb{R}$ convex, and $f_2 : \mathcal{O} \rightarrow \mathbb{R}$ weakly convex on some neighbourhood of each $x \in \mathcal{O}$. If in Algorithm 2 the stopping test tolerance ToI is set to zero, then any cluster point \bar{x} of the sequence of stability centers $\{x^{\ell k}\}$ satisfies the necessary optimality condition (3.3).*

If $\text{ToI} > 0$, then the algorithm stops after finitely many steps $k \in \mathbb{N}$ with an approximate critical point $x^{\ell k}$ of (3.3).

To have an intuition of why the above theorem is valid, the reader may think of adding a dummy convex nonlinear convex function $c(x) = c_1(x) - 0$ to (B.1) and rely on the results from Sections 4 and 5. Indeed, by selecting a constant $M > 0$ large enough and function c such that $c(x) \leq -M < 0$ for all $x \in X$, we can see that Algorithm 1 applied to (B.1) with the additional and superfluous constraint $c(x) \leq 0$ boils down to Algorithm 2. Furthermore, in this artificial setting, the above convergence result follows directly from Theorem 5.8, item iii).

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Algorithm 2 Proximal Bundle Method for Convexly-Constrained CwC programs

Step 0 (Initialization) Let $x^0 \in X$, $\kappa \in (0, \frac{1}{2})$, $\kappa \leq \mu^0$, and $\text{To1} \geq 0$ be given.
 Call the oracles to compute $f_i(x^0)$ and (generalized) subgradients $s_{f_i}^0$, $i = 1, 2$.
 Define $k := \ell_k = 0$ and $\mathcal{B}_f^0 := \{0\}$.

Step 1 (Trial point) Compute x^{k+1} the (x -part) solution of the QP

$$\begin{cases} \min_{x, r} & r + \frac{\mu^k}{2} \|x - x^{\ell_k}\|^2 \\ \text{s.t.} & \bar{f}_1^j(x) - \bar{f}_2^{\ell_k}(x) - r \leq 0 \quad \forall j \in \mathcal{B}_f^k \\ & x \in X, \quad r \in \mathbb{R}. \end{cases}$$

Step 2 (Stopping test) If $\|x^{k+1} - x^{\ell_k}\| \leq \text{To1}$, then stop and return x^{ℓ_k} .

Step 3 (Oracles call) Compute $f_i(x^{k+1})$, and subgradients $s_{f_i}^{k+1}$, $i = 1, 2$.

Step 4 (Descent test)

(a) If $f(x^{k+1}) \leq f(x^{\ell_k}) - \frac{\kappa}{2} \|x - x^{\ell_k}\|^2$ holds, then declare a *serious step*: define $\ell_{k+1} := k + 1$, choose $\mathcal{B}_f^{k+1} \subset \{0, \dots, k + 1\}$ with $\{k + 1\} \in \mathcal{B}_f^{k+1}$ and arbitrarily select $\mu^{k+1} \in (0, \mu^k]$.

(b) Else, declare a *null step*: define $\ell_{k+1} := \ell_k$ and choose $\mathcal{B}_f^{k+1} \subset \{0, \dots, k + 1\}$ with $\bar{\mathcal{B}}_f^k \cup \{k + 1, \ell_k\} \subset \mathcal{B}_f^{k+1}$ ($\bar{\mathcal{B}}_f^k$ as in (4.10)).

Compute $\nu^k := 2 \max \left\{ \frac{\bar{f}_2^{\ell_k}(x^{k+1}) - f_2(x^{k+1})}{\|x^{k+1} - x^{\ell_k}\|^2}, 0 \right\}$. If $\nu^k \geq \mu^k - 2\kappa$, set $\mu^{k+1} = \nu^k + 1$; otherwise $\mu^{k+1} = \mu^k$.

Step 5 (Loop) Set $k := k + 1$ and go back to Step 1.

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